

Labelings for Decreasing Diagrams

Harald Zankl

Bertram Felgenhauer

Aart Middeldorp

Institute of Computer Science
University of Innsbruck
Austria

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Quiz



Wikimedia

Overview

- Preliminaries
- Decreasing Diagrams
- Labelings for Decreasing Diagrams
- Experiments
- Conclusion
- Announcement

Preliminaries

Definition (ARS)

$$\mathcal{A} = (A, \rightarrow)$$

Definition (confluence)

$$*\leftarrow \cdot \rightarrow * \subseteq \rightarrow * \cdot * \leftarrow$$

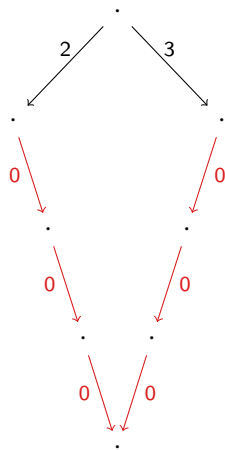
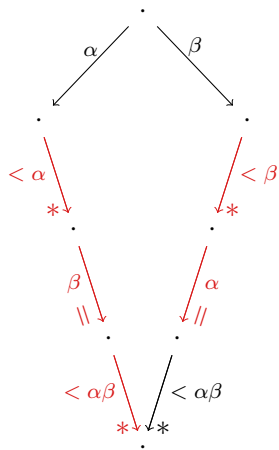
Definition (local confluence)

$$\leftarrow \cdot \rightarrow \subseteq \rightarrow * \cdot * \leftarrow$$

Theorem (van Oostrom, 1994)

every locally decreasing ARS is confluent

Decreasing Diagrams – Examples



Term Rewriting

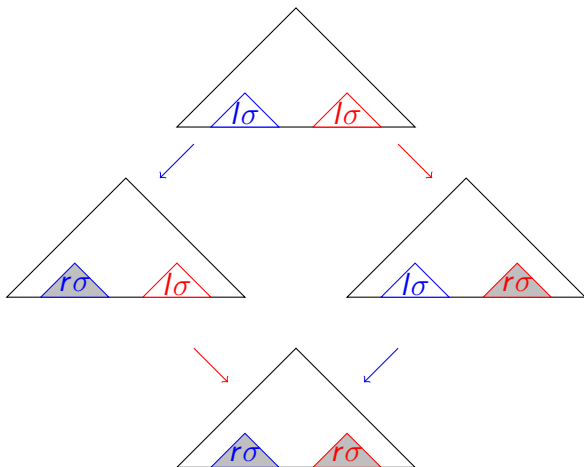
$$t \leftarrow s \rightarrow u$$

Three possibilities (modulo symmetry)

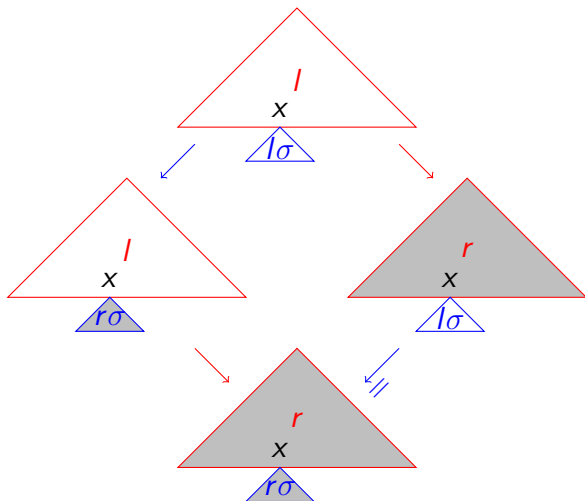
$$t \xrightarrow{p, l \rightarrow r} \leftarrow s[l\sigma]_p = s = s[l\sigma]_q \xrightarrow{l \rightarrow q, l \rightarrow r} u$$

- $p \parallel q$ (parallel overlap)
- $p < q$ and $q \geq \mathcal{P}\text{os}_V(s[l]_p)$ (variable overlap)
- $p \leq q$ and $q \in \mathcal{P}\text{os}_{\mathcal{F}}(s[l]_p)$ (critical overlap)

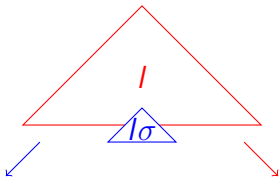
parallel



variable (linear)

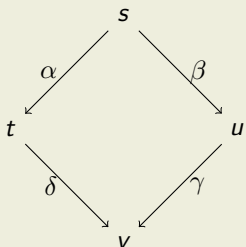


critical

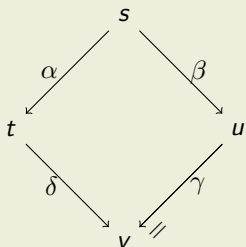


?

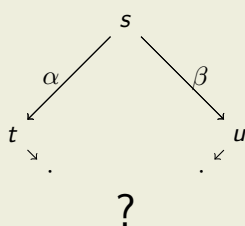
linear TRSs (3 kinds of local peaks)



(a) parallel



(b) variable linear



(c) critical

Definition (labeling)

$$\Gamma = s \rightarrow_{p,l} t \quad \Delta = u \rightarrow_{q,l'} v$$

$(l, \geq, >)$ labeling if $\geq \cdot > \cdot \geq \subseteq >$

- $l(\Gamma) \geq l(\Delta) \longrightarrow l(C[\Gamma\sigma]) \geq l(C[\Delta\sigma])$
- $l(\Gamma) > l(\Delta) \longrightarrow l(C[\Gamma\sigma]) > l(C[\Delta\sigma])$

L-labeling

Definition (L-labeling)

labeling ℓ is L-labeling if $\alpha \geq \gamma$, $\beta \geq \delta$ (parallel, variable linear)

Theorem

\mathcal{R} linear, critical diagrams decreasing for L-labeling $\longrightarrow \mathcal{R}$ confluent

Example

$(\ell_{rl}, \geq_{\mathbb{N}}, >_{\mathbb{N}})$ is L-labeling ($\ell_{rl}: \mathcal{R} \rightarrow \mathbb{N}$)

$(\ell_{sn}, \rightarrow_{\mathcal{R}}^*, \rightarrow_{\mathcal{S}/\mathcal{R}}^+)$ is L-labeling ($\ell_{sn}(s \rightarrow t) = s$)

if $\rightarrow_{\mathcal{S}} \subseteq \rightarrow_{\mathcal{R}}$ and \mathcal{S}/\mathcal{R} is terminating

Lemma

ℓ_1, ℓ_2 L-labelings $\longrightarrow \ell_1 \times \ell_2$ L-labeling

Example (van Oostrom, 2008)

1 : **nats** \rightarrow 0 : inc(nats)

2 : hd(x : y) \rightarrow x

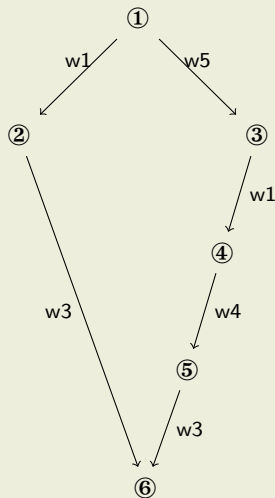
3 : tl(x : y) \rightarrow y

4 : inc(x : y) \rightarrow s(x) : inc(y)

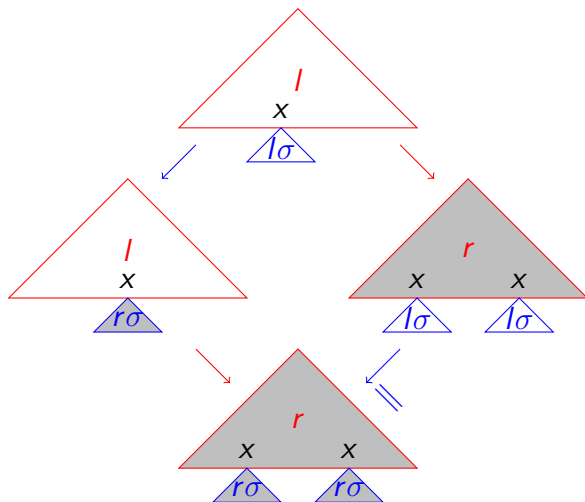
5 : inc(tl(**nats**)) \rightarrow tl(inc(nats))

- ① inc(tl(nats))
- ② inc(tl(0 : inc(nats)))
- ③ tl(inc(nats))
- ④ tl(inc(0 : inc(nats)))
- ⑤ tl(s(0) : inc(inc(nats)))
- ⑥ inc(inc(nats))

$w1 >_{\mathbb{N}} w3, w4$ shows confluence

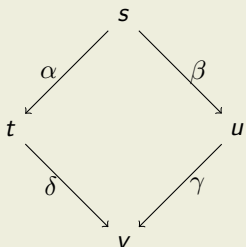


variable (left-linear)

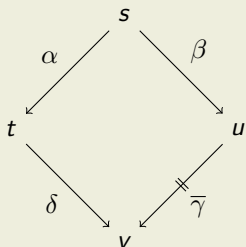


LL-labeling

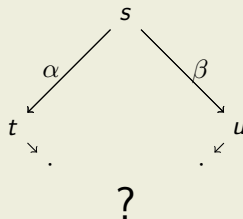
left-linear TRSs (3 kinds of local peaks)



(a) parallel



(b) variable left-linear



(c) critical

Definition (*LL*-labeling)

L-labeling ℓ is *LL*-labeling if $\alpha > \gamma_i$, $\beta \geq \delta$ (variable left-linear)

Theorem

\mathcal{R} left-linear, critical diagrams decreasing for LL-labeling $\longrightarrow \mathcal{R}$ confluent

Example

ℓ_{rl} is not LL-labeling

ℓ_{sn} is LL-labeling if $\mathcal{R}_d/\mathcal{R}_{nd}$ terminating

Definition (LL-labeling)

L-labeling ℓ is **weak** LL-labeling if $\alpha \geq \gamma_i, \beta \geq \delta$ (variable left-linear)

Example

ℓ_{rl} is weak LL-labeling

every LL-labeling is weak LL-labeling

Lemma

LL-labeling ℓ_1 , weak LL-labeling $\ell_2 \longrightarrow \ell_1 \times \ell_2, \ell_2 \times \ell_1$ LL-labelings

First result

Corollary

\mathcal{R} left-linear, $\mathcal{R}_d/\mathcal{R}_{nd}$ terminating, critical diagrams decreasing for weak LL-labeling $\longrightarrow \mathcal{R}$ confluent

Example (OO03 cont'd)

$$1: x + (y + z) \rightarrow (x + y) + z$$

$$2: (x + y) + z \rightarrow x + (y + z)$$

$$3: s(x) + y \rightarrow x + s(y)$$

$$4: x + s(y) \rightarrow s(x) + y$$

$$5: x \times s(y) \rightarrow x + (x \times y)$$

$$6: s(x) \times y \rightarrow (x \times y) + y$$

$$7: x + y \rightarrow y + x$$

$$8: x \times y \rightarrow y \times x$$

$$9: sq(x) \rightarrow x \times x$$

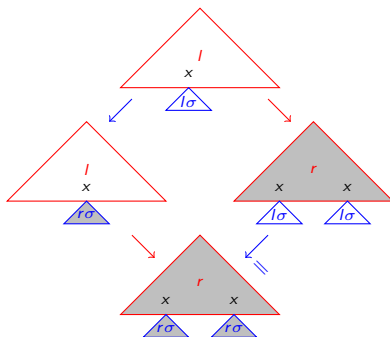
$$10: sq(s(x)) \rightarrow (x \times x) + s(x + x)$$

$\mathcal{R}_d/\mathcal{R}_{nd}$ terminating

$$+_N(x, y) = x + y \quad s_N(x) = x + 1 \quad \times_N(x, y) = x^2 + xy + y^2 \quad sq_N(x) = 3x^2 + 1$$

weak LL-labeling $l_{r1}(8) = l_{r1}(9) = 2, l_{r1}(6) = l_{r1}(10) = 1, l_{r1}(\cdot) = 0$

Towards a second result



Example

$1 : f(g(x, a)) \rightarrow g(f(x), f(x))$
 $2 : a \rightarrow b$
 $3 : b \rightarrow a$
 $4 : h(x) \rightarrow h(x)$

$\mathcal{R}_{>}^* = \{f_1(g_1(x)) \rightarrow g_1(f_1(x)), f_1(g_1(x)) \rightarrow g_2(f_1(x))\}$

$\mathcal{R}_{=}^* = \{h_1(x) \rightarrow h_1(x)\}$

Definition

$$\ell_\star(s \rightarrow_{p,l \rightarrow r} t) = \star(s, p)$$

Lemma

$\mathcal{R}_>^\star / \mathcal{R}_=^\star$ terminating $\longrightarrow (\ell_\star, \rightarrow_{\mathcal{R}_=^\star}^\star, \rightarrow_{\mathcal{R}_>^\star / \mathcal{R}_=^\star}^+) \text{ LL-labeling}$

Corollary

\mathcal{R} left-linear, $\mathcal{R}_>^\star / \mathcal{R}_=^\star$ terminating, ℓ weak LL-labeling, critical peaks decreasing for $\ell_\star \times \ell \longrightarrow \mathcal{R}$ confluent

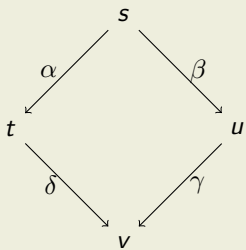
Example (OO03)

- | | |
|---|--|
| 1: $x + (y + z) \rightarrow (x + y) + z$ | 2: $(x + y) + z \rightarrow x + (y + z)$ |
| 3: $s(x) + y \rightarrow x + s(y)$ | 4: $x + s(y) \rightarrow s(x) + y$ |
| 5: $x \times s(y) \rightarrow x + (x \times y)$ | ... |

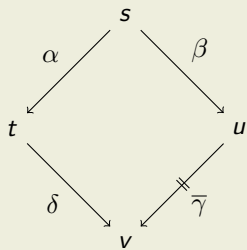
$\mathcal{R}_>^\star$ contains nonterminating $\times_1(x) \rightarrow +_2(\times_1(x))$

LL-labeling

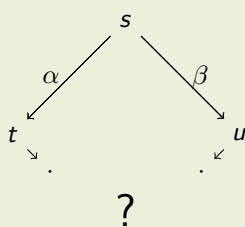
left-linear TRSs (3 kinds of local peaks)



(a) parallel



(b) variable left-linear



(c) critical

Definition (*LL*-labeling)

L-labeling ℓ is *LL*-labeling if $\alpha > \gamma_i$, $\beta \geq \delta$ (variable left-linear)

$$\alpha \geq \gamma_1, \alpha > \gamma_i \ (i > 1)$$

Example (OO03)

 $\mathcal{R}_{>}^*$

$$\text{sq}_1(s_1(x)) \rightarrow +_1(\times_1(x))$$

$$\text{sq}_1(s_1(x)) \rightarrow +_1(\times_2(x))$$

$$\text{sq}_1(s_1(x)) \rightarrow +_2(s_1(+_1(x)))$$

$$\text{sq}_1(s_1(x)) \rightarrow +_2(s_1(+_2(x)))$$

$$\text{sq}_1(x) \rightarrow \times_1(x)$$

$$\text{sq}_1(x) \rightarrow \times_2(x)$$

$$\times_2(y) \rightarrow +_1(\times_2(y))$$

$$\times_2(y) \rightarrow +_2(y)$$

$$\times_1(x) \rightarrow +_1(x)$$

$$\times_1(x) \rightarrow +_2(\times_1(x))$$

$$\times_1(x) \rightarrow \times_2(x)$$

$$\times_2(y) \rightarrow \times_1(y)$$

$$+_1(x) \rightarrow +_2(x)$$

$$+_2(y) \rightarrow +_1(y)$$

$$+_1(x) \rightarrow +_1(s_1(x))$$

$$+_1(x) \rightarrow +_1(+_1(x))$$

$$+_2(z) \rightarrow +_2(+_2(z))$$

$$+_2(y) \rightarrow +_2(s_1(y))$$

$$\times_1(x) \rightarrow +_2(\times_1(x))$$

 $\mathcal{R}_{=}^*$

$$\times_1(s_1(x)) \rightarrow +_1(\times_1(x))$$

$$\times_2(s_1(y)) \rightarrow +_2(\times_2(y))$$

$$+_1(s_1(x)) \rightarrow +_1(x)$$

$$+_1(+_1(x)) \rightarrow +_1(x)$$

$$+_1(+_2(y)) \rightarrow +_2(+_1(y))$$

$$+_2(+_2(z)) \rightarrow +_2(z)$$

$$+_2(s_1(y)) \rightarrow +_2(y)$$

$$+_2(+_1(y)) \rightarrow +_1(+_2(y))$$

$$\times_2(y) \rightarrow +_1(\times_2(y))$$

$$\text{sq}_{1\mathbb{N}}(x) = x + 2 \quad \times_{1\mathbb{N}}(x) = \times_{2\mathbb{N}}(x) = x + 1 \quad +_{1\mathbb{N}}(x) = +_{2\mathbb{N}}(x) = s_1(x) = x$$

Quiz



Google Maps

CSI

Experiments

53 left-linear TRSs (30 years of confluence literature)

method	pre	$CR(l_{rl})$	$CR(l_{sn})$	CR
rule labeling	40	35	–	35
$SN(\mathcal{R}_d/\mathcal{R}_{nd})$	45	40	37	42
$SN(\mathcal{R}_>^*/\mathcal{R}_=^*)$	45	42	34	42
$SN(\mathcal{R}_>^{**}/\mathcal{R}_=^{**})$	48	45	36	45
ACP				48
CSI				49

Conclusion

Summary

- decreasing diagrams
- (linear) $l_{rl}, l_{sn}, l_{\star}$
- (left-linear) $SN(\mathcal{R}_d/\mathcal{R}_{nd}), SN(\mathcal{R}_{>}^*/\mathcal{R}_{=}^*), SN(\mathcal{R}_{>}^{**}/\mathcal{R}_{=}^{**})$

Future Work

- study parallel reduction

Implementation

Label with \mathbb{N} (for lexicographic combination)

l_{rl} ✓ l_{sn} ✓ l_{\star} ✓

l_{rl}

(Aoto, 2010), (Hirokawa & Middeldorp, 2010)

l_{sn}

peak $t \leftarrow s \rightarrow u$

partial termination proof of $\text{CDS}(\mathcal{R})/\mathcal{R} \longrightarrow \mathcal{S}/\mathcal{R}'$

$l_{sn}(t_i \rightarrow t_{i+1}) = 1$ if $s \rightarrow_{\mathcal{S}}^* t_i$

$l_{sn}(t_i \rightarrow t_{i+1}) = 0$ otherwise

l_{\star}

prove termination of $\mathcal{R}_{>}^{\star\star}/\mathcal{R}_{=}^{\star\star}$ with matrix interpretations

Confluence Competition 2012



Takahito Aoto

Nao Hirokawa

Harald Zankl

`http://coco.nue.riec.tohoku.ac.jp`