



# Functional Programming

## Week 2 – Tree Shaped Data and Datatypes

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## Last Lecture

- algorithm (can be informal) vs. program (concrete programming language)
- Haskell script (code, program, ...), e.g., `program.hs`  
`fahrenheitToCelsius f = (f - 32) * 5 / 9`  
consists of function definitions that describe input-output behaviour
- function- and parameter-names have to start with lowercase letters
- read-eval-print loop (REPL):  
load script, enter expressions and let these be evaluated

```
$ ghci program.hs
... welcome message ...
Main> fahrenheitToCelsius (3 + 20) - 7
-12.0
Main> ... further expressions ...
...
Main> :q
```

# Structured Data

# Different Representations of Data

- some (abstract) element can be represented in various ways

- example: numbers

- roman:
- decimal:
- binary:
- English:
- tally list:

XI  
11  
1011  
eleven  
|||||

- fact: algorithms depend on concrete representation

- example: addition

- decimal + binary: process digits of both numbers from right to left

$$\begin{array}{r} 7823 \\ + 909 \\ \hline 8732 \end{array}$$

- tally list: just write the two numbers side-by-side
- roman: algorithm?
- English: not well-suited

(||| + || = ||||)  
(IV + IX = XIII)  
(twenty-nine + two = thirty-one)

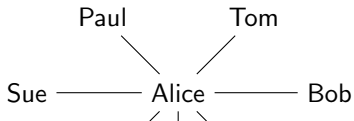
- in Haskell: numbers are built-in, representation not revealed to user

## Different Representations of Data – Continued

- representation must be chosen appropriately
- example: person
  - photographer:



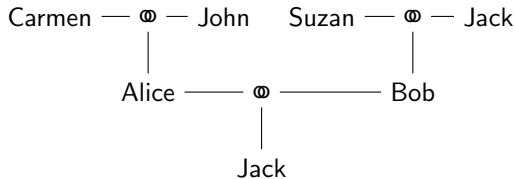
- social analysis:



- advertizing:

Bob (bob@foo.com, employee, hobbies: photography, jazz music, ...)

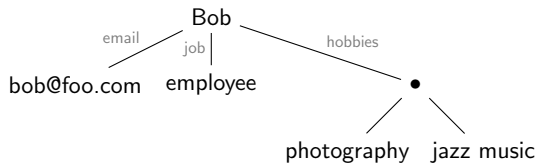
- genealogist:



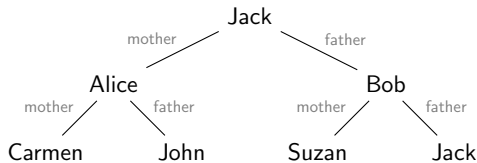
# Tree Shaped Data

- in functional programming most of the data is **tree shaped**
- a **tree**
  - has exactly one **root** node
  - can have several subtrees; nodes without subtrees are **leaves**
  - nodes and edges can be labeled
- in computer science, trees are usually displayed upside-down
- examples from previous slide

- advertizing:

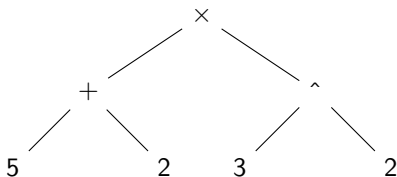


- genealogist:



# Expressions = Trees

- mathematical expressions can be represented as trees
- example
  - expression in textual form:  $(5 + 2) \times 3^2$
  - expression as tree

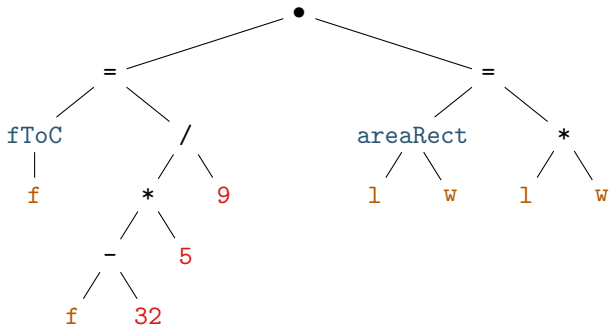


- remarks
  - the process of converting text into tree form is called **parsing**
  - operator precedences ( $\wedge$  binds stronger than  $\times$ , and  $\times$  binds stronger than  $+$ ) and parentheses are only required for parsing
    - parsing  $(5 + 2) \times (3^2)$  results in tree above
    - $5 + 2 \times 3^2$  and  $((5 + 2) \times 3)^2$  represent other trees
  - algorithm of calculator
    - convert textual input into tree
    - evaluate the tree bottom-up, i.e., start at leaves and end at root

# Programs = Trees

- programs can be represented as trees, too: **abstract syntax tree**
- example
  - program in textual form

```
-- some comment
fToC f = (f - 32) * 5 / 9
areaRect l w = l * w
```
  - abstract syntax tree (draft)



- comments and parentheses are no longer present in syntax tree



## Tree Shaped Data

- many programs deal with tree shaped data
- examples
  - calculator evaluates expression tree
  - compiler translates abstract syntax tree into machine code
  - search engine translates query into HTML (tree shaped)
  - contact application manages tree shaped personal data
  - file systems are organised as trees
- trees as **mental model** or representation of data is often suitable
- good news: processing tree shaped data is well-supported in functional programming
- next lecture: define functions on trees
- this lecture: restriction of trees via **types**

# Types

# Types

- functions are often annotated by their domain and codomain, e.g.,
  - $(!) : \mathbb{N} \rightarrow \mathbb{N}$
  - $(/) : \mathbb{R} \times (\mathbb{R} \setminus \{0\}) \rightarrow \mathbb{R}$
  - $\log_2 : \mathbb{R}_{>0} \rightarrow \mathbb{R}$
- domain and codomain provide useful information
  - domain: what are allowed inputs to a function
  - codomain: what are potential outputs of the function
- aim: specify domains and codomains of (Haskell-)functions
- notions
  - **elements** or **values**
    - maths: 5, 8,  $\pi$ ,  $-\frac{3}{4}$ , ...
    - Haskell: 5, 8, 3.141592653589793, -0.75, ..., "hello", 'c', ...
  - **sets of elements** to specify domain or codomain, in Haskell: **types**
    - maths:  $\mathbb{N}$ ,  $\mathbb{Z}$ ,  $\mathbb{Q}$ ,  $\mathbb{R}$ ,  $\mathbb{Q} \setminus \{0\}$ , ...
    - Haskell: **Integer** ( $\mathbb{Z}$ ), **Double** ( $\mathbb{R}$ ), **String**, **Char**, ...

# Typing Judgements

- in maths, we write statements like  $7 \in \mathbb{Z}$ ,  $7 \in \mathbb{R}$ ,  $0.75 \notin \mathbb{Z}$
- similarly in Haskell, we can express that a value or expression has a certain type via **typing judgements**
  - format: `expression :: type`
  - examples
    - `7 :: Integer` or `7 :: Double`
    - `'c' :: Char`
- that an expression indeed has the specified type is checked by the Haskell compiler
  - if an expression has not the given type, a type error is displayed
  - examples which raise an error
    - `7 :: String` or `0.75 :: Integer` or `'c' :: String`
    - `(7 :: Integer) :: Double`
  - remarks
    - unlike in maths where  $\mathbb{N} \subseteq \mathbb{Z} \subseteq \mathbb{Q}$ , in Haskell the types `Integer` and `Double` are **not** subtypes of each other
    - although some expressions can have both types (e.g., `7 + 5`), in general numbers of different types have to be converted explicitly
    - once a typing judgement is applied, the type of that expressions is fixed

## Typing of Haskell Expressions

- not only values but also functions have a type, e.g.,
  - `(/) :: Double -> Double -> Double`
  - `(+) :: Integer -> Integer -> Integer`
  - `(+) :: Double -> Double -> Double`
  - `head :: String -> Char`

### remarks

- a function can have multiple types, e.g., `(+)`
- limited expressivity, e.g. `(/) :: Double -> Double \ {0} -> Double` not allowed
- type checking enforces that in all function applications, type of arguments matches input-types of function
- example: consider expression `expr1 / expr2`
  - recall: `(/) :: Double -> Double -> Double`
  - it will be checked that both `expr1` and `expr2` have type `Double`
  - type of the overall expression `expr1 / expr2` will then be `Double`
- examples
  - `5 + 3 / 2`
  - `5 + '3'` or `5.2 + 0.8 :: Integer`



## Static Typing

- Haskell performs static typing
- **static** typing: types will be checked before evaluation  
(by contrast, dynamic typing checks types during evaluation)
- when loading Haskell script
  - check types of all function definitions `someFun x ... z = expr`:  
check that lhs `someFun x ... z` has same type as rhs `expr`
  - consequence: expressions cannot change their type during evaluation
- when entering expression in REPL: type check expression before evaluation
- benefits
  - no type checking required during evaluation
  - no type errors during evaluation

# Built-In Types – A First Overview

- numbers
  - `Integer` – arbitrary-precision integers
  - `Int` – fixed-precision integers with range at least  $\{-2^{28}, \dots, 2^{28} - 1\}$  (`-100`, `0`, `999`)
  - `Float` – single-precision floating-point numbers (`-12.34`, `5.78e36`)
  - `Double` – double-precision floating-point numbers
- characters and text
  - `Char` – a single character (`'a'`, `'Z'`, `' '`)
  - `String` – text of arbitrary length (`""`, `"a"`, `"The answer is 42."`)
  - some characters have to be **escaped** via the backslash-symbol `\`:
    - `'\t'` and `'\n'` – tabulator and new-line
    - `'\"'` and `'\''` – double- and single quote
    - `'\\'` – the backslash character
    - example: in the program

```
text = "Please say \"hello\"\n\nwhenever you enter the room"
```

the string `text` corresponds to the following two lines:  
Please say "hello"  
whenever you enter the room
- `Bool` – yes/no-decisions or truth-values (`True`, `False`)

# Datatypes



## Current State

- each value and function in Haskell has a type
- types are used to define input and output of function
- example: `fahrenheitToCelsius :: Double -> Double`
- built-in types for numbers, strings, and truth values
- missing: how to define types that describe tree shaped data?
- solution: definition of (algebraic) **datatypes**

## Datatype Definitions

- recall: a tree consists of a (labelled) root and 0 or more subtrees
- a **datatype** definition defines a set of trees by specifying all possible labelled roots together with a list of allowed subtrees
- Haskell scripts can contain many **datatype definitions** of the form

```
data TName =  
    CName1 type1_1 ... type1_N1  
  | ...  
  | CNameM typeM_1 ... typeM_NM  
deriving Show
```

where

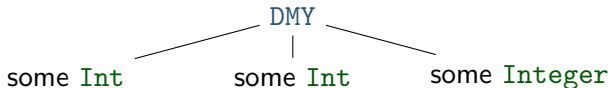
- **data** is a Haskell keyword to define a new datatype
- **TName** is the name of the new type; **type-names** always start with capital letters
- **CName1, ..., CNameM** are the labels of the permitted roots; these are called **constructors** and have to start with capital letters
- **typeI\_J** can be any Haskell type, including **TName** itself
- **|** is used as separator between different constructors
- **deriving Show** is required for displaying values of type **TName**

## Example Datatype Definition – Date

```
data Date = -- name of type
  DMY      -- name of constructor
    Int    -- day
    Int    -- month
    Integer -- year
```

### deriving Show

- here, there is only one constructor: `DMY`
- for day and month the precision of `Int` is sufficient
- the values of the type `Date` are exactly trees of the form



- in Haskell, these trees are built via the constructor `DMY`; `DMY` is a function of type `Int -> Int -> Integer -> Date` that is **not evaluated**
- example value of type `Date`: `DMY 16 10 2023`

## Example Datatype Definition – Person

```
data Person = -- name of type
  Person      -- constructor name can be same as type name
    String    -- first name
    String    -- last name
    Bool      -- married
    Date      -- birthday
```

### deriving Show

- reuse of previously defined types is permitted, in particular `Date`
- this leads to trees with more than one level of subtrees
- example program that defines a person (and an auxiliary date)

```
today = DMY 16 10 2023
```

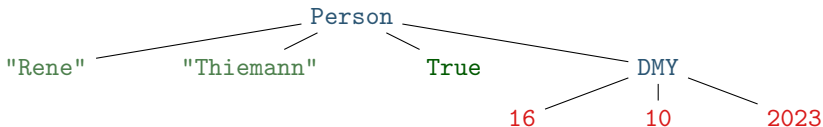
```
myself = Person "Rene" "Thiemann" True today
```

```
-- is the same as
```

```
myself = Person "Rene" "Thiemann" True (DMY 16 10 2023)
```

## Trees and Their Textual Representation

- in Haskell, trees have to be entered in a textual form, and trees are also output in textual form
- to define a tree with root constructor `C` and subtrees `t1`, ..., `tN`
  - one writes `C (t1) ... (tN)`;
  - if some `tI` is not a composed expression, then one can omit the parenthesis around `tI`;
  - this format is the same as for function applications
- example



- `Person "Rene" "Thiemann" True (DMY 16 10 2023)` ✓
- `Person "Rene" "Thiemann" (5 > 3) (DMY 16 (length "0123456789") 2023)` ✓
- `Person ("Rene", "Thiemann", True, DMY (16, 10, 2023))` ✗
- `Person "Rene" "Thiemann" True DMY 16 10 2023` ✗

## Example Datatype Definition – Vehicle

```
data Brand = Audi | BMW | Fiat | Opel deriving Show
```

```
data Vehicle =
```

```
    Car Brand Double -- horsepower
```

```
    | Bicycle
```

```
    | Truck Int -- number of wheels
```

```
deriving Show
```

- `Brand` just defines 4 car brands; all "trees" of type `Brand` consist of a single node; such datatypes are called **enumerations**
- there are three kinds of `Vehicles`, each having a different list of types
- example expressions of type `Vehicle`:

```
Car Fiat (60 + 1)
```

```
Car Audi 149.5
```

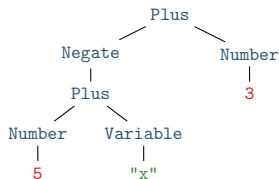
```
Bicycle
```

```
Truck (-7) -- types don't enforce all sanity checks
```

## Example Datatype Definition – Expr

```
data Expr =  
  Number Integer  
| Variable String  
| Plus Expr Expr  
| Negate Expr  
deriving Show
```

- type `Expr` models arithmetic expressions with addition and negation
- `Expr` is a **recursive** datatype: `Expr` is defined via `Expr` itself
- recursive datatypes contain values (trees) of arbitrary large height
  - expression  $-(5 + x) + 3$  in Haskell (as value of type `Expr`):  
`Plus (Negate (Plus (Number 5) (Variable "x"))) (Number 3)`
  - expression as tree



## Example Datatype Definition – Lists

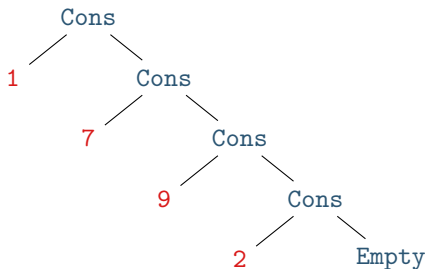
- lists are just a special kind of trees, e.g., lists of integers

```
data List =  
  Empty  
  | Cons Integer List  
deriving Show
```

- example representation of list [1, 7, 9, 2]

- in Haskell: `Cons 1 (Cons 7 (Cons 9 (Cons 2 Empty)))`

- as tree:





## Summary

- mental model: data = tree shaped data
- type = set of values; restricts shape of trees
- built-in types for numbers and strings
- user-definable datatypes, e.g., for expressions, lists, persons

```
data TName =  
    CName1 type1_1 ... type1_N1  
  | ...  
  | CNameM typeM_1 ... typeM_NM  
deriving Show
```

- next lecture: function definitions on trees