

# Proving Equality of Simply Typed Term Rewriting Systems

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# Program Transformation by Templates

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Chiba et al. 2005, 2006, 2008, 2010.

- Pattern matching
- Correct templates
  - Equivalent Transformation of TRSs (Toyama 1991)
  - Inductionless induction
- RAPT
  - Automated Verification of the correctness of transformations

# Higher-order Program Transformation

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Based on 1st-order term rewriting

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fun sum xs = foldr (op +) 0 xs
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Based on 1st-order term rewriting

⇒ Difficult to apply higher-order programs (e.g. `map`, `filter`, `foldr` ...)

```
fun sum xs = foldr (op +) 0 xs
```

```
fun sum xs = sum1 xs 0
```

```
and sum1 [] y = y
```

```
  | sum1 (x::xs) y = sum1 xs (y+x);
```

# Simply Typed Term Rewriting System

(Yamada 2001)

$$(\text{sum } xs) \quad \rightarrow \quad (\text{foldr } + \ 0 \ xs)$$

$$(\text{foldr } f \ e \ []) \quad \rightarrow \quad e$$

$$(\text{foldr } f \ e \ (: \ x \ xs)) \rightarrow (f \ x \ (\text{foldr } f \ e \ xs))$$

$$(+ \ 0 \ y) \quad \rightarrow \quad y$$

$$(+ \ (s \ x) \ y) \quad \rightarrow \quad (s \ (+ \ x \ y))$$

$$\mathcal{F}^{\text{Nat}} = \{0\}, \quad \mathcal{F}^{\text{Nat} \rightarrow \text{Nat}} = \{s\}, \quad \mathcal{F}^{\text{Nat} \times \text{Nat} \rightarrow \text{Nat}} = \{+\}, \quad \mathcal{F}^{\text{List}} = \{[]\},$$

$$\mathcal{F}^{\text{Nat} \times \text{List} \rightarrow \text{List}} = \{:\}, \quad \mathcal{F}^{(\text{Nat} \times \text{Nat} \rightarrow \text{Nat}) \times \text{Nat} \times \text{List} \rightarrow \text{List}} = \{\text{foldr}\}, \quad \mathcal{F}^{\text{List} \rightarrow \text{Nat}} = \{\text{sum}\}$$

# Purpose

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Extending the framework to STTRS

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## Extending the framework to STTRS

1. Equivalent transformation
2. Transformation templates
3. Pattern matching
4. Correct templates



# Contents

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1. Introduction
2. Equivalent Transformation
3. Proof
4. Conclusion and Future works

# The equality of STTRSs

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$$\mathcal{R} \simeq_{\mathcal{G}} \mathcal{R}' \stackrel{\text{def}}{\iff} \forall s \in \mathsf{T}(\mathcal{G}). \forall t \in \mathsf{T}(\mathcal{C}). s \xrightarrow{*}_{\mathcal{R}} t \text{ iff } s \xrightarrow{*}_{\mathcal{R}'} t$$

# The equality of STTRSs

$$\mathcal{R} \simeq_{\mathcal{G}} \mathcal{R}' \stackrel{\text{def}}{\Leftrightarrow} \forall s \in \mathsf{T}(\mathcal{G}). \forall t \in \mathsf{T}(\mathcal{C}). s \xrightarrow{*}_{\mathcal{R}} t \text{ iff } s \xrightarrow{*}_{\mathcal{R}'} t$$

$$\left\{ \begin{array}{ll} (\text{sum } xs) & \rightarrow (\text{foldr } + \ 0 \ xs) \\ (\text{foldr } f \ e \ []) & \rightarrow e \\ (\text{foldr } f \ e \ (: \ x \ xs)) & \rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\ (+ \ 0 \ y) & \rightarrow y \\ (+ \ (\text{s } x) \ y) & \rightarrow (\text{s } (+ \ x \ y)) \end{array} \right\}$$

$$\left\{ \begin{array}{ll} (\text{sum } xs) & \rightarrow (\text{sum1 } xs) \\ (\text{sum1 } [] \ z) & \rightarrow z \\ (\text{sum1 } (: \ x \ xs) \ z) & \rightarrow (\text{sum1 } xs \ (+ \ z \ s)) \\ (+ \ 0 \ y) & \rightarrow y \\ (+ \ (\text{s } x) \ y) & \rightarrow (\text{s } (+ \ x \ y)) \end{array} \right\}$$

# The equality of STTRSs

$$\mathcal{R} \simeq_{\mathcal{G}} \mathcal{R}' \stackrel{\text{def}}{\Leftrightarrow} \forall s \in \mathsf{T}(\mathcal{G}). \forall t \in \mathsf{T}(\mathcal{C}). s \xrightarrow{*}_{\mathcal{R}} t \text{ iff } s \xrightarrow{*}_{\mathcal{R}'} t$$

$$\left\{ \begin{array}{l} (\text{sum } xs) \quad \rightarrow (\text{foldr } + \ 0 \ xs) \\ (\text{foldr } f \ e \ [ \ ]) \quad \rightarrow e \\ (\text{foldr } f \ e \ (: \ x \ xs)) \rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\ (+ \ 0 \ y) \quad \rightarrow y \\ (+ \ (\text{s } x) \ y) \quad \rightarrow (\text{s } (+ \ x \ y)) \end{array} \right\}$$

$$\simeq_{\{\text{sum}, +, [ \ ], : , \text{s}, 0\}}$$

$$\left\{ \begin{array}{l} (\text{sum } xs) \quad \rightarrow (\text{sum1 } xs) \\ (\text{sum1 } [ \ ] \ z) \quad \rightarrow z \\ (\text{sum1 } (: \ x \ xs) \ z) \rightarrow (\text{sum1 } xs \ (+ \ z \ s)) \\ (+ \ 0 \ y) \quad \rightarrow y \\ (+ \ (\text{s } x) \ y) \quad \rightarrow (\text{s } (+ \ x \ y)) \end{array} \right\}$$

# Equivalent Transformation

$\mathcal{R}_0$ : left-linear CS over  $\mathcal{F}_0$ ,  $\mathcal{E}$ : set of equations over  $\mathcal{F}_0$

## • Introduction

–  $(f \ x_1 \ \dots \ x_n)$  is linear

$$\mathcal{R}_k \xRightarrow{I} \mathcal{R}_k \cup \{(f \ x_1 \ \dots \ x_n) \rightarrow r\} (= \mathcal{R}_{k+1})$$

–  $f \notin \mathcal{F}_k$ ,

$$\mathcal{F}_{k+1} = \mathcal{F}_k \cup \{f\}$$

–  $r \in T(\mathcal{F}_k, \mathcal{V})$ , and

–  $x_1, \dots, x_n$  are basic

## • Addition

$$\mathcal{R}_k \xRightarrow{A} \mathcal{R}_k \cup \{l \rightarrow r\} (= \mathcal{R}_{k+1}) \quad - \quad l \overset{*}{\leftrightarrow}_{\mathcal{R}_k \cup \mathcal{E}} r$$

## • Elimination

$$\mathcal{R}_k \xRightarrow{E} \mathcal{R}_k \setminus \{l \rightarrow r\} (= \mathcal{R}_{k+1})$$

# Example

$$\left\{ \begin{array}{ll} (\text{sum } xs) & \rightarrow (\text{foldr } + \ 0 \ xs) \\ (\text{foldr } f \ e \ []) & \rightarrow e \\ (\text{foldr } f \ e \ (: \ x \ xs)) & \rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\ (+ \ 0 \ y) & \rightarrow y \\ (+ \ (\text{s } x) \ y) & \rightarrow (\text{s } (+ \ x \ y)) \end{array} \right\}$$

$$\left\{ \begin{array}{ll} (+ \ x \ (+ \ y \ z)) & \approx (+ \ (+ \ x \ y) \ z) \\ (+ \ x \ 0) & \approx x \end{array} \right\}$$

# Example

$$\left\{ \begin{array}{ll} (\text{sum } xs) & \rightarrow (\text{foldr } + \ 0 \ xs) \\ (\text{foldr } f \ e \ []) & \rightarrow e \\ (\text{foldr } f \ e \ (: \ x \ xs)) & \rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\ (+ \ 0 \ y) & \rightarrow y \\ (+ \ (\text{s } x) \ y) & \rightarrow (\text{s } (+ \ x \ y)) \end{array} \right\} \bullet \mathcal{R}_0 = \mathcal{R}_{\text{sum}}$$

$$\left\{ \begin{array}{l} (+ \ x \ (+ \ y \ z)) \approx (+ \ (+ \ x \ y) \ z) \\ (+ \ x \ 0) \approx x \end{array} \right\}$$

# Example

$$\left\{ \begin{array}{ll} (\text{sum } xs) & \rightarrow (\text{foldr } + \ 0 \ xs) \\ (\text{foldr } f \ e \ []) & \rightarrow e \\ (\text{foldr } f \ e \ (: \ x \ xs)) & \rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\ (+ \ 0 \ y) & \rightarrow y \\ (+ \ (\text{s } x) \ y) & \rightarrow (\text{s } (+ \ x \ y)) \end{array} \right\} \begin{array}{l} \bullet \mathcal{R}_0 = \mathcal{R}_{\text{sum}} \\ \bullet \mathcal{R}_0 \xrightarrow{I} \mathcal{R}_0 \cup \{(\text{sum1 } xs \ y) \rightarrow (+ \ y \ (\text{sum } xs))\} = \mathcal{R}_1 \end{array}$$

$$\left\{ \begin{array}{ll} (+ \ x \ (+ \ y \ z)) & \approx (+ \ (+ \ x \ y) \ z) \\ (+ \ x \ 0) & \approx x \end{array} \right\}$$



# Example

$$\left\{ \begin{array}{ll} (\text{sum } xs) & \rightarrow (\text{foldr } + \ 0 \ xs) \\ (\text{foldr } f \ e \ []) & \rightarrow e \\ (\text{foldr } f \ e \ (: \ x \ xs)) & \rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\ (+ \ 0 \ y) & \rightarrow y \\ (+ \ (\text{s } x) \ y) & \rightarrow (\text{s } (+ \ x \ y)) \end{array} \right\} \begin{array}{l} \bullet \mathcal{R}_0 = \mathcal{R}_{\text{sum}} \\ \bullet \mathcal{R}_0 \xRightarrow{I} \mathcal{R}_0 \cup \{(\text{sum1 } xs \ y) \rightarrow (+ \ y \ (\text{sum } xs))\} = \mathcal{R}_1 \\ \bullet \mathcal{R}_1 \xRightarrow{A} \mathcal{R}_1 \cup \{(\text{sum1 } [] \ y) \rightarrow y\} = \mathcal{R}_2 \end{array}$$

$$\left\{ \begin{array}{l} (+ \ x \ (+ \ y \ z)) \approx (+ \ (+ \ x \ y) \ z) \\ (+ \ x \ 0) \approx x \end{array} \right\}$$

# Example

$$\left\{ \begin{array}{ll}
 (\text{sum } xs) & \rightarrow (\text{foldr } + 0 xs) \\
 (\text{foldr } f e []) & \rightarrow e \\
 (\text{foldr } f e (: x xs)) & \rightarrow (f x (\text{foldr } f e xs)) \\
 (+ 0 y) & \rightarrow y \\
 (+ (s x) y) & \rightarrow (s (+ x y))
 \end{array} \right\}$$

$$\left. \begin{array}{l}
 \bullet \mathcal{R}_0 = \mathcal{R}_{\text{sum}} \\
 \bullet \mathcal{R}_0 \xRightarrow{I} \mathcal{R}_0 \cup \{(\text{sum1 } xs y) \rightarrow (+ y (\text{sum } xs))\} = \mathcal{R}_1 \\
 \bullet \mathcal{R}_1 \xRightarrow{A} \mathcal{R}_1 \cup \{(\text{sum1 } [] y) \rightarrow y\} = \mathcal{R}_2 \\
 \bullet \mathcal{R}_2 \xRightarrow{A} \mathcal{R}_2 \cup \{(\text{sum1 } (: x xs) y) \rightarrow (\text{sum1 } xs (+ y x))\} \\
 = \mathcal{R}_3
 \end{array} \right\}$$

$$\left\{ \begin{array}{l}
 (+ x (+ y z)) \approx (+ (+ x y) z) \\
 (+ x 0) \approx x
 \end{array} \right\}$$

$$\begin{aligned}
 & (\text{sum1 } (: x xs) y) \rightarrow_{\mathcal{R}_2} (+ y (\text{sum } (: x xs))) \\
 & \rightarrow_{\mathcal{R}_2} (+ y (\text{foldr } + 0 (: x xs))) \\
 & \rightarrow_{\mathcal{R}_2} (+ y (+ x (\text{foldr } + 0 xs))) \\
 & \approx (+ (+ y x) (\text{foldr } + 0 xs)) \\
 & \leftarrow_{\mathcal{R}_2} (\text{sum1 } xs (+ y x))
 \end{aligned}$$

# Example

$$\left\{ \begin{array}{ll}
 (\text{sum } xs) & \rightarrow (\text{foldr } + 0 xs) \\
 (\text{foldr } f e []) & \rightarrow e \\
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 (+ 0 y) & \rightarrow y \\
 (+ (s x) y) & \rightarrow (s (+ x y))
 \end{array} \right\}$$

- $\mathcal{R}_0 = \mathcal{R}_{\text{sum}}$
- $\mathcal{R}_0 \xRightarrow{I} \mathcal{R}_0 \cup \{(\text{sum1 } xs y) \rightarrow (+ y (\text{sum } xs))\} = \mathcal{R}_1$
- $\mathcal{R}_1 \xRightarrow{A} \mathcal{R}_1 \cup \{(\text{sum1 } [] y) \rightarrow y\} = \mathcal{R}_2$
- $\mathcal{R}_2 \xRightarrow{A} \mathcal{R}_2 \cup \{(\text{sum1 } (: x xs) y) \rightarrow (\text{sum1 } xs (+ y x))\} = \mathcal{R}_3$ 

$$\begin{aligned}
 & (\text{sum1 } (: x xs) y) \rightarrow_{\mathcal{R}_2} (+ y (\text{sum } (: x xs))) \\
 & \rightarrow_{\mathcal{R}_2} (+ y (\text{foldr } + 0 (: x xs))) \\
 & \rightarrow_{\mathcal{R}_2} (+ y (+ x (\text{foldr } + 0 xs))) \\
 & \approx (+ (+ y x) (\text{foldr } + 0 xs)) \\
 & \leftarrow_{\mathcal{R}_2} (\text{sum1 } xs (+ y x))
 \end{aligned}$$
- $\mathcal{R}_3 \xRightarrow{A} \mathcal{R}_3 \cup \{(\text{sum } xs) \rightarrow (\text{sum1 } xs 0)\} = \mathcal{R}_4$ 

$$\begin{aligned}
 & (\text{sum } xs) \leftarrow_{\mathcal{R}_3} (+ 0 (\text{sum } xs)) \\
 & \leftarrow_{\mathcal{R}_3} (\text{sum1 } xs 0)
 \end{aligned}$$

$$\left\{ \begin{array}{ll}
 (+ x (+ y z)) \approx (+ (+ x y) z) \\
 (+ x 0) \approx x
 \end{array} \right\}$$

# Example

$$\left\{ \begin{array}{l}
 (\text{sum } xs) \quad \rightarrow (\text{foldr } + \ 0 \ xs) \\
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 (\text{foldr } f \ e \ (: \ x \ xs)) \rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\
 (+ \ 0 \ y) \quad \rightarrow y \\
 (+ \ (\text{s } x) \ y) \quad \rightarrow (\text{s } (+ \ x \ y))
 \end{array} \right\} \bullet \begin{array}{l}
 \mathcal{R}_0 = \mathcal{R}_{\text{sum}} \\
 \mathcal{R}_0 \xRightarrow{I} \mathcal{R}_0 \cup \{(\text{sum1 } xs \ y) \rightarrow (+ \ y \ (\text{sum } xs))\} = \mathcal{R}_1 \\
 \mathcal{R}_1 \xRightarrow{A} \mathcal{R}_1 \cup \{(\text{sum1 } [] \ y) \rightarrow y\} = \mathcal{R}_2 \\
 \mathcal{R}_2 \xRightarrow{A} \mathcal{R}_2 \cup \{(\text{sum1 } (: \ x \ xs) \ y) \rightarrow (\text{sum1 } xs \ (+ \ y \ x))\} \\
 = \mathcal{R}_3 \\
 (\text{sum1 } (: \ x \ xs) \ y) \rightarrow_{\mathcal{R}_2} (+ \ y \ (\text{sum } (: \ x \ xs))) \\
 \rightarrow_{\mathcal{R}_2} (+ \ y \ (\text{foldr } + \ 0 \ (: \ x \ xs))) \\
 \rightarrow_{\mathcal{R}_2} (+ \ y \ (+ \ x \ (\text{foldr } + \ 0 \ xs))) \\
 \approx (+ \ (+ \ y \ x) \ (\text{foldr } + \ 0 \ xs)) \\
 \leftarrow_{\mathcal{R}_2} (\text{sum1 } xs \ (+ \ y \ x)) \\
 \bullet \mathcal{R}_3 \xRightarrow{A} \mathcal{R}_3 \cup \{(\text{sum } xs) \rightarrow (\text{sum1 } xs \ 0)\} = \mathcal{R}_4 \\
 (\text{sum } xs) \leftarrow_{\mathcal{R}_3} (+ \ 0 \ (\text{sum } xs)) \\
 \leftarrow_{\mathcal{R}_3} (\text{sum1 } xs \ 0) \\
 \bullet \mathcal{R}_4 \xRightarrow[E]{*} \mathcal{R}'
 \end{array}$$

$$\left\{ \begin{array}{l}
 (+ \ x \ (+ \ y \ z)) \approx (+ \ (+ \ x \ y) \ z) \\
 (+ \ x \ 0) \approx x
 \end{array} \right\}$$

# Example

$$\left\{ \begin{array}{l}
 (\text{sum } xs) \quad \rightarrow (\text{foldr } + \ 0 \ xs) \\
 (\text{foldr } f \ e \ []) \quad \rightarrow e \\
 (\text{foldr } f \ e \ (: \ x \ xs)) \rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\
 (+ \ 0 \ y) \quad \rightarrow y \\
 (+ \ (\text{s } x) \ y) \quad \rightarrow (\text{s } (+ \ x \ y))
 \end{array} \right\}$$

- $\mathcal{R}_0 = \mathcal{R}_{\text{sum}}$
- $\mathcal{R}_0 \xrightarrow{I} \mathcal{R}_0 \cup \{(\text{sum1 } xs \ y) \rightarrow (+ \ y \ (\text{sum } xs))\} = \mathcal{R}_1$
- $\mathcal{R}_1 \xrightarrow{A} \mathcal{R}_1 \cup \{(\text{sum1 } [] \ y) \rightarrow y\} = \mathcal{R}_2$
- $\mathcal{R}_2 \xrightarrow{A} \mathcal{R}_2 \cup \{(\text{sum1 } (: \ x \ xs) \ y) \rightarrow (\text{sum1 } xs \ (+ \ y \ x))\} = \mathcal{R}_3$ 

$$\begin{aligned}
 & (\text{sum1 } (: \ x \ xs) \ y) \rightarrow_{\mathcal{R}_2} (+ \ y \ (\text{sum } (: \ x \ xs))) \\
 & \rightarrow_{\mathcal{R}_2} (+ \ y \ (\text{foldr } + \ 0 \ (: \ x \ xs))) \\
 & \rightarrow_{\mathcal{R}_2} (+ \ y \ (+ \ x \ (\text{foldr } + \ 0 \ xs))) \\
 & \approx (+ \ (+ \ y \ x) \ (\text{foldr } + \ 0 \ xs)) \\
 & \leftarrow_{\mathcal{R}_2} (\text{sum1 } xs \ (+ \ y \ x))
 \end{aligned}$$

$$\left\{ \begin{array}{l}
 (+ \ x \ (+ \ y \ z)) \approx (+ \ (+ \ x \ y) \ z) \\
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$$\left\{ \begin{array}{l}
 (\text{sum } xs) \quad \rightarrow (\text{sum1 } xs) \\
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 (+ \ 0 \ y) \quad \rightarrow y \\
 (+ \ (\text{s } x) \ y) \quad \rightarrow (\text{s } (+ \ x \ y))
 \end{array} \right\}$$

- $\mathcal{R}_3 \xrightarrow{A} \mathcal{R}_3 \cup \{(\text{sum } xs) \rightarrow (\text{sum1 } xs \ 0)\} = \mathcal{R}_4$ 

$$\begin{aligned}
 & (\text{sum } xs) \leftarrow_{\mathcal{R}_3} (+ \ 0 \ (\text{sum } xs)) \\
 & \leftarrow_{\mathcal{R}_3} (\text{sum1 } xs \ 0)
 \end{aligned}$$
- $\mathcal{R}_4 \xrightarrow[E]{*} \mathcal{R}'$

# Verifying Equality of TRSs (1st-order)

## Theorem

- $\mathcal{R}$  is a left-linear CS over  $\mathcal{G}$
- $\mathcal{R}'$  is a TRS over  $\mathcal{G}'$
- $\mathcal{E}$  is a set of equations over  $\mathcal{G}$
- $\mathcal{R} \xrightarrow{I}^* \cdot \xrightarrow{A}^* \cdot \xrightarrow{E}^* \mathcal{R}'$  under  $\mathcal{E}$
- $\mathcal{R}, \mathcal{G} \vdash_{ind} \mathcal{E}$
- $CR(\mathcal{R}), SC(\mathcal{R}, \mathcal{G})$
- $SC(\mathcal{R}', \mathcal{G}')$

$$\implies \mathcal{R} \simeq_{\mathcal{G} \cap \mathcal{G}'} \mathcal{R}'$$

# Conjectures

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$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

1.  $\text{CR}(\mathcal{R})$  implies  $\text{CR}(\mathcal{R}_I)$
2. Let  $\mathcal{R}$  and  $\mathcal{R}_I$  be STTRSs over  $\mathcal{G}$  and  $\mathcal{F}$ .  $\forall s \in \text{T}(\mathcal{F}). \exists t \in \text{T}(\mathcal{G})$  s.t.  
 $s \xrightarrow[\mathcal{R}_I]{*} t$

# Proof

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$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

1.  $\overset{*}{\leftrightarrow}_{\mathcal{R}} = \overset{*}{\leftrightarrow}_{\mathcal{R}_I}$  on  $\mathbb{T}(\mathcal{G})$
2.  $\overset{*}{\leftrightarrow}_{\mathcal{R}_I} = \overset{*}{\leftrightarrow}_{\mathcal{R}_A}$  on  $\mathbb{T}(\mathcal{F})$
3.  $\overset{*}{\leftrightarrow}_{\mathcal{R}_I} = \overset{*}{\leftrightarrow}_{\mathcal{R}'}$  on  $\mathbb{T}(\mathcal{G}')$
4.  $\mathcal{R} \simeq_{\mathcal{G} \cap \mathcal{G}'} \mathcal{R}'$



# $\overset{*}{\leftrightarrow} \mathcal{R} = \overset{*}{\leftrightarrow} \mathcal{R}_I$ on $T(\mathcal{G})$

$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

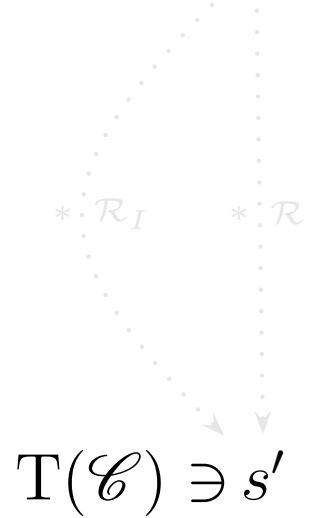
• Introduction

$$\mathcal{R}_k \xrightarrow[I]{*} \mathcal{R}_k \cup \{(f \ x_1 \ \dots \ x_n) \rightarrow r\} (= \mathcal{R}_{k+1})$$

$$\mathcal{F}_{k+1} = \mathcal{F}_k \cup \{f\}$$

- $(f \ x_1 \ \dots \ x_n)$  is linear
- $f \notin \mathcal{F}_k$ ,
- $r \in T(\mathcal{F}_k, \mathcal{V})$ , and
- $x_1, \dots, x_n$  are basic

$$T(\mathcal{G}) \ni s \xleftarrow[\mathcal{R}_I]{*} t \in T(\mathcal{G})$$



=



- $SC(\mathcal{R}, \mathcal{G})$
- $\mathcal{R} \subseteq \mathcal{R}_I$
- $CR(\mathcal{R}_I)$  and  $T(\mathcal{C}) \subseteq NF(\mathcal{R}_I)$

# $\overset{*}{\leftrightarrow} \mathcal{R} = \overset{*}{\leftrightarrow} \mathcal{R}_I$ on $T(\mathcal{G})$

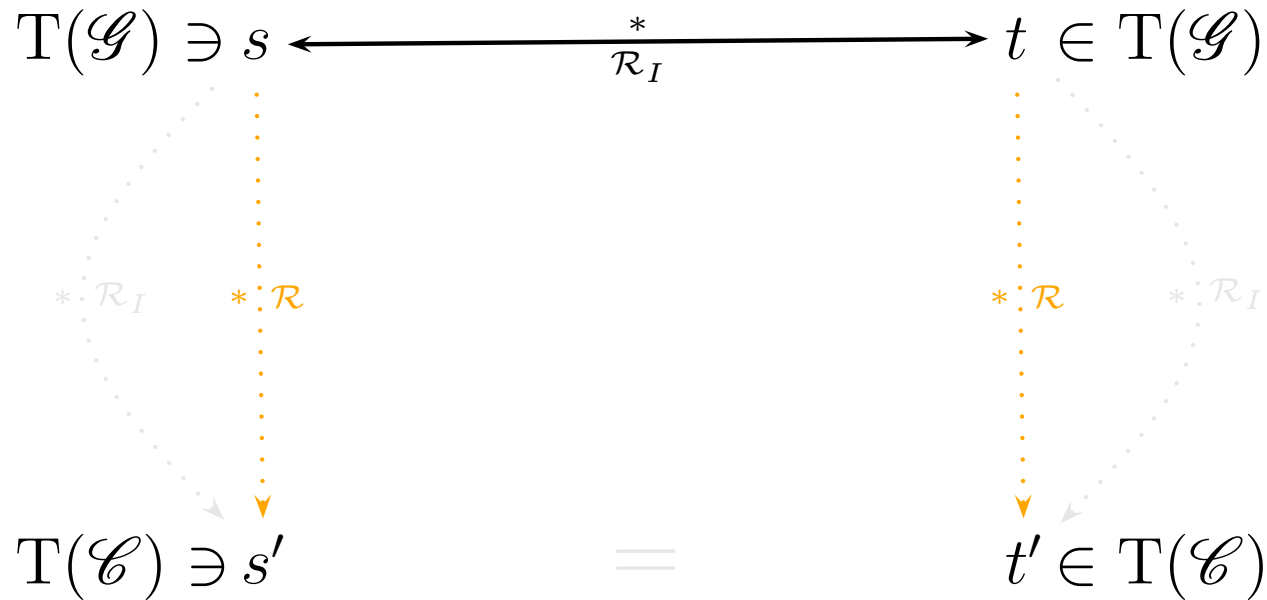
$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

• Introduction

$$\mathcal{R}_k \xrightarrow[I]{*} \mathcal{R}_k \cup \{(f \ x_1 \ \dots \ x_n) \rightarrow r\} (= \mathcal{R}_{k+1})$$

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- $(f \ x_1 \ \dots \ x_n)$  is linear
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- $SC(\mathcal{R}, \mathcal{G})$
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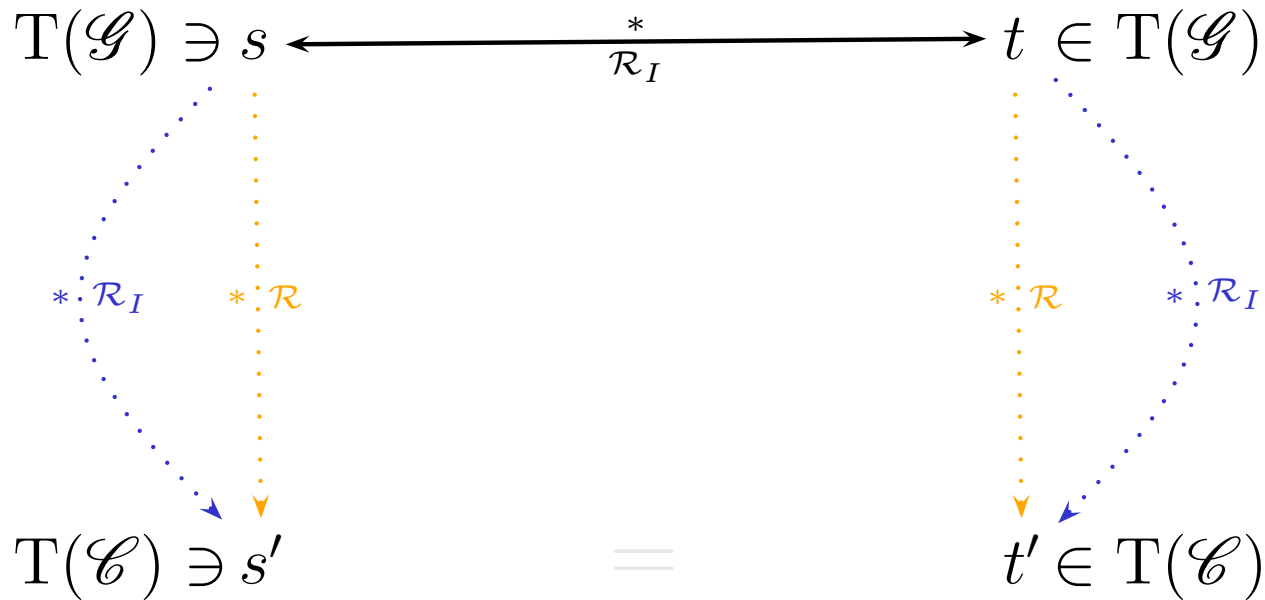
$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

• Introduction

$$\mathcal{R}_k \xrightarrow[I]{*} \mathcal{R}_k \cup \{(f \ x_1 \ \dots \ x_n) \rightarrow r\} (= \mathcal{R}_{k+1})$$

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- $SC(\mathcal{R}, \mathcal{G})$
- $\mathcal{R} \subseteq \mathcal{R}_I$
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# $\overset{*}{\leftrightarrow} \mathcal{R} = \overset{*}{\leftrightarrow} \mathcal{R}_I$ on $T(\mathcal{G})$

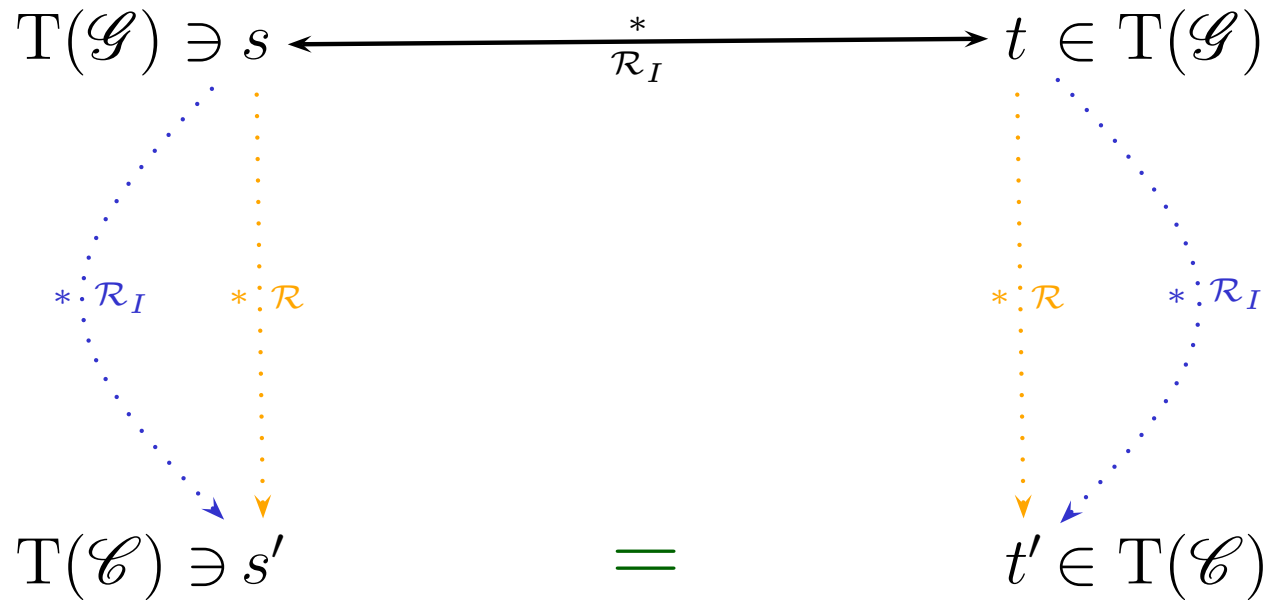
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• Introduction

$$\mathcal{R}_k \xrightarrow[I]{*} \mathcal{R}_k \cup \{(f \ x_1 \ \dots \ x_n) \rightarrow r\} (= \mathcal{R}_{k+1})$$

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- $SC(\mathcal{R}, \mathcal{G})$
- $\mathcal{R} \subseteq \mathcal{R}_I$
- $CR(\mathcal{R}_I)$  and  $T(\mathcal{C}) \subseteq NF(\mathcal{R}_I)$

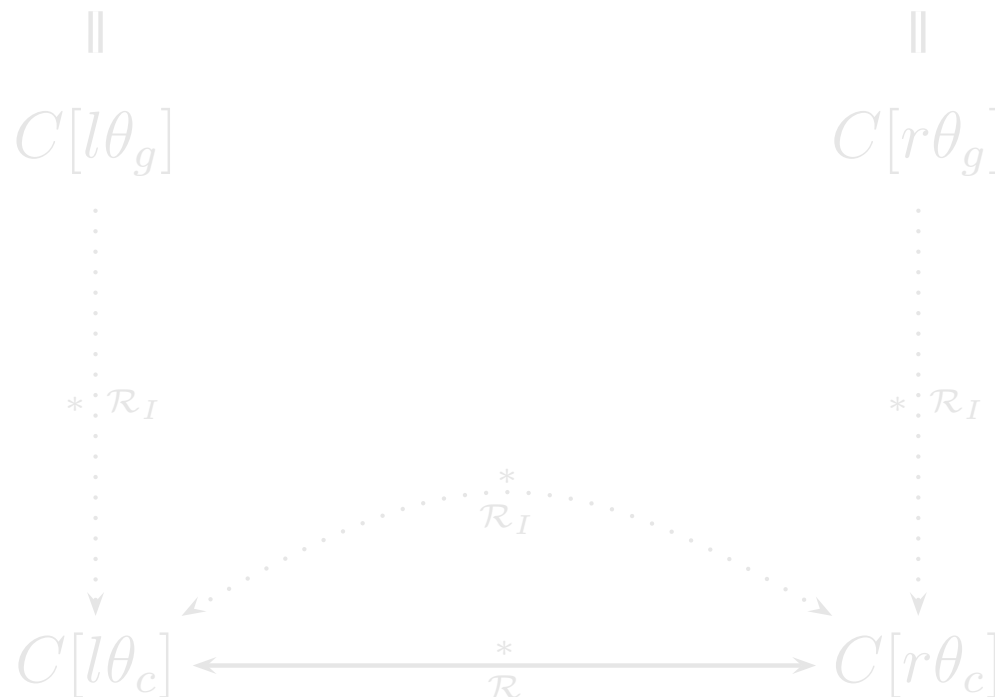
# $\overset{*}{\leftrightarrow} \mathcal{R}_I = \overset{*}{\leftrightarrow} \mathcal{R}_A$ on $T(\mathcal{F})$

$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

• Addition

$$\mathcal{R}_k \xrightarrow[A]{*} \mathcal{R}_k \cup \{l \rightarrow r\} (= \mathcal{R}_{k+1}) \quad - \quad l \overset{*}{\leftrightarrow}_{\mathcal{R}_k \cup \mathcal{E}} r$$

$$T(\mathcal{F}) \ni s \xleftarrow{\varepsilon} t \in T(\mathcal{F})$$



• The definition of reduction relation where

-  $\theta_g : \mathcal{V} \rightarrow T(\mathcal{F})$

•  $SC(\mathcal{R}_I, \mathcal{F})$  where

-  $\theta_c : \mathcal{V} \rightarrow T(\mathcal{C})$

•  $\mathcal{R}, \mathcal{G} \vdash_{ind} \mathcal{E}$

•  $\mathcal{R} \subseteq \mathcal{R}_I$

# $\overset{*}{\leftrightarrow} \mathcal{R}_I = \overset{*}{\leftrightarrow} \mathcal{R}_A$ on $T(\mathcal{F})$

$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

• Addition

$$\mathcal{R}_k \xrightarrow[A]{*} \mathcal{R}_k \cup \{l \rightarrow r\} (= \mathcal{R}_{k+1}) \quad - \quad l \overset{*}{\leftrightarrow}_{\mathcal{R}_k \cup \mathcal{E}} r$$

$$T(\mathcal{F}) \ni s \xleftarrow{\varepsilon} t \in T(\mathcal{F})$$



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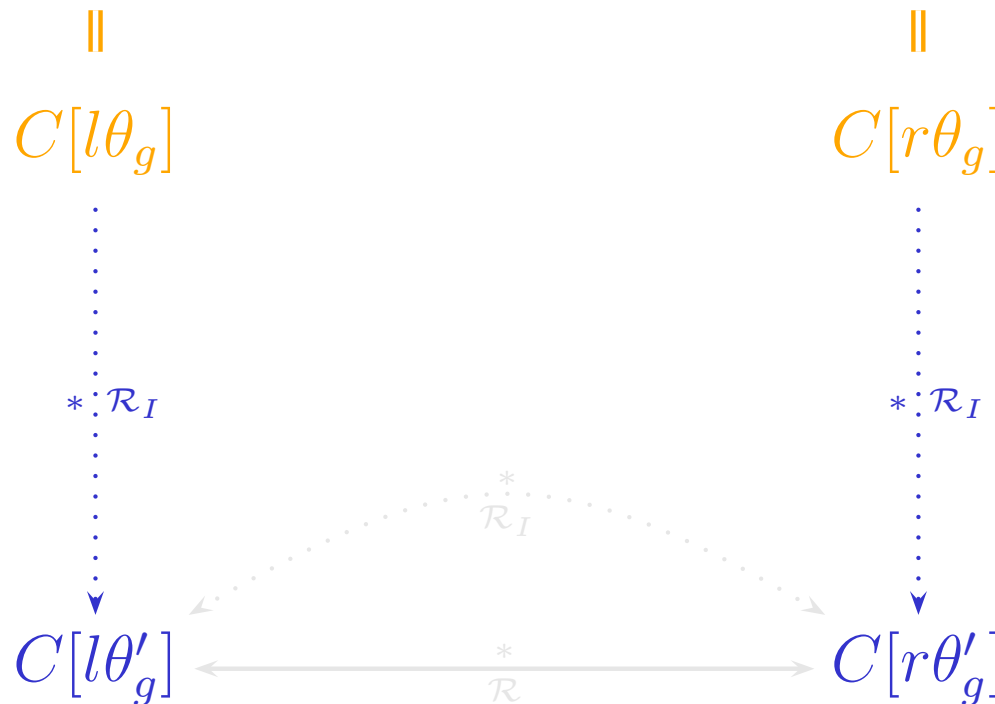
# $\overset{*}{\leftrightarrow} \mathcal{R}_I = \overset{*}{\leftrightarrow} \mathcal{R}_A$ on $T(\mathcal{F})$

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• The definition of reduction relation where

-  $\theta_g : \mathcal{V} \rightarrow T(\mathcal{F})$

• Conjecture 2 where

-  $\theta'_g : \mathcal{V} \rightarrow T(\mathcal{G})$

•  $\mathcal{R}, \mathcal{G} \vdash_{ind} \mathcal{E}$

•  $\mathcal{R} \subseteq \mathcal{R}_I$

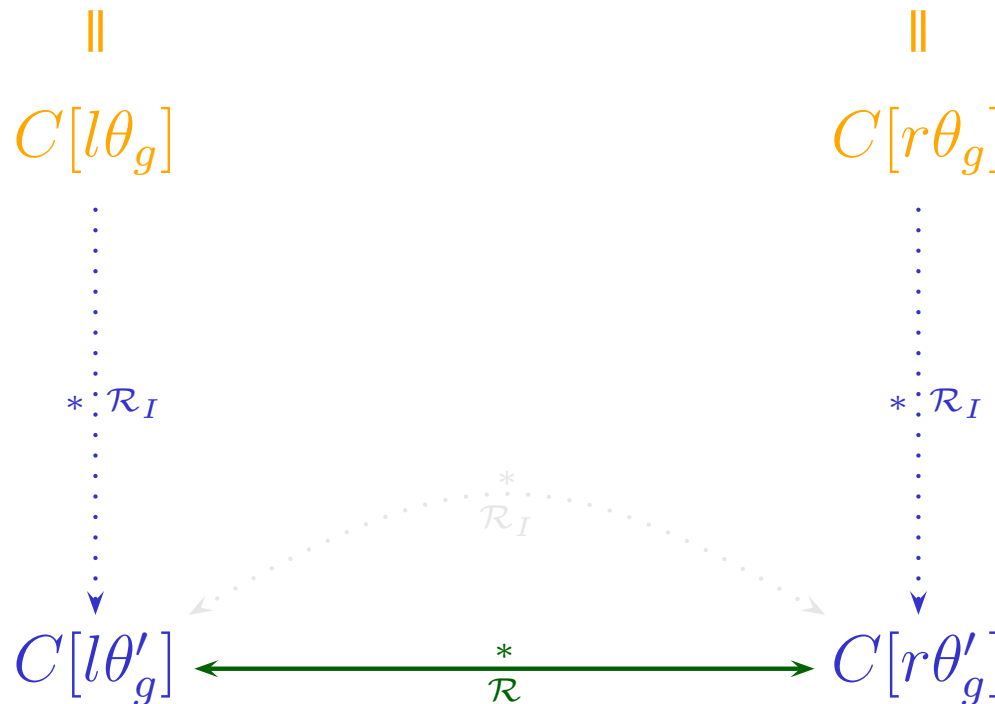
# $\overset{*}{\leftrightarrow} \mathcal{R}_I = \overset{*}{\leftrightarrow} \mathcal{R}_A$ on $T(\mathcal{F})$

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• Addition

$$\mathcal{R}_k \xrightarrow[A]{*} \mathcal{R}_k \cup \{l \rightarrow r\} (= \mathcal{R}_{k+1}) \quad - \quad l \overset{*}{\leftrightarrow}_{\mathcal{R}_k \cup \mathcal{E}} r$$

$$T(\mathcal{F}) \ni s \xleftarrow{\varepsilon} t \in T(\mathcal{F})$$



• The definition of reduction relation where

-  $\theta_g : \mathcal{V} \rightarrow T(\mathcal{F})$

• Conjecture 2 where

-  $\theta'_g : \mathcal{V} \rightarrow T(\mathcal{G})$

•  $\mathcal{R}, \mathcal{G} \vdash_{ind} \mathcal{E}$

•  $\mathcal{R} \subseteq \mathcal{R}_I$



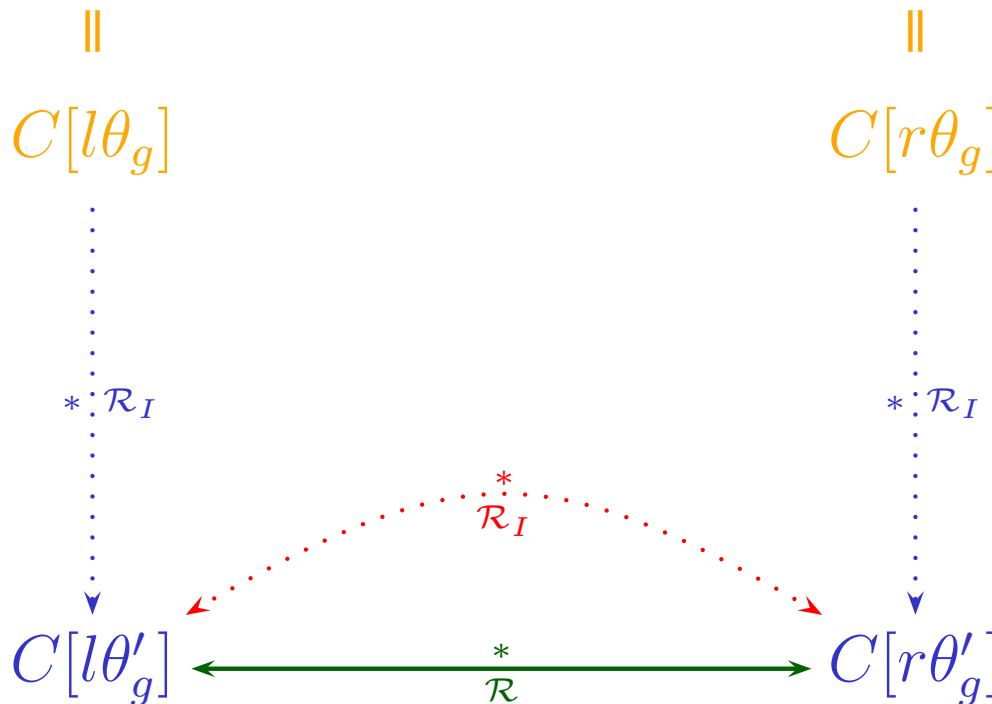
# $\overset{*}{\leftrightarrow} \mathcal{R}_I = \overset{*}{\leftrightarrow} \mathcal{R}_A$ on $T(\mathcal{F})$

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• Addition

$$\mathcal{R}_k \xrightarrow[A]{*} \mathcal{R}_k \cup \{l \rightarrow r\} (= \mathcal{R}_{k+1}) \quad - \quad l \overset{*}{\leftrightarrow}_{\mathcal{R}_k \cup \mathcal{E}} r$$

$$T(\mathcal{F}) \ni s \xleftarrow{\varepsilon} t \in T(\mathcal{F})$$



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-  $\theta_g : \mathcal{V} \rightarrow T(\mathcal{F})$

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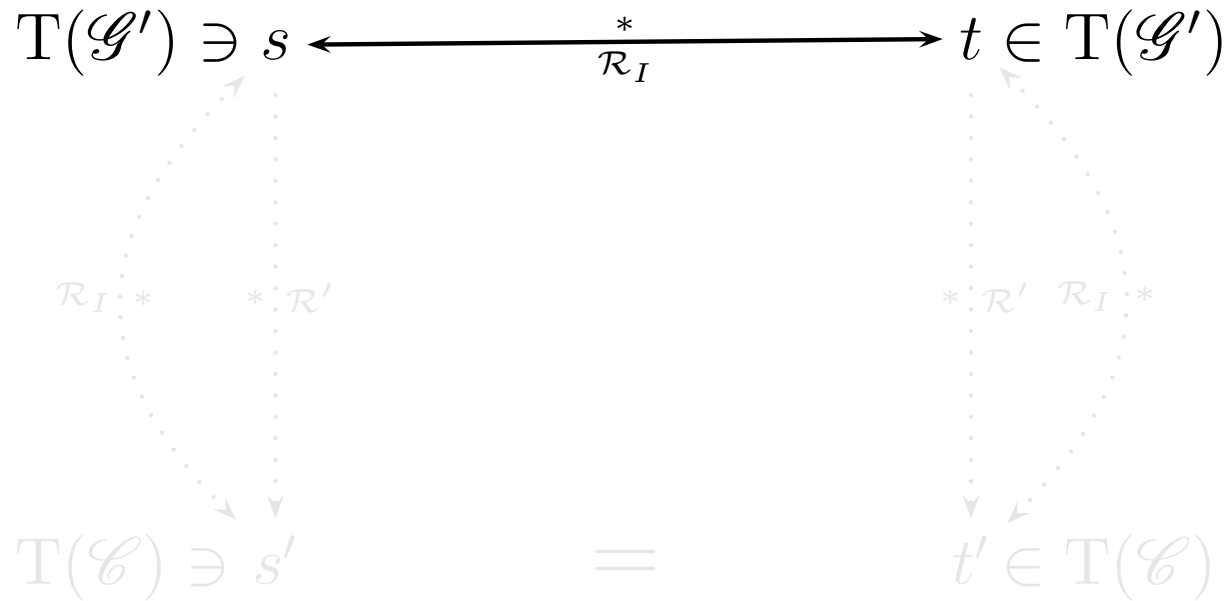
•  $\mathcal{R} \subseteq \mathcal{R}_I$

# $\overset{*}{\leftrightarrow} \mathcal{R}_I = \overset{*}{\leftrightarrow} \mathcal{R}'$ on $T(\mathcal{G}')$

$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

• Elimination

$$\mathcal{R}_k \xrightarrow[E]{*} \mathcal{R}_k \setminus \{l \rightarrow r\}$$



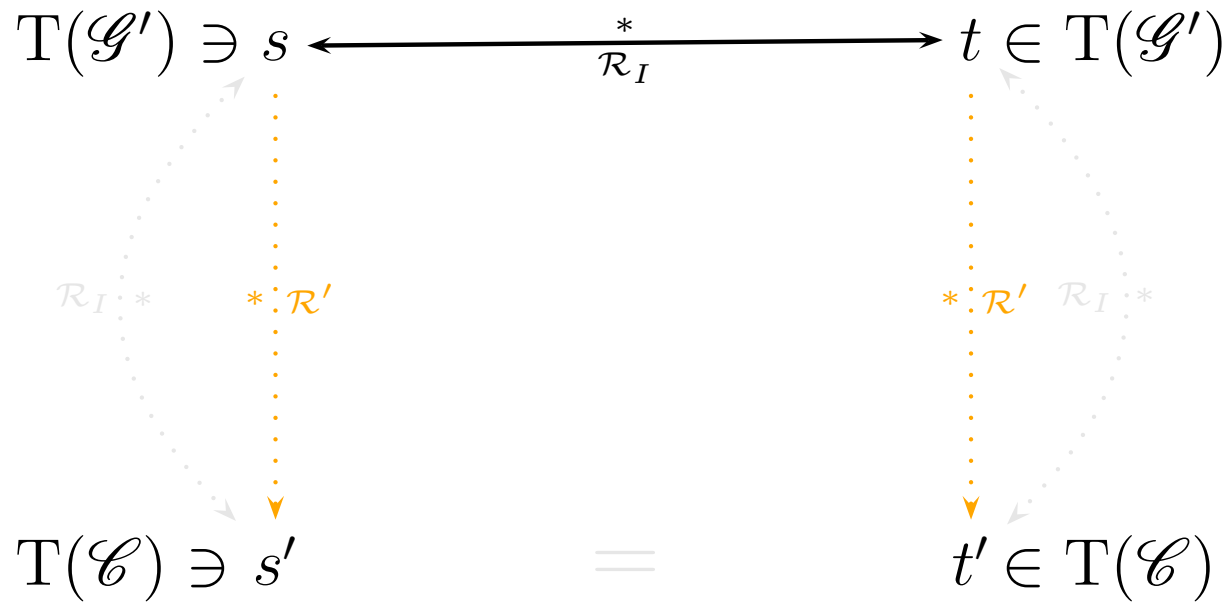
- $SC(\mathcal{R}', \mathcal{G})$ ,
- $\overset{*}{\leftrightarrow} \mathcal{R}' \subseteq \overset{*}{\leftrightarrow} \mathcal{R}_A = \overset{*}{\leftrightarrow} \mathcal{R}_I$
- $CR(\mathcal{R}_I)$ ,  
 $T(\mathcal{C}) \subseteq NF(\mathcal{R}_I)$

# $\overset{*}{\leftrightarrow} \mathcal{R}_I = \overset{*}{\leftrightarrow} \mathcal{R}'$ on $T(\mathcal{G}')$

$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

• Elimination

$$\mathcal{R}_k \xrightarrow[E]{*} \mathcal{R}_k \setminus \{l \rightarrow r\}$$



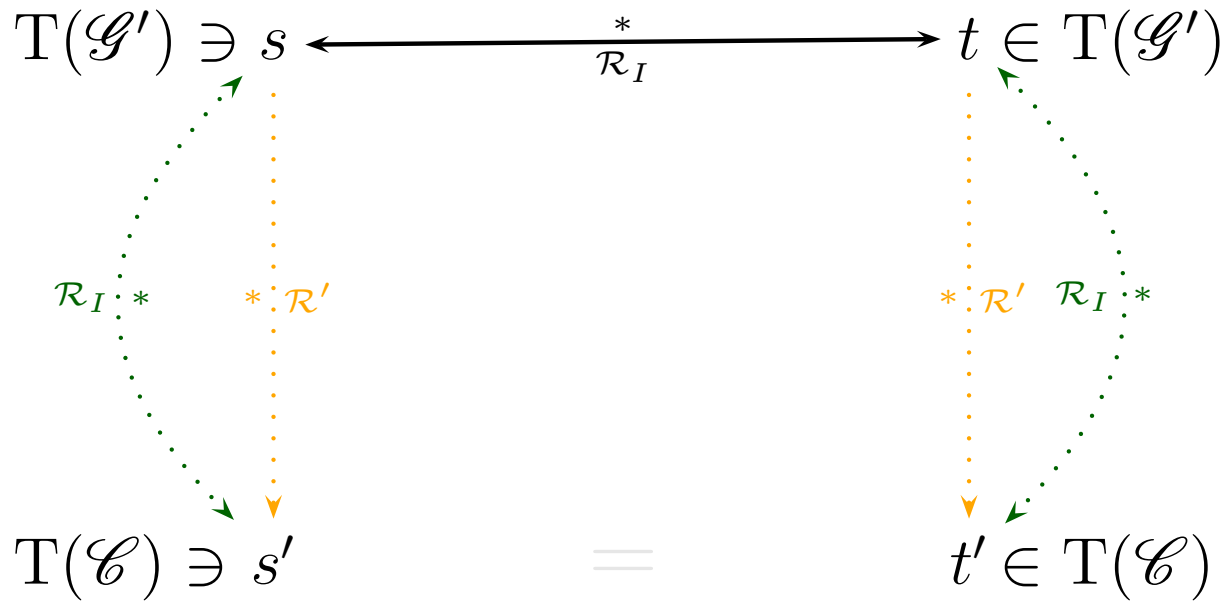
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$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

• Elimination

$$\mathcal{R}_k \xrightarrow[E]{*} \mathcal{R}_k \setminus \{l \rightarrow r\}$$



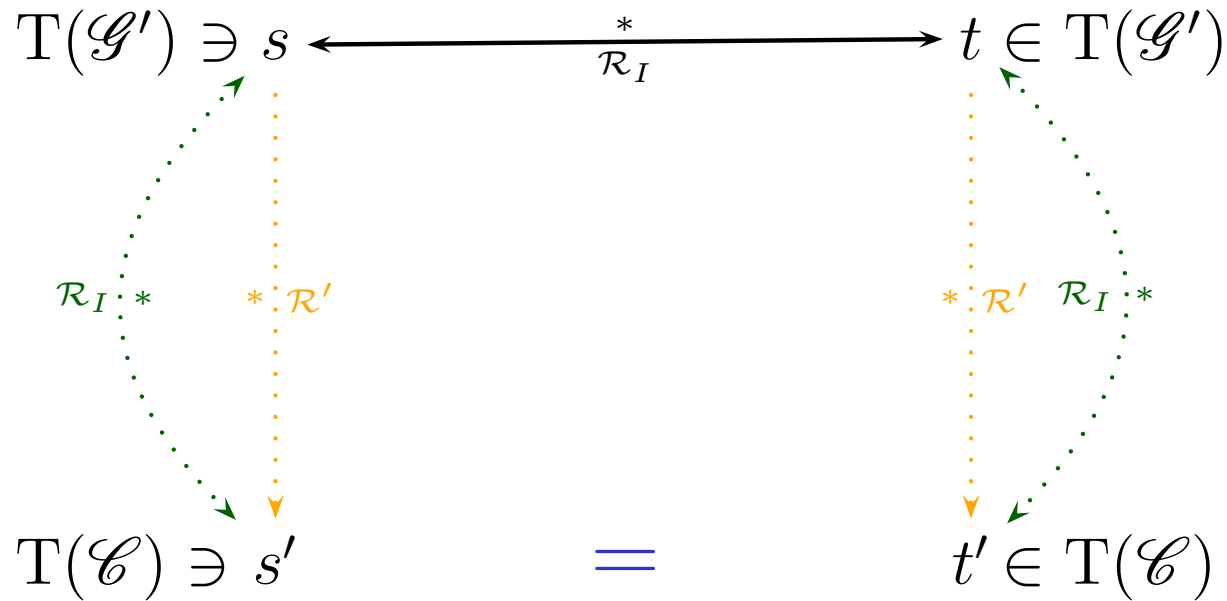
- $SC(\mathcal{R}', \mathcal{G})$ ,
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$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

• Elimination

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- $SC(\mathcal{R}', \mathcal{G})$ ,
- $\overset{*}{\leftrightarrow} \mathcal{R}' \subseteq \overset{*}{\leftrightarrow} \mathcal{R}_A = \overset{*}{\leftrightarrow} \mathcal{R}_I$
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 $T(\mathcal{C}) \subseteq NF(\mathcal{R}_I)$

# $\mathcal{R} \simeq_{\mathcal{G} \cap \mathcal{G}'} \mathcal{R}' (\subseteq)$

From 2 to 4,  $\overset{*}{\leftrightarrow}_{\mathcal{R}} = \overset{*}{\leftrightarrow}_{\mathcal{R}'}$  on  $T(\mathcal{G}' \cap \mathcal{G}')$

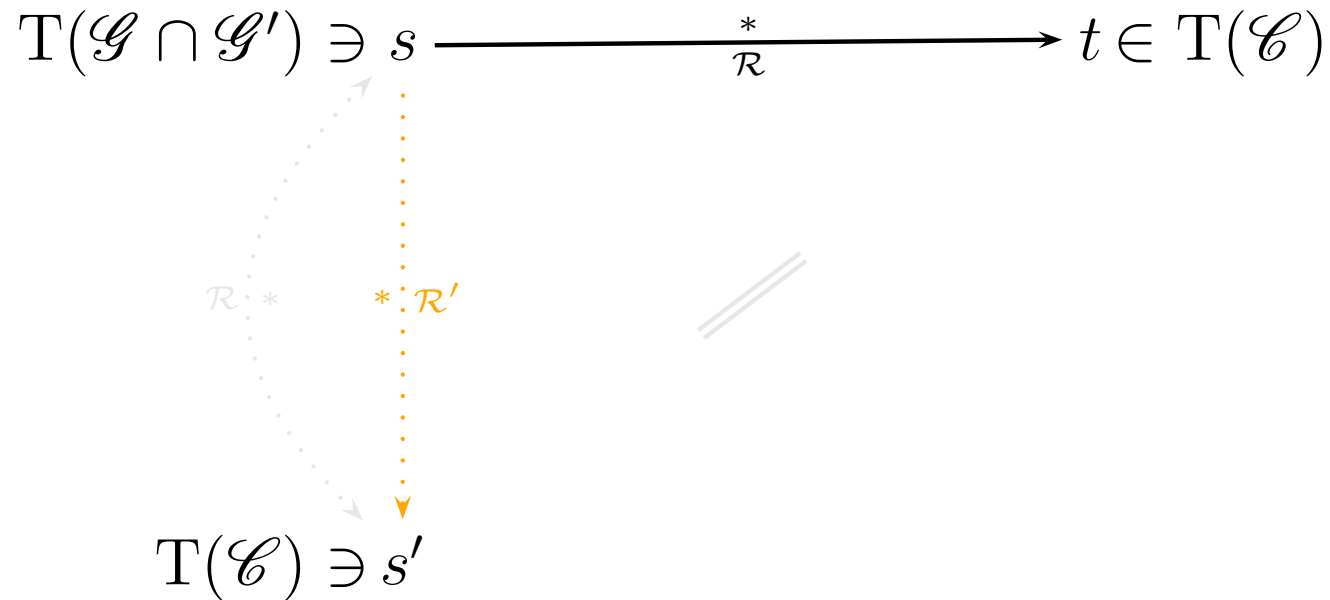
$$T(\mathcal{G} \cap \mathcal{G}') \ni s \xrightarrow[\mathcal{R}]{*} t \in T(\mathcal{C})$$



- $SC(\mathcal{R}', \mathcal{G})$ ,
- $\overset{*}{\leftarrow}_{\mathcal{R}'} \subseteq \overset{*}{\leftrightarrow}_{\mathcal{R}'} = \overset{*}{\leftrightarrow}_{\mathcal{R}}$
- $CR(\mathcal{R})$ ,  
 $T(\mathcal{C}) \subseteq NF(\mathcal{R})$

# $\mathcal{R} \simeq_{\mathcal{G} \cap \mathcal{G}'} \mathcal{R}' (\subseteq)$

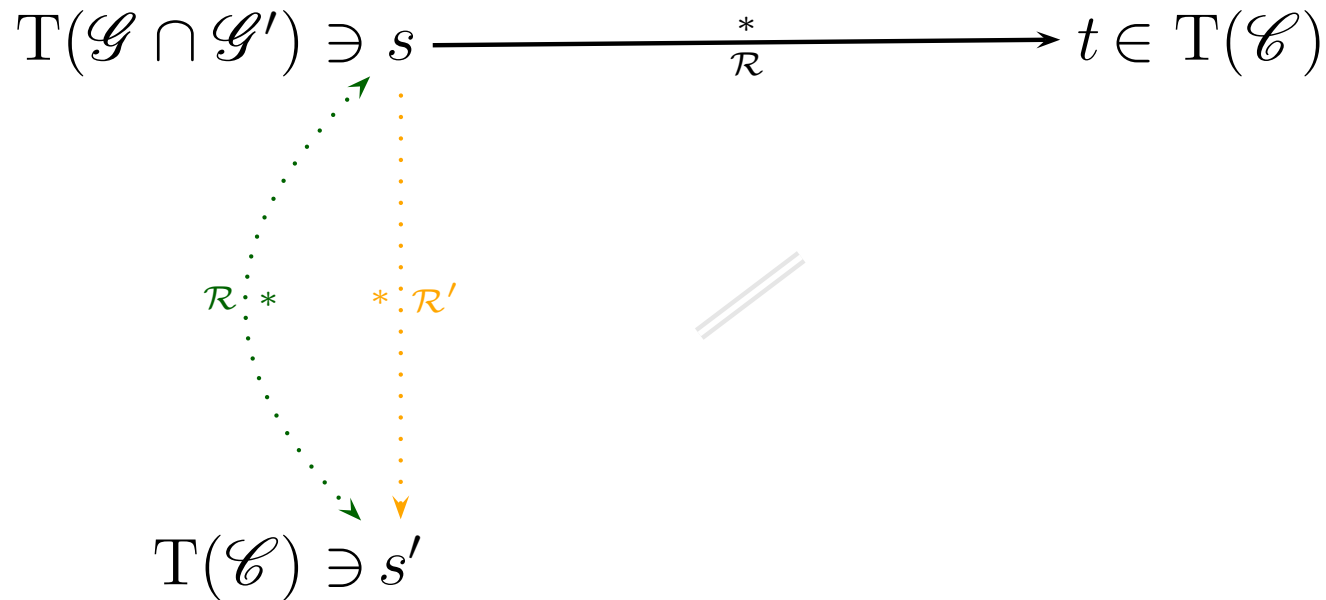
From 2 to 4,  $\overset{*}{\leftrightarrow}_{\mathcal{R}} = \overset{*}{\leftrightarrow}_{\mathcal{R}'}$  on  $T(\mathcal{G}' \cap \mathcal{G}')$



- $SC(\mathcal{R}', \mathcal{G})$ ,
- $\overset{*}{\leftarrow}_{\mathcal{R}'} \subseteq \overset{*}{\leftrightarrow}_{\mathcal{R}'} = \overset{*}{\leftrightarrow}_{\mathcal{R}}$
- $CR(\mathcal{R})$ ,  
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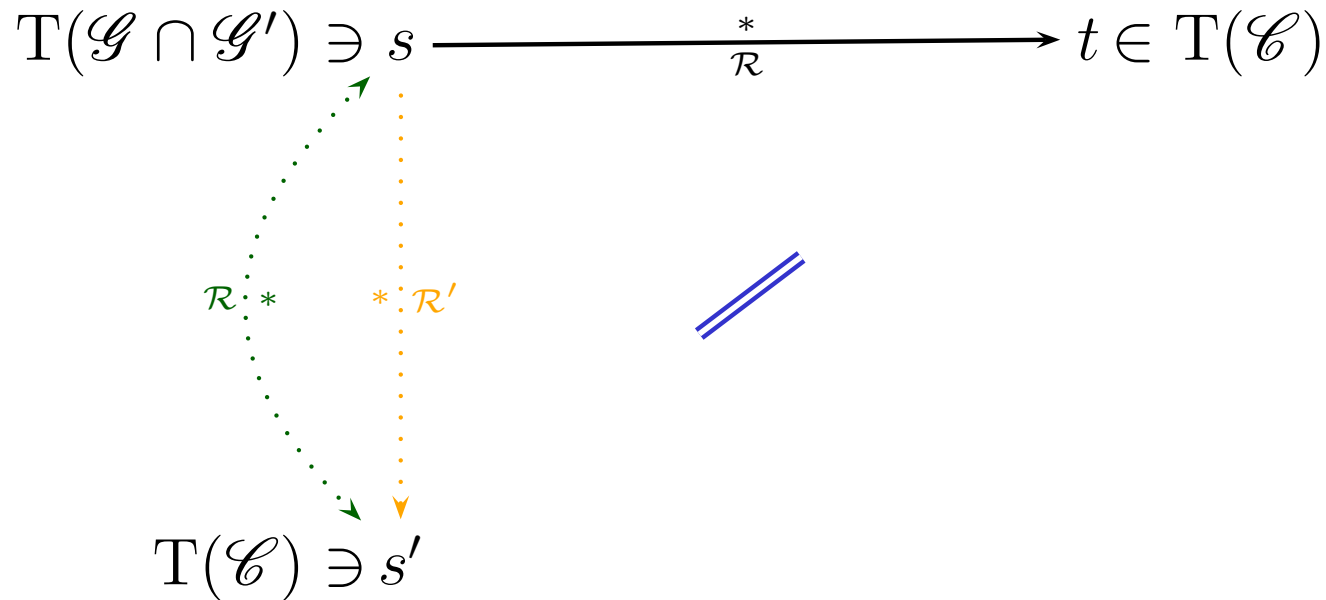


- $SC(\mathcal{R}', \mathcal{G})$ ,
- $\overset{*}{\leftarrow} \mathcal{R}' \subseteq \overset{*}{\leftrightarrow} \mathcal{R}' = \overset{*}{\leftrightarrow} \mathcal{R}$
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# $\mathcal{R} \simeq_{\mathcal{G} \cap \mathcal{G}'} \mathcal{R}' (\subseteq)$

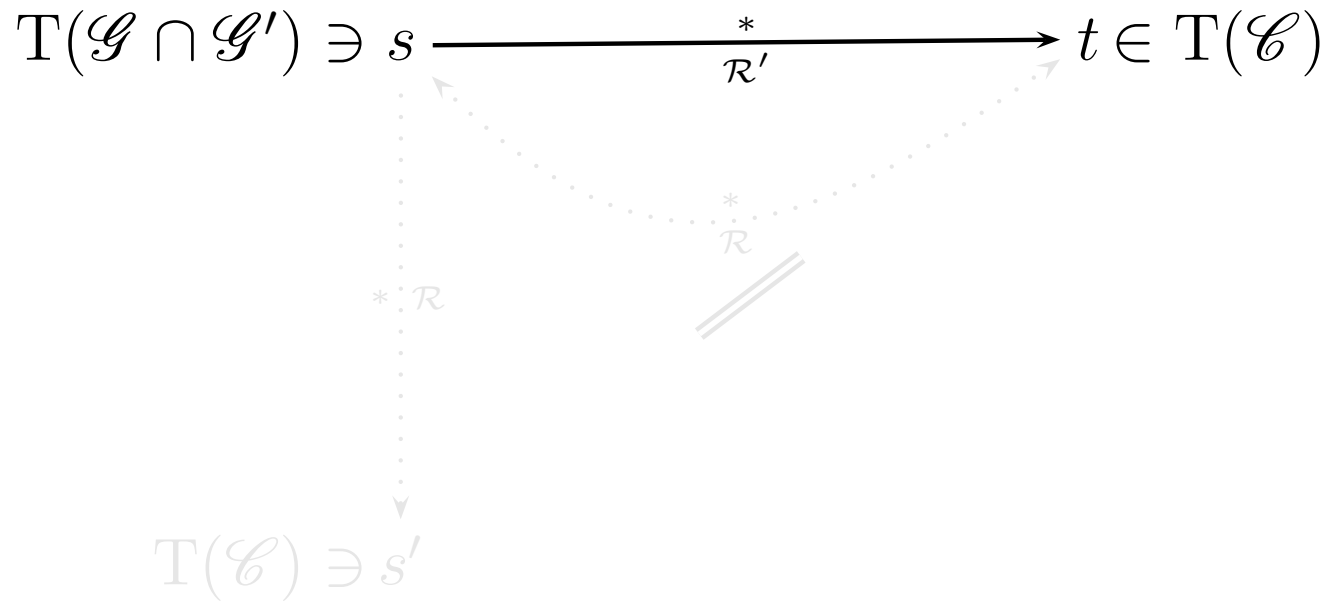
From 2 to 4,  $\overset{*}{\leftrightarrow}_{\mathcal{R}} = \overset{*}{\leftrightarrow}_{\mathcal{R}'}$  on  $T(\mathcal{G}' \cap \mathcal{G}')$



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# $\mathcal{R} \simeq_{\mathcal{G} \cap \mathcal{G}'} \mathcal{R}' (\supseteq)$

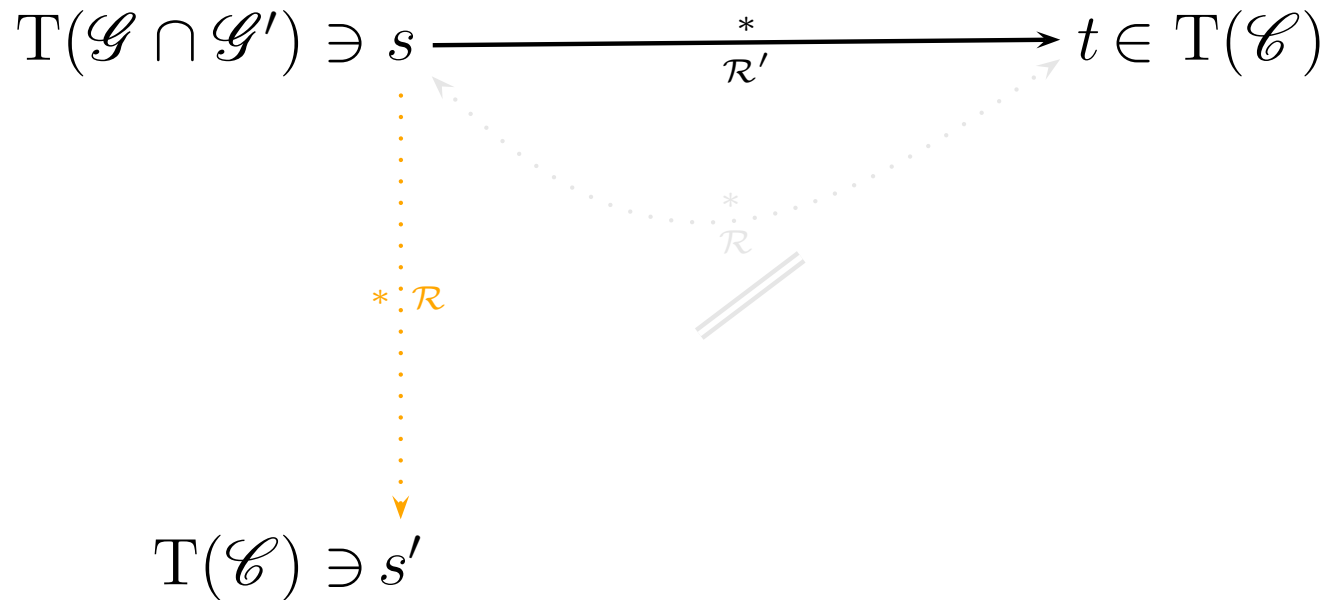
From 2 to 4,  $\overset{*}{\leftrightarrow}_{\mathcal{R}} = \overset{*}{\leftrightarrow}_{\mathcal{R}'}$  on  $T(\mathcal{G}' \cap \mathcal{G}')$



- $SC(\mathcal{R}, \mathcal{G})$ ,
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# $\mathcal{R} \simeq_{\mathcal{G} \cap \mathcal{G}'} \mathcal{R}' (\supseteq)$

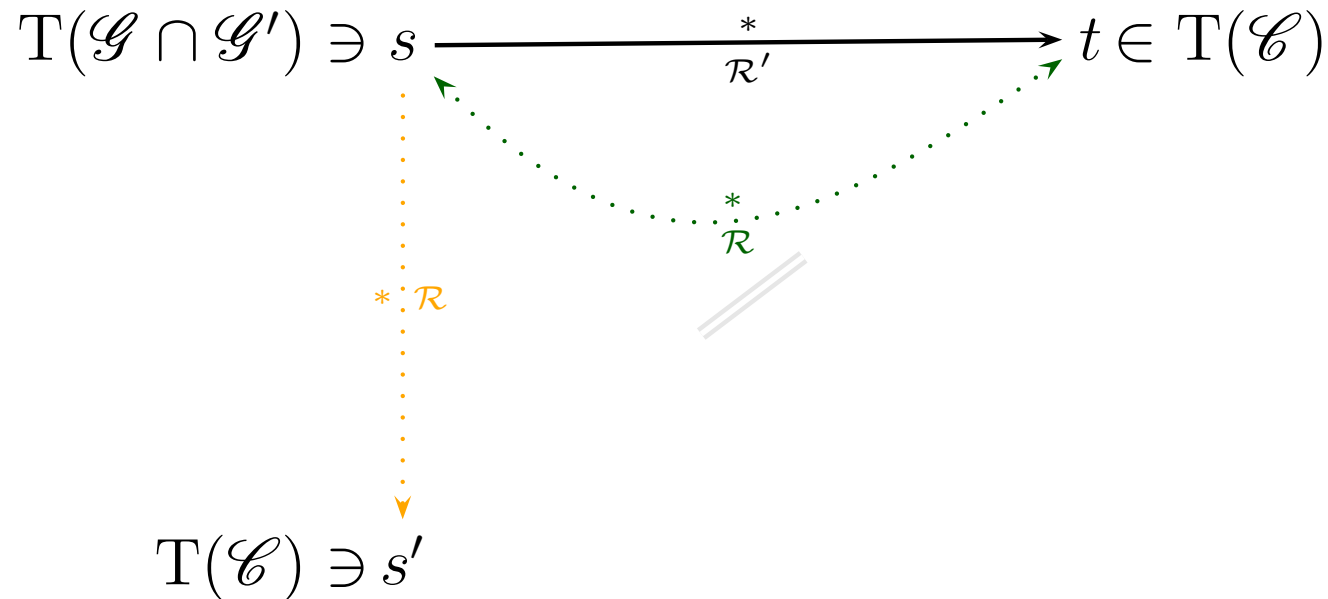
From 2 to 4,  $\overset{*}{\leftrightarrow}_{\mathcal{R}} = \overset{*}{\leftrightarrow}_{\mathcal{R}'}$  on  $T(\mathcal{G}' \cap \mathcal{G}')$



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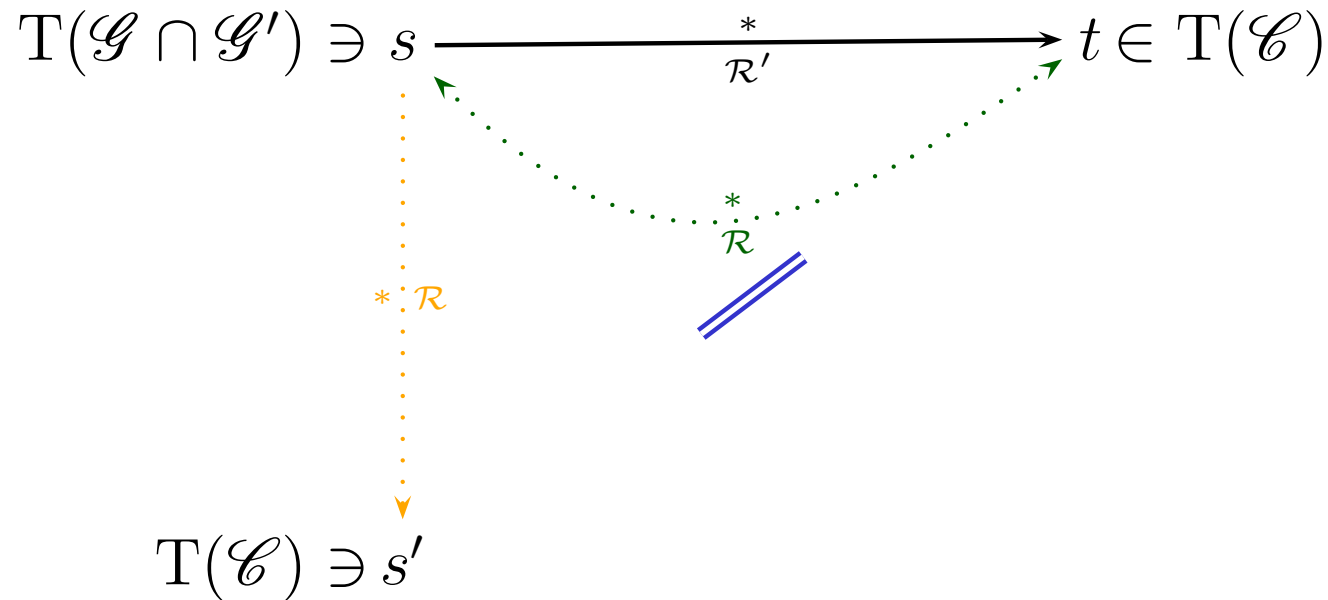
From 2 to 4,  $\overset{*}{\leftrightarrow}_{\mathcal{R}} = \overset{*}{\leftrightarrow}_{\mathcal{R}'}$  on  $T(\mathcal{G}' \cap \mathcal{G}')$



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# Conjectures

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$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

1.  $\text{CR}(\mathcal{R})$  implies  $\text{CR}(\mathcal{R}_I)$
2. Let  $\mathcal{R}$  and  $\mathcal{R}_I$  be STTRSs over  $\mathcal{G}$  and  $\mathcal{F}$ .  $\forall s \in T(\mathcal{F}). \exists t \in T(\mathcal{G})$  s.t.  
 $s \xrightarrow{\mathcal{R}_I}^* t$

# Higher-order sufficient completeness

(Aoto, Yamada and Toyama, 2004)

## Definition

An STTRS  $\mathcal{R}$  is said to be *higher-order sufficiently complete*  $\text{HSC}(R)$  if for any basic simply typed term  $s \in \mathbb{T}^b(\mathcal{F}, \mathcal{V}^h)$  there exists  $t \in \mathbb{T}^b(\mathcal{F}, \mathcal{V}^h)$  such that  $s \xrightarrow{*}_{\mathcal{R}} t$  and that either  $t$  has an expanded variable occurrence or  $t$  is a ground constructor term.

$$\begin{aligned} (\text{foldr } f \ e \ []) &\rightarrow e \\ (\text{foldr } f \ e \ (: x \ xs)) &\rightarrow (f \ x \ (\text{foldr } f \ e \ xs)) \\ (+ \ 0 \ y) &\rightarrow y \\ (+ \ s(x) \ y) &\rightarrow (s \ (+ \ x \ y)) \end{aligned}$$

$$(\text{foldr } f \ 0 \ [1, 2, 3]) \rightarrow_{\mathcal{R}} (f \ 1 \ (\text{foldr } f \ 0 \ [2, 3])),$$

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$$(\text{foldr } + \ 0 \ [1, 2, 3]) \rightarrow_{\mathcal{R}} (+ \ 1 \ (\text{foldr } + \ 0 \ [2, 3]))$$



# Conjectures

---

$$\mathcal{R} \xrightarrow[I]{*} \mathcal{R}_I \xrightarrow[A]{*} \mathcal{R}_A \xrightarrow[E]{*} \mathcal{R}'$$

Let  $\mathcal{R}$  and  $\mathcal{R}_I$  be STTRSs over  $\mathcal{G}$  and  $\mathcal{F}$ .  $\forall s \in T(\mathcal{F}). \exists t \in T(\mathcal{G})$  s.t.  $s \xrightarrow{\mathcal{R}_I}^* t$

# Proof (sketch)

HSC( $\mathcal{R}$ ) and WN( $\mathcal{R}$ )

$$\begin{array}{ccccc}
 \tilde{s}' \in T(\mathcal{G}, \{x\}) & \xrightarrow[\mathcal{R}]{*} & t \in T(\mathcal{C}, \{x\}) & & \\
 \downarrow \sigma (= \{x \mapsto f\}) & & \downarrow \sigma & & \\
 s \in T(\mathcal{F}) & \xrightarrow[\mathcal{R}_I \setminus \mathcal{R}]{*} & s' (= \tilde{s}'\sigma) & \xrightarrow[\mathcal{R}]{*} & t\sigma & \xrightarrow[\mathcal{R}_I \setminus \mathcal{R}]{*} & \hat{t} \in T(\mathcal{G})
 \end{array}$$

# Conclusion

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- Equivalent transformation for STTRSs
- Verification of the equality of STTRSs

# Conclusion

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- Transformation templates
- Examples