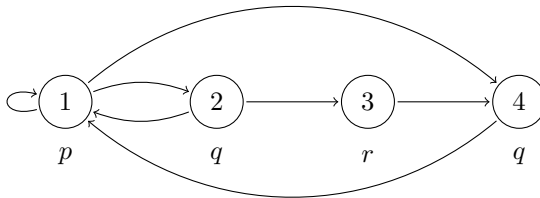


Solved exercises must be marked and solutions (as a single PDF file) uploaded in [OLAT](#). The (strict) deadline is 7 am on January 9.

## Exercises

- (3) 1. Consider the language  $L = \{x \in \{a, b\}^\omega \mid |x|_a = \infty \text{ and } |x|_b = \infty\}$  from [slide 9](#).
- (a) Show that no DBA with only two states accepts  $L$ .
  - (b) Give a minimal deterministic generalized Büchi automaton (GBA) accepting  $L$ .
  - (c) Apply the construction on [slide 11](#) to transform your GBA of (b) into an NBA.
- (2) 2. Determine which of the following equivalences hold. For every equivalence that does not hold give a counterexample.
- (a)  $G \varphi \equiv X G \varphi$
  - (b)  $F \neg(\varphi R \psi) \equiv \neg G \varphi U \neg G \psi$
  - (c)  $(\varphi \rightarrow \psi) W \chi \equiv \neg \varphi W \chi \vee \psi W \chi$
- (2) 3. Consider the following model  $\mathcal{M}$ :



- (a) For each of the LTL formulas  $\phi_1 = X(p U q)$ ,  $\phi_2 = G F p \wedge G F q \wedge G F r$ , and  $\phi_3 = G(X r \rightarrow q)$ , find a path  $\pi_i$  starting at state 1 such that  $\pi_i \models \phi_i$  for each  $i \in \{1, 2, 3\}$ .
  - (b) Give an LTL formula that distinguishes states 2 and 4 of  $\mathcal{M}$ .
- (3) 4. Let  $\mathcal{V} \in \{F, G\}^*$  be a sequence consisting of the temporal operators F and G. Show that  $F G \varphi \equiv \mathcal{V} F G \varphi$  and  $G F \varphi \equiv \mathcal{V} G F \varphi$ .

## Bonus Exercise

- (5) 5. Implement an algorithm that decides whether a given (hardcoded) DBA accepts a specific word  $u \cdot v^\omega$ , where  $u, v \in \Sigma^*$ . The submission should include the source code and a short description of how to run it.