ISR 2010

Introduction to Term Rewriting

lecture 9

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Sunday

introduction, examples, abstract rewriting, equational reasoning, term rewriting

Monday

termination, completion

Tuesday

completion, termination

Wednesday

confluence, modularity, strategies

Thursday

exam, advanced topics

Outline

Definition

- rewrite strategy $\mathcal{S}$ is mapping that assigns to every reducible term t nonempty set of finite nonempty rewrite sequences starting from t
- rewrite strategy $\mathcal{S}$ is deterministic if $\mathcal{S}(t)$ contains exactly one rewrite sequence for every reducible term t
- rewrite strategy $\mathcal{S}$ normalizes term t if there are no infinite $\mathcal{S}$ rewrite sequences starting from t
- rewrite strategy $\mathcal{S}$ is normalizing if it normalizes every term that has normal form
- rewrite strategy $\mathcal{S}$ is perpetual if every maximal $\mathcal{S}$ rewrite sequence starting from any non-terminating term is infinite

Lemma

for terminating TRSs every strategy is normalizing and perpetual

Example

- rewrite rules

$$
\begin{array}{llll}
0+0 \to 0 & 1+0 \to 1 & \cdots & 9+0 \to 9 \\
0+1 \to 1 & 1+1 \to 2 & \cdots & 9+1 \to 1:0 \\
0+2 \to 2 & 1+2 \to 3 & \cdots & 9+2 \to 1:1 \\
0+3 \to 3 & 1+3 \to 4 & \cdots & 9+3 \to 1:2 \\
0+4 \to 4 & 1+4 \to 5 & \cdots & 9+4 \to 1:3 \\
0+5 \to 5 & 1+5 \to 6 & \cdots & 9+5 \to 1:4 \\
0+6 \to 6 & 1+6 \to 7 & \cdots & 9+6 \to 1:5 \\
0+7 \to 7 & 1+7 \to 8 & \cdots & 9+7 \to 1:6 \\
0+8 \to 8 & 1+8 \to 9 & \cdots & 9+8 \to 1:7 \\
0+9 \to 9 & 1+9 \to 1:0 & \cdots & 9+9 \to 1:8
\end{array}
$$

$$x+(y:z) \to y:(x+z) \qquad 0:x \to x$$

$$(x:y)+z \to x:(y+z) \qquad x:(y:z) \to (x+y):z$$

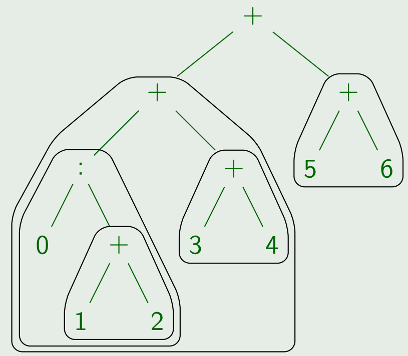
- term

$$((0:(1+2))+(3+4))+(5+6)$$

Example (cont'd)

term $\boxed{\boxed{0:\boxed{1+2}}+\boxed{3+4}}+\boxed{5+6}$

tree representation



parallel/leftmost outermost/innermost strategies

Example (cont'd)

- leftmost outermost strategy   12 redexes

$$((0:(1+2))+(3+4))+(5+6)$$
$$\to (0:((1+2)+(3+4)))+(5+6)$$
$$\to 0:(((1+2)+(3+4))+(5+6))$$
$$\to ((1+2)+(3+4))+(5+6) \to (3+(3+4))+(5+6)$$
$$\to (3+7)+(5+6) \to (1:0)+(5+6) \to 1:(0+(5+6))$$
$$\to 1:(0+(1:1)) \to 1:(1:(0+1)) \to (1+1):(0+1)$$
$$\to 2:(0+1) \to 2:1$$

- leftmost innermost strategy   10 redexes

$$(0:(1+2))+(3+4)+(5+6)$$
$$\to ((0:3)+(3+4))+(5+6) \to (3+(3+4))+(5+6)$$
$$\to (3+7)+(5+6) \to (1:0)+(5+6) \to (1:0)+(1:1)$$
$$\to 1:(0+(1:1)) \to 1:(1:(0+1)) \to 1:(1:1)$$
$$\to (1+1):1 \to 2:1$$

Example (cont'd)

- parallel outermost strategy   12 redexes in 9 steps

$$((0:(1+2))+(3+4))+(5+6)$$
$$\mathrel{+\!\!\!\!\!\rightarrow} (0:((1+2)+(3+4)))+(1:1)$$
$$\to 0:(((1+2)+(3+4))+(1:1))$$
$$\to ((1+2)+(3+4))+(1:1)$$
$$\to 1:(((1+2)+(3+4))+1)$$
$$\mathrel{+\!\!\!\!\!\rightarrow} 1:((3+7)+1) \to 1:((1:0)+1) \to 1:(1:(0+1))$$
$$\to (1+1):(0+1) \mathrel{+\!\!\!\!\!\rightarrow} 2:1$$

- parallel innermost strategy   10 redexes in 8 steps

$$((0:(1+2))+(3+4))+(5+6)$$
$$\mathrel{+\!\!\!\!\!\rightarrow} ((0:3)+7)+(1:1) \to (3+7)+(1:1)$$
$$\to (1:0)+(1:1) \to 1:(0+(1:1)) \to 1:(1:(0+1))$$
$$\to 1:(1:1) \to (1+1):1 \to 2:1$$

Definition

development of set of redex positions Q in term t is rewrite sequence starting from t in which all contracted redex positions descend from position in Q

Example

- rewrite rules

$$0 + y \to y \qquad 0 \times y \to 0$$
$$\mathsf{s}(x) + y \to \mathsf{s}(x + y) \qquad \mathsf{s}(x) \times y \to (x \times y) + y$$

- rewrite sequences

$$\underline{\mathsf{s}(0) \times (0 \times 0)} \to (0 \times \underline{(0 \times 0)}) + (0 \times 0) \to (0 \times 0) + (0 \times 0) \quad ☺$$
$$\underline{\mathsf{s}(0) \times (0 \times 0)} \to \underline{(0 \times (0 \times 0))} + (0 \times 0) \to (0 \times 0) + (0 \times 0) \quad ☹$$
$$\mathsf{s}(0) \times \underline{(0 \times 0)} \to \underline{\mathsf{s}(0) \times 0} \to (0 \times 0) + 0 \quad ☺$$

Theorem

developments are finite

Definition

development $A\colon s \to^* t$ of set of redex positions $Q \subseteq \mathcal{P}os(s)$ is complete if $Q \backslash A = \varnothing$

Lemma

for orthogonal TRSs: $s \to^* t$ *is complete development* $\iff$ $s \twoheadrightarrow\!\!\!\!\circ\;\, t$

Theorem

all complete developments of Q are permutation equivalent

Definition

for orthogonal TRSs full substitution strategy performs complete development of all redexes

Example

- rewrite rules

$$0 + y \to y \qquad 0 \times y \to 0$$
$$\mathsf{s}(x) + y \to \mathsf{s}(x + y) \qquad \mathsf{s}(x) \times y \to (x \times y) + y$$

- full substitution strategy

$$\begin{aligned}
&\mathsf{s}(\mathsf{s}(0)) \times (\mathsf{s}(0) + \mathsf{s}(\mathsf{s}(0))) \\
&\quad \twoheadrightarrow\!\!\!\!\circ\;\, (\mathsf{s}(0) \times \mathsf{s}(0 + \mathsf{s}(\mathsf{s}(0)))) + \mathsf{s}(0 + \mathsf{s}(\mathsf{s}(0))) \\
&\quad \twoheadrightarrow\!\!\!\!\circ\;\, ((0 \times \mathsf{s}(\mathsf{s}(\mathsf{s}(0)))) + \mathsf{s}(\mathsf{s}(\mathsf{s}(0)))) + \mathsf{s}(\mathsf{s}(\mathsf{s}(0))) \\
&\quad \to (0 + \mathsf{s}(\mathsf{s}(\mathsf{s}(0)))) + \mathsf{s}(\mathsf{s}(\mathsf{s}(0))) \\
&\quad \to \mathsf{s}(\mathsf{s}(\mathsf{s}(0))) + \mathsf{s}(\mathsf{s}(\mathsf{s}(0))) \to^+ \mathsf{s}(\mathsf{s}(\mathsf{s}(\mathsf{s}(\mathsf{s}(\mathsf{s}(0))))))
\end{aligned}$$

Outline

Theorem

for orthogonal TRSs

- *full substitution and parallel outermost strategies are normalizing*
- *innermost strategies are perpetual*
- *leftmost outermost strategy is not normalizing*

Example

$$\mathsf{a} \to \mathsf{b} \qquad \mathsf{c} \to \mathsf{c} \qquad \mathsf{f}(x, \mathsf{b}) \to \mathsf{b}$$

- leftmost outermost $\mathsf{f}(\mathsf{c}, \mathsf{a}) \to \mathsf{f}(\mathsf{c}, \mathsf{a}) \to \mathsf{f}(\mathsf{c}, \mathsf{a}) \to \cdots$
- leftmost innermost $\mathsf{f}(\mathsf{c}, \mathsf{a}) \to \mathsf{f}(\mathsf{c}, \mathsf{a}) \to \mathsf{f}(\mathsf{c}, \mathsf{a}) \to \cdots$
- parallel outermost $\mathsf{f}(\mathsf{c}, \mathsf{a}) \mathrel{+\!\!\!\!\!\longrightarrow} \mathsf{f}(\mathsf{c}, \mathsf{b}) \to \mathsf{b}$
- parallel innermost $\mathsf{f}(\mathsf{c}, \mathsf{a}) \mathrel{+\!\!\!\!\!\longrightarrow} \mathsf{f}(\mathsf{c}, \mathsf{b}) \to \mathsf{f}(\mathsf{c}, \mathsf{b}) \to \cdots$
- full substitution $\mathsf{f}(\mathsf{c}, \mathsf{a}) \mathrel{\circ\!\!\!\!\!\twoheadrightarrow} \mathsf{f}(\mathsf{c}, \mathsf{b}) \mathrel{\circ\!\!\!\!\!\twoheadrightarrow} \mathsf{b}$

Definition

TRS is left-normal if variables do not precede function symbols in left-hand sides

Example

- $x + (y : z) \to y : (x + z)$ ☹
- $(x : y) + z \to x : (y + z)$ ☺

Theorem

leftmost outermost strategy is normalizing for orthogonal left-normal TRSs

Remark

easy but important result: Combinatory Logic is left-normal

$$\mathsf{I}\,x \to x \qquad \mathsf{K}\,x\,y \to x \qquad \mathsf{S}\,x\,y\,z \to x\,z\,(y\,z)$$

Innsbruck

PhD Positions in Innsbruck

1. 3-year FWF project position (advisor: Dr. René Thiemann)
   - certification of termination methods in Isabelle/HOL
2. 4-year university position from October 1, 2010
3. 4-year university position from April 1, 2011

Requirements

- master degree in computer science or mathematics
- interest in computational logic – term rewriting – theorem proving

Further Information

- Aart Middeldorp

Outline

- Strategies
  - Definitions
  - Results
  - Optimality
- Strategy Annotations
- Further Reading

Observation

parallel outermost and full-substitution are not optimal because they perform useless steps

Example

- rewrite rules

$$0 + y \to y \qquad 0 \times y \to 0$$
$$\mathsf{s}(x) + y \to \mathsf{s}(x + y) \qquad \mathsf{s}(x) \times y \to (x \times y) + y$$

- rewrite sequence

$$(0 \times \mathsf{s}(0)) \times (0 + \mathsf{s}(0)) \twoheadrightarrow\!\!\!\!\!\!+ \; 0 \times \mathsf{s}(0) \to 0$$

- redex $0 + \mathsf{s}(0)$ is not needed

Definition

redex Δ in term t is needed if descendant of Δ is contracted in every rewrite sequence from t to normal form

Theorem

for orthogonal TRSs

- *every reducible term has needed redex*
- *needed reduction is normalizing*

Unfortunately

for orthogonal TRSs it is undecidable whether redex is needed

Remark

decidable approximations based on powerful tree automata techniques exist

Lemma

for left-normal orthogonal TRSs leftmost outermost redex is needed

Outline

- Strategies
  - Definitions
  - Results
  - Optimality
- Strategy Annotations
- Further Reading

Example

rewrite rules

$$\alpha:\ \mathsf{and}(x, \mathsf{T}) \to x \qquad \gamma:\ \mathsf{or}(\mathsf{T}, x) \to \mathsf{T} \qquad \epsilon:\ \infty \to \infty$$

$$\beta:\ \mathsf{and}(x, \mathsf{F}) \to \mathsf{F} \qquad \delta:\ \mathsf{or}(\mathsf{F}, x) \to x$$

strategy annotation A

$$A(\mathsf{and}) = [2, \alpha, \beta, 1] \qquad A(\mathsf{or}) = [1, \gamma, \delta, 2] \qquad A(\infty) = [\epsilon]$$

evaluation

$$\begin{aligned}\mathsf{or}(\mathsf{and}(\infty, \mathsf{F}), \mathsf{or}(\mathsf{T}, \infty)) &\to \mathsf{or}(\mathsf{F}, \mathsf{or}(\mathsf{T}, \infty)) \\ &\to \mathsf{or}(\mathsf{T}, \infty) \\ &\to \mathsf{T}\end{aligned}$$

Definitions

- strategy annotation for function symbol f is finite list $A(f)$ containing
  - argument positions of f
  - (labels of) rewrite rules for f
- strategy annotation $A(f)$ for function symbol f is full if $A(f)$ contains all argument positions of f and all rewrite rules for f
- strategy annotation $A(f)$ for function symbol f is in-time if argument positions are listed in $A(f)$ before rewrite rules that need them
- rewrite rule $f(s_1, \ldots, s_n) \to t$ needs argument position i if
  - s_i is non-variable
  - s_i is variable that appears in $s_1, \ldots, s_{i-1}, s_{i+1}, \ldots, s_n$

Example

rewrite rules

$$\alpha:\ \mathsf{if}(\mathsf{T}, x, y) \to x \qquad \beta:\ \mathsf{if}(\mathsf{F}, x, y) \to y \qquad \gamma:\ \mathsf{if}(z, x, x) \to x$$

strategy annotations

	full	in-time
$A_1(\mathsf{if}) = [1, \alpha, \beta, 2, 3, \gamma]$	☺	☺
$A_2(\mathsf{if}) = [\alpha, \beta, 1, 2, 3, \gamma]$	☺	☹
$A_3(\mathsf{if}) = [1, \alpha, \beta, \gamma, 2, 3]$	☺	☹
$A_4(\mathsf{if}) = [1, \alpha, \beta]$	☹	☺

Definition

$$\text{find-redex}_A(t) = \text{find-redex}_A(t, A(\text{root}(t)))$$

$$\text{find-redex}_A(t, [\,]) = \bot$$

$$\text{find-redex}_A(t, [\ell \to r \mid L]) = \begin{cases} (\epsilon, r\sigma) & \text{if } t = \ell\sigma \\ \text{find-redex}_A(t, L) & \text{otherwise} \end{cases}$$

$$\text{find-redex}_A(t, [i \mid L]) = \begin{cases} (ip, u) & \text{if find-redex}_A(t|_i) = (p, u) \\ \text{find-redex}_A(t, L) & \text{otherwise} \end{cases}$$

Theorem

$\forall$ *full strategy annotation* A $\quad\forall$ *term* t

$$\text{find-redex}_A(t) = \bot \iff t \text{ is normal form}$$

Example

rewrite rules

$$\alpha\colon\ \mathsf{and}(x,\mathsf{T}) \to x \qquad \gamma\colon\ \mathsf{or}(\mathsf{T},x) \to \mathsf{T} \qquad \epsilon\colon\ \infty \to \infty$$
$$\beta\colon\ \mathsf{and}(x,\mathsf{F}) \to \mathsf{F} \qquad \delta\colon\ \mathsf{or}(\mathsf{F},x) \to x$$

strategy annotation A

$$A(\mathsf{and}) = [\,2,\alpha,\beta,1\,] \qquad A(\mathsf{or}) = [\,1,\gamma,\delta,2\,] \qquad A(\infty) = [\,\epsilon\,]$$

evaluation

$$\text{find-redex}_A(\mathsf{or}(\mathsf{and}(\infty,\mathsf{F}),\mathsf{or}(\mathsf{T},\infty)))$$
$$= \text{find-redex}_A(\mathsf{or}(\mathsf{and}(\infty,\mathsf{F}),\mathsf{or}(\mathsf{T},\infty)),[\,1,\gamma,\delta,2\,]) = (1,\mathsf{F})$$
$$\text{find-redex}_A(\mathsf{and}(\infty,\mathsf{F}))$$
$$= \text{find-redex}_A(\mathsf{and}(\infty,\mathsf{F}),[\,2,\alpha,\beta,1\,])$$
$$= \text{find-redex}_A(\mathsf{and}(\infty,\mathsf{F}),[\,\alpha,\beta,1\,])$$
$$= \text{find-redex}_A(\mathsf{and}(\infty,\mathsf{F}),[\,\beta,1\,]) = (\epsilon,\mathsf{F})$$
$$\text{find-redex}_A(\mathsf{F}) = \bot$$

Example

rewrite rules

$$\alpha\colon\ \mathsf{if}(\mathsf{T},x,y) \to x \qquad \beta\colon\ \mathsf{if}(\mathsf{F},x,y) \to y \qquad \gamma\colon\ \mathsf{if}(z,x,x) \to x$$

strategy annotation

$$A(\mathsf{if}) = [\,1,\alpha,2,3,\gamma\,]$$

$$\text{find-redex}_A(\mathsf{if}(\mathsf{F},\mathsf{T},\mathsf{F})) = \bot \qquad \mathsf{if}(\mathsf{F},\mathsf{T},\mathsf{F}) \to \mathsf{T}$$

Definition

$$\text{strategy } \mathcal{S}_A(t) = \begin{cases} \varnothing & \text{if find-redex}_A(t) = \bot \\ \{t \to t[u]_p\} & \text{if find-redex}_A(t) = (p,u) \end{cases}$$

Corollary

∀ *full strategy annotation* A

$\mathcal{S}_A$ *is rewrite strategy*

Definition

$$\text{normalize}_A(t) = \text{normalize}_A(t, A(\text{root}(t)))$$
$$\text{normalize}_A(t,[\,]) = t$$
$$\text{normalize}_A(t,[\,\ell \to r \mid L\,]) = \begin{cases} \text{normalize}_A(r\sigma) & \text{if } t = \ell\sigma \\ \text{normalize}_A(t,L) & \text{otherwise} \end{cases}$$
$$\text{normalize}_A(t,[\,i \mid L\,]) = \text{normalize}_A(t[\text{normalize}_A(t|_i)]_i, L)$$

Theorem

∀ *in-time strategy annotation* A ∀ *term* t

$\mathcal{S}_A$ *normalizes* $t \iff$ $\text{normalize}_A(t)$ *is normal form of* t

JITty

`http://www.cwi.nl/~vdpol/jitty/`

Outline

- Strategies
- Strategy Annotations
- Further Reading

- Just-in-Time: On Strategy Annotations
  Jaco van de Pol
  Proc. 1st WRS, ENTCS 57, pp. 41–63, 2001
- A Sequential Reduction Strategy
  Sergio Antoy and Aart Middeldorp
  TCS 165(1), pp. 75 – 95, 1996