

- [1] (a) The set  $A = \{xay \mid x, y \in \{a\}^* \text{ and } x \neq y\}$  is regular since it consists of all strings over  $\{a\}^*$  having at least two  $a$ 's, so  $A = L(aaa^*)$ .
- (b) The set  $\{xy \mid x, y \in \Sigma^* \text{ and } x \neq y\}$  consists of all non-empty strings over  $\Sigma$  since we can take  $x = \epsilon$ . Since the empty string belongs to  $\{xy \mid x, y \in \Sigma^* \text{ and } x = y\}$ , it follows that the set  $B$  consists of all strings over  $\Sigma$ , which is obviously regular.

- [2] (a)  $\psi = \exists \ell. \neg P_a(\ell) \wedge \neg P_b(\ell) \wedge (\forall x. \neg P_a(x) \wedge \neg P_b(x) \rightarrow \ell \leq x) \wedge \exists X. X(0) \wedge (\forall x. \forall y. y = x + 1 \wedge y \leq \ell \rightarrow (X(x) \leftrightarrow \neg X(y))) \wedge X(\ell)$
- (b) First consider the subformula  $\psi: y < x \vee Y(y)$  with  $\text{FV}(\psi) = (x, y, Y)$ . We have

$$L_a(y < x) = \begin{pmatrix} 0 \\ 0 \end{pmatrix}^* \begin{pmatrix} 0 \\ 1 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \end{pmatrix}^* \begin{pmatrix} 1 \\ 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \end{pmatrix}^*$$

$$L_a(Y(y)) = \left[ \begin{pmatrix} 0 \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ 1 \end{pmatrix} \right]^* \begin{pmatrix} 1 \\ 1 \end{pmatrix} \left[ \begin{pmatrix} 0 \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ 1 \end{pmatrix} \right]^*$$

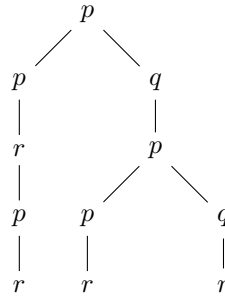
and compute

$$L_1 = \text{drop}_3^{-1}(L_a(y < x)) = \begin{pmatrix} 0 \\ 0 \\ * \end{pmatrix}^* \begin{pmatrix} 0 \\ 1 \\ * \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ * \end{pmatrix}^* \begin{pmatrix} 1 \\ 0 \\ * \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ * \end{pmatrix}^*$$

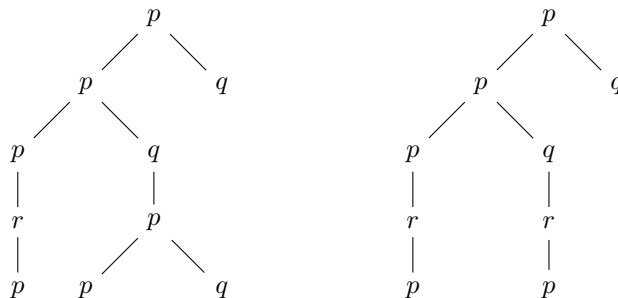
$$L_2 = \text{drop}_1^{-1}(L_a(Y(y))) = \left[ \begin{pmatrix} * \\ 0 \\ 0 \end{pmatrix} + \begin{pmatrix} * \\ 0 \\ 1 \end{pmatrix} \right]^* \begin{pmatrix} * \\ 1 \\ 1 \end{pmatrix} \left[ \begin{pmatrix} * \\ 0 \\ 0 \end{pmatrix} + \begin{pmatrix} * \\ 0 \\ 1 \end{pmatrix} \right]^*$$

We have  $L_a(\psi) = (L_1 \cup L_2) \cap L(\mathcal{A}_{2,1})$ . Here  $\mathcal{A}_{2,1}$  is a DFA to check the admissibility condition. For  $\varphi$  we obtain  $L_a(\varphi) = \text{stz}(\text{drop}_2(L_a(\psi)))$ .

- [3] (a) The string  $bbab$  does not belong to  $L(M)$  because after the first  $b$  the automaton is in state  $r$  but  $\Delta(r, b) = \perp$ . The string  $abab$  belongs to  $L(M)$  because of the accepting run



The string  $aaba$  does not belong to  $L(M)$  because the only runs



are not accepting.

(b) The ABA  $M = (Q, \{a, b\}, \Delta, s, \emptyset)$  with  $Q = \{s\} \cup \{a\} \times \{1, 2, 3\} \cup \{b\} \times \{1, 2\}$  and

$$\Delta(s, a) = (a, 2) \wedge (b, 2)$$

$$\Delta(s, b) = (a, 3) \wedge (b, 1)$$

$$\Delta((a, 3), a) = (a, 2)$$

$$\Delta((a, 2), a) = (a, 1)$$

$$\Delta((a, 1), a) = \top$$

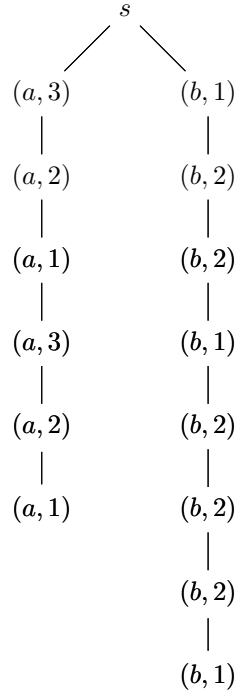
$$\Delta((a, 3), b) = \Delta((a, 2), b) = \Delta((a, 1), b) = (a, 3)$$

$$\Delta((b, 2), b) = (b, 1)$$

$$\Delta((b, 1), b) = \top$$

$$\Delta((b, 2), a) = \Delta((b, 1), a) = (b, 2)$$

accepts the given set. The following is an accepting run of  $M$  on the string  $baabaaab^\omega$ :



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1. True.
2. False.
3. True.
4. True.
5. False.
6. True.
7. True.
8. True.
9. True.
10. True.