Lastname: $\qquad$
Firstname: $\qquad$
Matriculation Number:

| Exercise | Points | Score |
| :---: | :---: | :---: |
| Program Analysis Including Modules and I/O | 20 |  |
| Programming with Lists | 32 |  |
| Datatypes and Higher-Order Functions | 26 |  |
| Evaluation and Types | 12 |  |
| $\sum$ | 90 |  |

- You have 90 minutes to solve the exercises.
- The exam consists of 4 exercises, for a total of 90 points (so there is 1 point per minute).
- The available points per exercise are written in the margin.
- Don't remove the staple (Heftklammer) from the exam.
- Don't write your solution in red color.
- Textual answers can be formulated in either English or German.

Exercise 1: Program Analysis Including Modules and I/O
Consider the following program. It chooses a random non-negative number and then asks the user to guess it.
module Main(main, finalize) where
import System.Random(randomIO) -- randomIO :: IO Integer
main $=$ do
num <- randomIO
putStrLn "Try to guess my number" guessingGame (abs num) 1
guessingGame :: Integer -> Integer -> IO String
guessingGame num $\mathrm{n}=$ do
str <- getLine
$\mathrm{x}<-$ (read str : : Integer)
let reason $=$ if $\mathrm{x}<\mathrm{num}$ then "small" else "large"
if $\mathrm{x}==$ num then finish n
else do
putStrLn \$ "Your guess was too " ++ reason ++ ". Try again" guessingGame num $n+1$
finish $\mathrm{n}=$ putStrLn $\$$ "You guessed my number using " ++ show n ++ " tries"
This program contains four mistakes that cause compilation errors.

- Identify these mistakes by providing line numbers,
- briefly explain the problem of each mistake, and
- explain how to correct the mistakes.

Note that all four mistakes are independent of one another.
(a) Mistake \#1

Solution: Line 1, finalize cannot be exported as it is not defined in this program. Either remove finalize or rename it to finish.
(b) Mistake \#2

Solution: Line 10, the type of guessingGame is wrong, it must be Integer -> Integer -> IO ().
(c) Mistake \#3

Solution: Line 13 , read is not a monadic operation, hence $\mathrm{x}<-$ read str must be replaced by let $\mathrm{x}=$ read str
(d) Mistake \#4

Solution: Line 18, the expression is parsed as (guessingGame num $n$ ) +1 , but it should have been guessingGame num $(n+1)$, so parentheses or a $\$$ are required.

## Exercise 2: Programming with Lists

Periodic functions can be represented by a list of initial values vs $=\left[v \_1, \ldots, v \_k\right]$ and a non-empty list of values ws $=\left[w \_1, \ldots, w \_m\right]$ that is repeated over and over again.
The function periodic : : [a] -> [a] -> Int $\rightarrow$ a is then defined as follows: periodic vs ws $n$ is the n-th element of the infinite list [v_1,..., v_k,w_1,...,w_m,w_1,...,w_m,w_1,...,w_m, ...] for every non-negative integer $n$. For example, if $f=$ periodic $[4,2][5,1,3]$ then the results of evaluating $f$ for arguments $0, \ldots, 8$ are shown in the following table:

| f 0 | f 1 | f 2 | f 3 | f 4 | f 5 | f 6 | f 7 | f 8 | $\ldots$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 4 | 2 | 5 | 1 | 3 | 5 | 1 | 3 | 5 | $\ldots$ |

Your task is to develop different (equivalent) implementations of periodic. You may freely use all Prelude functions. In particular, take, drop, splitAt, (!!), (++), filter, map and lookup might be useful.
(a) Define a function infList for the infinite list described above, i.e., infList vs ws should evaluate to $\left[\mathrm{v}_{-} 1, \ldots, \mathrm{v}_{-} \mathrm{k}, \mathrm{w}_{-} 1, \ldots, \mathrm{w}_{-} \mathrm{m}, \mathrm{w}_{-} 1, \ldots, \mathrm{w}_{-} \mathrm{m}, \ldots\right]$, and provide the most general type of infList. Further define a function periodicInf as an implementation of periodic that is based on infList.
(b) Define a function periodicN that implements periodic without constructing an infinite list and without using any predefined functions on lists (except for the list constructors). Evaluating periodicN vs ws n should require approximately n steps.
(c) Define a function periodicFast :: [a] -> [a] -> Int $\rightarrow$ a as an implementation of periodic. You should make use of the periodicity so that evaluating periodicFast vs ws $n$ does not need many more than length vs + length ws steps, even if $n$ is very large.
Hint: Haskell contains functions div and mod to compute the quotient and the remainder of an integer division respectively.
(d) Assume that a function $\mathrm{g}:$ : Int $\rightarrow$ a is periodic, i.e. there exists a finite list vs and a non-empty finite list ws such that $\mathrm{g}=$ periodic vs ws. Assume also that the list vs ++ ws is distinct, i.e., it contains no duplicates.
Define a function getLists : : Eq a => (Int $\rightarrow$ a) $->$ ([a], [a]) such that getLists g reconstructs vs and ws from g. In particular getLists (periodic vs ws) == (vs, ws) should be satisfied whenever vs and ws are two finite lists such that vs ++ ws is distinct and ws is non-empty.
Remark: you can get half of the points for this part if you instead implement an easier function getListSimple : : Eq a $\Rightarrow$ (Int $\rightarrow$ a) $\rightarrow$ [a]. Here we assume that the input is a periodic function $\mathrm{g}=$ periodic [] ws and only ws is computed via getListSimple. In particular, the property getListSimple (periodic [] ws) == ws should be satisfied whenever ws is a finite, non-empty and distinct list.

```
Solution:
infList :: [a] -> [a] -> [a]
infList vs ws = vs ++ infList ws ws
periodicInf vs ws n = infList vs ws !! n
periodicN (x : _) _ O = x
periodicN (_ : vs) ws n = periodicN vs ws (n - 1)
periodicN [] ws n = periodicN ws ws n
periodicFast vs ws n
    | n < lvs = vs !! n
    | otherwise = ws !! ((n - lvs) `mod` length ws)
    where lvs = length vs
getListSimple :: Eq a => (Int -> a) -> [a]
getListSimple g = g 0 : takeWhile (/= g 0) (map g [1..])
-- solution based on lookup / take / splitAt
getLists :: Eq a => (Int -> a) -> ([a], [a])
getLists g = search 1 where
    gis = map (\ i -> (f i, i)) [0..]
    search n = case lookup (g n) (take (n - 1) gis) of
            Just i -> let (vs, long) = splitAt i (map fst gis)
                in (vs, take (n - i) long)
        Nothing -> search (n + 1)
-- solution based on elem / take / span / filter
getLists :: Eq a => (Int -> a) -> ([a], [a])
getLists g = let
    xs = map g [0..]
    duplIndex = head (filter (\ i -> g i `elem` take i xs) [0..])
    vsWs = take duplIndex xs
    in span (/= g duplIndex) vsWs
```

Exercise 3: Datatypes and Higher-Order Functions
Consider the following program.

```
import Data.List(sort, sortOn)
-- sort :: Ord a => [a] -> [a]
-- sortOn :: Ord b => (a -> b) -> [a] -> [a]
-- sortOn f xs provides a sorted list ys of xs,
-- such that f (ys !! (i - 1)) <= f (ys !! i) for all 1 <= i < length ys;
-- sort and sortOn are closely related: sort = sortOn id
import Data.Char(toUpper)
-- toUpper :: Char -> Char
type Name = [String] -- a name might be composed, e.g., John Paul van de Boes
data Employee = Empl
    Name
    Int -- age
    Float -- salary
nameOf (Empl name _ _) = name
mapEmp f g h (Empl name age salary) = Empl (f name) (g age) (h salary)
```

(a) Write down the most general types of nameOf and of mapEmp.

```
Solution:
nameOf :: Employee -> Name
mapEmp ::
    (Name -> Name) ->
    (Int -> Int) ->
    (Float -> Float) ->
    (Employee -> Employee)
```

(b) Assume we want to write a function raiseSalary : : Employee -> Employee where the new salary is computed by the formula

$$
\text { new-salary }=\text { old-salary }+ \text { age } \times 10
$$

Further assume our implementation uses the following structure.
raiseSalary = mapEmp undefined undefined undefined
Either replace each undefined by a suitable $\lambda$-expression, or argue why raiseSalary cannot be implemented via mapEmp.

Solution: It cannot be implemented via mapEmp, since the new salary depends on the age and the old salary. However, in mapEmp the new salary is computed by a function that only gets the old salary as input.
(c) Assume we want to define a function toUpperEmployees : : [Employee] -> [Employee] that changes all names of all employees in a list so that they are written with uppercase letters. Choose a suitable implementation ( 4 points for the correct solution, 1 point for making no choice, 0 points for marking a wrong solution)
$\square$ toUpperEmployees $=$ toUpper
$\square$ toUpperEmployees $=\operatorname{map}(\operatorname{mapEmp} \quad(m a p ~ t o U p p e r))$
$\square$ toUpperEmployees $=\operatorname{map}(m a p E m p(m a p(m a p ~ t o U p p e r)) ~ i d ~ i d)$
$\square$ toUpperEmployees $=$ map (mapEmp (map toUpper) id id)
(d) Assume we want to define a function sortedUppercaseNames : : [Employee] -> [Name] that returns a sorted list of the names of all employees in a list converted to uppercase. The sorting should also be done using the uppercase names. There are four different attempts to implement sortedUppercaseNames (sun for brevity).

```
sun1 = map nameOf . sort . toUpperEmployees
sun2 = sort . toUpperEmployees . map nameOf
sun3 = map nameOf . sortOn nameOf . toUpperEmployees
sun4 = map nameOf . sortOn (nameOf . toUpperEmployees)
```

For each of the functions sun1, sun2, sun3 and sun4, indicate whether they are correct implementations of sortedUppercaseNames or not; and for the incorrect ones, give a brief description of the problem.

## Solution:

- sun1 does not compile, since Employee does not instantiate Ord.
- sun2 does not compile, since map nameOf produces a list of names, but toUpperEmployees expects a list of employees as input.
- sun3 is correct.
- sun4 does not compile for several reasons. For instance, toUpperEmployees requires a list of employees as input, but sortOn will invoke this function on singleton employees; or the result of toUpperEmployees, a list of employees, is passed to nameOf, which expects a singleton employee. Moreover, even if it compiled, the final result would not be in uppercase, as sortOn produces a permutation of the input list.


## Exercise 4: Evaluation and Types

In each multiple choice question, exactly one statement is correct. Marking the correct statement is worth 4 points, giving no answer counts 1 point, and marking multiple or a wrong statement results in 0 points.

Consider the following program.

```
foo \(\mathrm{n}=\) bar 01 n
bar x y n
    | \(\mathrm{n}=0=0=\mathrm{x}\)
    | otherwise \(=\) bar \(y(x+y)(n-1)\)
```

(a) What is the most general type of foo?foo : : Int -> Intfoo : : (Eq a, Num a) => a $->$ a
$\square$ foo : : (Eq a, Num a, Num b) => a $\rightarrow$ b
$\square$ foo : : (Eq a, Num a, Eq b, Num b) => a $->$ b
(b) What is the result of invoking foo n for some positive natural number n ?$\square+1+2+\ldots+n$
$\square \mathrm{n} * \mathrm{n}$
$\square$ The n-th element of the list of Fibonacci numbers $0,1,1,2,3,5,8,13,21, \ldots$None of the above
(c) Assume that we evaluate foo $\mathrm{n}:$ : Int for some positive n : : Int.

Choose the correct statement.The memory consumption is constant for both innermost evaluation and lazy evaluation.The memory consumption grows linearly in n for both innermost evaluation and lazy evaluation.The memory consumption grows linearly in n when using innermost evaluation, but is constant when using lazy evaluation.
$\square$ The memory consumption is constant when using innermost evaluation, but grows linearly in $n$ when using lazy evaluation.

