Saigawa: A Confluence Tool

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Saigawa: A River in Kanazawa

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Theorem (Klein & Hirokawa '12)

Suppose $\mathcal{R}$ *and* $\mathcal{S}$ *strongly non-overlap each other and* $\mathcal{S}$ *is confluent.* $\mathcal{R} \cup \mathcal{S}$ *is confluent if and only if* $\mathsf{CP}_{\mathcal{S}}(\mathcal{R}) \subseteq \downarrow_{\mathcal{R} \cup \mathcal{S}}$.

demo

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$\Rightarrow$ capture large classes of TRSs with least effort