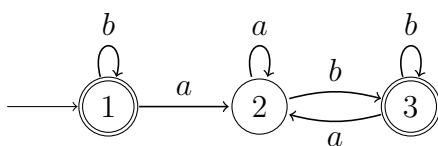


This exam consists of **four** exercises. The available points for each item are written in the margin. **Explain your answers!**

- [1] Determine whether the following sets over $\Sigma = \{a, b\}$ are regular or not. Prove your answers.
- (10) (a) $\{(a^n b)^m \mid n > m > 0\}$
- (10) (b) $\{x^n \mid \#a(x) \neq \#b(x) \text{ and } n > 0\}$

- (10) [2] (a) Construct a WSMO formula φ such that $L(\varphi) = L(M)$ for the following DFA M :



- (15) (b) Consider the WSMO formula $\varphi = \exists X. \exists y. y < x \wedge X(y)$. Give automata or regular expressions for the atomic subformulas and explain the operations needed to obtain the regularity of $L_a(\varphi)$.

- [3] Consider the LTL formula $\varphi = \neg(\neg a \text{ U X } b) \wedge \neg \text{G } b$.

- (10) (a) Transform φ into negation normal form ψ .
- (15) (b) Use the construction from the lecture to compute the alternating Büchi automaton A_ψ .
- (10) (c) Which of the following traces are accepted by A_ψ ?

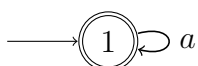
- i. $\{a\}\{b\}^\omega$ ii. $(\{a, b\} \emptyset \{a, b\})^\omega$ iii. \emptyset^ω

- (20) [4] Determine whether the following statements are true or false. (Providing explanation is optional.) Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

1. There exists a non-regular set that is accepted by an AFA.
2. MONA is a tool that implements algorithms for Büchi automata.
3. The class of ω -regular sets is effectively closed under intersection.
4. In a completely defined tree automaton, all states are productive.
5. Every Büchi automaton can be effectively transformed into a DBA.
6. Every regular set is accepted by a DFA having exactly one final state.
7. The automaton A_φ for the Presburger formula $\varphi: 3x - 2y = 1$ has at most 14 states.
8. The MSO formula $(\forall x. x = 0 \rightarrow X(x)) \wedge \neg(\forall x. X(x) \rightarrow \exists y. x < y \wedge X(y))$ is satisfiable.
9. For the WMSO formula $\varphi = \forall y. y < x \rightarrow X(y)$ we have

$$L_a(\varphi) = \binom{0}{0}^* \left[\binom{1}{0} + \binom{1}{1} \right] \left[\binom{0}{0} + \binom{0}{1} \right]^*$$

10. The equivalence relation \sim_M for the NBA M



over the alphabet $\Sigma = \{a, b\}$ has 2 equivalence classes.