

# Equational Reasoning for Termination of Rewriting

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# Overview

- Dependency Pairs
- Dependency Graph
- Waldmeister

## Definition (Dependency Pairs [AG00])

$$\text{DP}(\mathcal{R}) = \{l^\# \rightarrow u^\# \mid l \rightarrow r \in \mathcal{R}, r \triangleright u, \text{root}(u) \text{ defined}\}$$

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## Example

$$0 + y \rightarrow y$$

$$s(x) + y \rightarrow s(x + y)$$

$$0 \times y \rightarrow 0$$

$$s(x) \times y \rightarrow (x \times y) + y$$

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with DPs

$$s(x) +^\# y \rightarrow x +^\# y \quad s(x) \times^\# y \rightarrow (x \times y) +^\# y \quad s(x) \times^\# y \rightarrow x \times^\# y$$

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$$\begin{array}{ll} 0 + y \rightarrow y & 0 \times y \rightarrow 0 \\ s(x) + y \rightarrow s(x + y) & s(x) \times y \rightarrow (x \times y) + y \end{array}$$

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## Theorem (Termination)

$\mathcal{R}$  *terminating* iff no sequence  $s_0 \rightarrow_{\text{DP}(\mathcal{R})} t_0 \rightarrow_{\mathcal{R}}^* s_1 \rightarrow_{\text{DP}(\mathcal{R})} t_1 \rightarrow_{\mathcal{R}}^* \dots$

## Definition (Dependency Graph)

graph **DG** :=  $(N, E)$  with

- $N = \text{DP}(\mathcal{R})$
- $E = \{(s \rightarrow t, u \rightarrow v) \in N \times N \mid \exists \sigma, \tau : t\sigma \rightarrow_{\mathcal{R}}^* u\tau\}$



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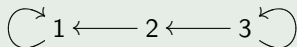
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$$1 : s(x) +^{\#} y \rightarrow x +^{\#} y \quad 2 : s(x) \times^{\#} y \rightarrow (x \times y) +^{\#} y \quad 3 : s(x) \times^{\#} y \rightarrow x \times^{\#} y$$



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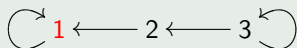
$$0 + y \geq y$$

$$0 \times y \geq 0$$

$$s(x) + y \geq s(x + y)$$

$$s(x) \times y \geq (x \times y) + y$$

$$1 : s(x) +^{\#} y > x +^{\#} y \quad 2 : s(x) \times^{\#} y \rightarrow (x \times y) +^{\#} y \quad 3 : s(x) \times^{\#} y \rightarrow x \times^{\#} y$$



## Theorem

$\forall \text{ SCC } \mathcal{S} \exists (\geq, >) \text{ s.t. } \mathcal{S} \cup \mathcal{R} \subseteq \geq \text{ and } \mathcal{S} \cap > \neq \emptyset \text{ then } \mathcal{R} \text{ terminating}$

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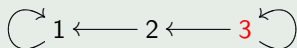
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## ren, cap [AG00]

- ren: rename variables (fresh variables)
- cap: replace term with defined root by fresh variable
- idea:  $\exists \sigma, \tau : t\sigma \rightarrow_{\mathcal{R}}^* u\tau$  implies  $\text{ren}(\text{cap}(t)) \stackrel{u}{\approx} u$

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$$\text{ren}( \quad F(h(x), h(x), x) \quad ) \stackrel{u}{\not\approx} F(0, 1, x)$$

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## Example

$$\mathcal{R} : \quad f(0, 1, x) \rightarrow f(h(x), h(x), x)$$

$$?: \quad \exists \sigma, \tau : F(h(x), h(x), x)\sigma \rightarrow_{\mathcal{R}}^* F(0, 1, x)\tau$$

$$F(h(y), h(z), w) \stackrel{u}{\not\approx} F(0, 1, x)$$

## Example

$$\mathcal{R} : \quad f(0, 1, x) \rightarrow f(h(x), h(x), x) \quad h(x) \rightarrow g(x)$$

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$$F(y, z, w) \stackrel{u}{\approx} F(0, 1, x)$$

ren, q $\delta$  [Mi02, HM05]

- q $\delta$ : replace term with defined root (wrt  $\mathcal{R}$ ) by fresh variable
- idea: use system  $\mathcal{R}$  to test  $\text{ren}(\text{cap}(t)) \stackrel{u}{\approx} u \ (\exists \sigma, \tau : t\sigma \rightarrow_{\mathcal{R}}^* u\tau)$   
idea: use system  $\mathcal{R}$  to test  $t \stackrel{u}{\approx} \text{ren}(\text{q}\delta(u)) \ (\exists \sigma, \tau : t\sigma \stackrel{*}{\leftarrow}_{\mathcal{R}} u\tau)$
- non-collapsing TRSs only

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$$\mathcal{R} : \quad f(0, 1, x) \rightarrow f(h(x), h(x), x) \quad h(x) \rightarrow g(x)$$

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- idea: use system  $\mathcal{R}$  to test  $t \stackrel{u}{\approx} \text{ren}(\text{q}\zeta(u)) \ (\exists \sigma, \tau : t\sigma \xrightarrow{\mathcal{R}}^* u\tau)$
- non-collapsing TRSs only

## Example

$$\mathcal{R} : \quad f(0, 1, x) \rightarrow f(h(x), h(x), x) \quad h(x) \rightarrow g(x)$$

$$F(h(x), h(x), x) \stackrel{u}{\not\approx} \text{ren}(\text{q}\zeta(F(0, 1, x)))$$

## Example

$$\mathcal{R} : \quad f(0, 1, x) \rightarrow f(h(x), h(x), x) \quad h(x) \rightarrow g(x)$$

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## ren, qso [Mi02, HM05]

- qso: replace term with defined root (wrt  $\mathcal{R}$ ) by fresh variable
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- idea: use system  $\mathcal{R}$  to test  $t \stackrel{u}{\approx} \text{ren}(\text{qso}(u)) \ (\exists \sigma, \tau : t\sigma \xrightarrow{\mathcal{R}}^* u\tau)$
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## Example

$$\mathcal{R} : \quad f(0, 1, x) \rightarrow f(h(x), h(x), x) \quad h(x) \rightarrow g(x)$$

$$F(h(x), h(x), x) \stackrel{u}{\not\approx} \text{ren}(F(0, 1, x))$$

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## Example

$$\mathcal{R} : \quad f(0, 1, x) \rightarrow f(h(x), h(x), x) \quad h(x) \rightarrow g(x)$$

$$F(h(x), h(x), x) \stackrel{u}{\not\approx} F(0, 1, y)$$

## Example

$$\mathcal{R} : \quad f(0, 1, x) \rightarrow f(h(a(x)), h(a(x)), x) \quad h(b(x)) \rightarrow 0 \quad h(b(x)) \rightarrow 1$$

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## tcap [GTS 05]

- $\text{tcap}(x) = y$  (fresh variable)
- $\text{tcap}(f(t_1, \dots, t_n)) = f(\text{tcap}(t_1), \dots, \text{tcap}(t_n))$   
if  $f(\text{tcap}(t_1), \dots, \text{tcap}(t_n))$  does not unify with lhs from  $\mathcal{R}$
- $\text{tcap}(f(t_1, \dots, t_n)) = y$  (fresh variable) otherwise

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$$F(h(x'), h(x'), x') \stackrel{u}{\approx} \text{ren}(q\sigma(F(0, 1, x')))$$

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$$F(y, z, w) \stackrel{u}{\approx} F(0, 1, x)$$

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## Equational Reasoning

- $t\sigma \rightarrow_{\mathcal{R}}^* u\tau$  implies  $t\sigma \leftrightarrow_{\mathcal{R}}^* u\tau$  means  $\mathcal{E}_{\mathcal{R}} \models t \approx u$
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- Waldmeister can refute  $\mathcal{E}_{\mathcal{R}} \models t \approx u$  (sometimes)



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$\implies$  yes, no, timeout

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$\implies$  yes, **no**, timeout

## Example (various/11)

$$f(0, 1, x) \rightarrow f(h(x), h(x), x)$$

$$h(0) \rightarrow 0$$

$$h(g(x, y)) \rightarrow y$$

with DPs

$$F(0, 1, x) \rightarrow F(h(x), h(x), x)$$

$$F(0, 1, x) \rightarrow H(x)$$

<b>NAME</b>	various11	
<b>MODE</b>	PROOF	
<b>SORTS</b>	ANY	
<b>SIGNATURE</b>	g: ANY ANY $\rightarrow$ ANY	0: $\rightarrow$ ANY
	h: ANY $\rightarrow$ ANY	1: $\rightarrow$ ANY
	F: ANY ANY ANY $\rightarrow$ ANY	f: ANY ANY ANY $\rightarrow$ ANY
<b>ORDERING</b>	LPO g > 1 > 0 > h > F > f	
<b>VARIABLES</b>	z, w, y, x : ANY	
<b>EQUATIONS</b>	$f(0, 1, x) = f(h(x), h(x), x)$ $h(0) = 0$ $h(g(x, y)) = y$ $F(x, x, x) = F(x, x, x)$	
<b>CONCLUSION</b>	$F(h(z), h(z), z) = F(0, 1, w)$	

```

Initial eqs: ( 1)  $x1 = h(g(x2, x1))$ 
                ( 2)  $h(0) = 0$ 
                ( 3)  $f(h(x1), h(x1), x1) = f(0, 1, x1)$ 
                ( 4)  $F(x1, x1, x1) = F(x1, x1, x1)$ 
                ( 5)  $eq(x1, x1) = true$ 
                ( 6)  $eq(F(h(x1), h(x1), x1), F(0, 1, x2)) = false$ 
Goals:      ( 1)  $true \neq false$ 

```

using narrowing to prove  $F(h(x1), h(x1), x1) \neq F(0, 1, x2)$

\*\*\*\*\* **COMPLETION – PROOF** \*\*\*\*\*

```

new rule:      1   $F(h(x1), h(x1), x1) ? F(0, 1, x2)$ 
new rule:      2   $h(g(x1, x2)) \rightarrow x2$ 
new rule:      3   $F(x1, x1, g(x2, x1)) ? F(0, 1, x3)$ 
new rule:      4   $h(0) \rightarrow 0$ 
new rule:      5   $F(0, 0, 0) ? F(0, 1, x1)$ 
new rule:      6   $f(h(x1), h(x1), x1) \rightarrow f(0, 1, x1)$ 
new rule:      7   $eq(x1, x1) \rightarrow true$ 
new rule:      8   $f(0, 0, 0) \rightarrow f(0, 1, 0)$ 
new equation:  1   $f(0, 1, g(x1, x2)) = f(x2, x2, g(x1, x2))$ 

```

**Refuted Goals:** No. 1:  $true \neq false$ , current  $true \neq false$

using narrowing to prove  $F(h(x1), h(x1), x1) \neq F(0, 1, x2)$

1 goal was specified, which was **refuted**.

**Waldmeister** states: System completed.

Experimental Results ( $T_1T_2$ )

	sccs	competition
edg***		
wdg		



# Experimental Results ( $T_1T_2$ )

	sccs	competition
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- 1391 TRSs from TPDB 5.0
- 5 sec timeout
- 0.1 sec timeout (per edge)

Experimental Results ( $T_1T_2$ )

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# Experimental Results ( $T_1T_2$ )

	sccs	competition
edg***	60	... not essential for competition
wdg	71	

- 1391 TRSs from TPDB 5.0
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# Experimental Results ( $T_1T_2$ )

	sccs	competition
edg***	60	but still ...
wdg	71	

- 1391 TRSs from TPDB 5.0
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Experimental Results ( $T_1T_2$ )

	sccs	competition
edg***	60	779
wdg	71	778

- 1391 TRSs from TPDB 5.0
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- 0.1 sec timeout (per edge)

# Experimental Results ( $T_1T_2$ )

	sccs	competition
edg***	60	779
wdg	71	778

- 1391 TRSs from TPDB 5.0
- 5 sec timeout
- 0.1 sec timeout (per edge)
- competition:
  - gain various/11, secret06/tpa01
  - lose two TRCSR/\*, various/14 (5 sec limit)
- AProVE, Jambox, TPA could (partly) solve these systems

# Conclusion

## Contribution

DP graph estimation based on equational reasoning

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## Related Work

- apart from [AG00, Mi02, HM05, GTS05]
- [KM09] → tree automata techniques
- [ZM09] → (increasing) interpretations



# Conclusion

## Contribution

DP graph estimation based on equational reasoning

## Related Work

- apart from [AG00, Mi02, HM05, GTS05] (and **innermost** refinements)
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# Conclusion

## Contribution

DP graph estimation based on equational reasoning

## Related Work

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## Future Work

- innermost termination (usable rules)
- other completion tools suitable?