

Pólya's Proof of the Weighted Arithmetic–Geometric Mean Inequality

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Abstract

This article provides a formalisation of the Weighted Arithmetic–Geometric Mean Inequality: given non-negative reals a_1, \dots, a_n and non-negative weights w_1, \dots, w_n such that $w_1 + \dots + w_n = 1$, we have

$$\prod_{i=1}^n a_i^{w_i} \leq \sum_{i=1}^n w_i a_i .$$

If the weights are additionally all non-zero, equality holds if and only if $a_1 = \dots = a_n$.

As a corollary with $w_1 = \dots = w_n = \frac{1}{n}$, the regular arithmetic–geometric mean inequality follows, namely that

$$\sqrt[n]{a_1 \dots a_n} \leq \frac{1}{n}(a_1 + \dots + a_n) .$$

I follow Pólya's elegant proof, which uses the inequality $1 + x \leq e^x$ as a starting point. Pólya claims that this proof came to him in a dream, and that it was ‘the best mathematics he had ever dreamt.’ [1, pp. 22–26]

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1 The Weighted Arithmetic–Geometric Mean Inequality

```
theory Weighted-Arithmetic-Geometric-Mean
  imports Complex-Main
begin
```

1.1 Auxiliary Facts

```
lemma root-powr-inverse':  $0 < n \implies 0 \leq x \implies \text{root } n \ x = x \text{ powr } (1/n)$ 
  by (cases  $x = 0$ ) (auto simp: root-powr-inverse)
```

```
lemma powr-sum-distrib-real-right:
  assumes  $a \neq 0$ 
  shows  $(\prod_{x \in X}. a \text{ powr } e \ x :: \text{real}) = a \text{ powr } (\sum_{x \in X}. e \ x)$ 
  using assms
  by (induction  $X$  rule: infinite-finite-induct) (auto simp: powr-add)
```

```
lemma powr-sum-distrib-real-left:
  assumes  $\bigwedge x. x \in X \implies a \ x \geq 0$ 
  shows  $(\prod_{x \in X}. a \ x \text{ powr } e :: \text{real}) = (\prod_{x \in X}. a \ x) \text{ powr } e$ 
  using assms
  by (induction  $X$  rule: infinite-finite-induct)
    (auto simp: powr-mult prod-nonneg)
```

```
lemma (in linordered-semidom) prod-mono-strict':
  assumes  $i \in A$ 
  assumes finite  $A$ 
  assumes  $\bigwedge i. i \in A \implies 0 \leq f \ i \wedge f \ i \leq g \ i$ 
  assumes  $\bigwedge i. i \in A \implies 0 < g \ i$ 
  assumes  $f \ i < g \ i$ 
  shows  $\text{prod } f \ A < \text{prod } g \ A$ 
proof -
  have  $\text{prod } f \ A = f \ i * \text{prod } f \ (A - \{i\})$ 
    using assms by (intro prod.remove)
  also have  $\dots \leq f \ i * \text{prod } g \ (A - \{i\})$ 
    using assms by (intro mult-left-mono prod-mono) auto
  also have  $\dots < g \ i * \text{prod } g \ (A - \{i\})$ 
    using assms by (intro mult-strict-right-mono prod-pos) auto
  also have  $\dots = \text{prod } g \ A$ 
    using assms by (intro prod.remove [symmetric])
  finally show ?thesis .
qed
```

```
lemma prod-ge-pointwise-le-imp-pointwise-eq:
  fixes  $f :: 'a \Rightarrow \text{real}$ 
  assumes finite  $X$ 
  assumes  $\text{ge}: \text{prod } f \ X \geq \text{prod } g \ X$ 
  assumes  $\text{nonneg}: \bigwedge x. x \in X \implies f \ x \geq 0$ 
```

```

assumes pos:  $\bigwedge x. x \in X \implies g\ x > 0$ 
assumes le:  $\bigwedge x. x \in X \implies f\ x \leq g\ x$  and x:  $x \in X$ 
shows  $f\ x = g\ x$ 
proof (rule ccontr)
  assume  $f\ x \neq g\ x$ 
  with le[of x] and x have  $f\ x < g\ x$ 
    by auto
  hence  $\text{prod } f\ X < \text{prod } g\ X$ 
    using x and le and nonneg and pos and  $\langle \text{finite } X \rangle$ 
    by (intro prod-mono-strict') auto
  with ge show False
    by simp
qed

```

```

lemma powr-right-real-eq-iff:
  assumes  $a \geq (0 :: \text{real})$ 
  shows  $a \text{ powr } x = a \text{ powr } y \iff a = 0 \vee a = 1 \vee x = y$ 
  using assms by (auto simp: powr-def)

```

```

lemma powr-left-real-eq-iff:
  assumes  $a \geq (0 :: \text{real})$   $b \geq 0$   $x \neq 0$ 
  shows  $a \text{ powr } x = b \text{ powr } x \iff a = b$ 
  using assms by (auto simp: powr-def)

```

```

lemma exp-real-eq-one-plus-iff:
  fixes  $x :: \text{real}$ 
  shows  $\exp x = 1 + x \iff x = 0$ 
proof (cases  $x = 0$ )
  case False
    define  $f :: \text{real} \Rightarrow \text{real}$  where  $f = (\lambda x. \exp x - 1 - x)$ 
    have deriv: (f has-field-derivative ( $\exp x - 1$ )) (at x) for x
      by (auto simp: f-def intro!: derivative-eq-intros)

    have  $\exists z. z > \min x\ 0 \wedge z < \max x\ 0 \wedge f(\max x\ 0) - f(\min x\ 0) =$ 
       $(\max x\ 0 - \min x\ 0) * (\exp z - 1)$ 
      using MVT2[of  $\min x\ 0$   $\max x\ 0$  f  $\lambda x. \exp x - 1$ ] deriv False
      by (auto simp: min-def max-def)
    then obtain z where  $z \in \{\min x\ 0 <..<\max x\ 0\}$ 
       $f(\max x\ 0) - f(\min x\ 0) = (\max x\ 0 - \min x\ 0) * (\exp z - 1)$ 
      by (auto simp: f-def)
    thus ?thesis using False
      by (cases  $x\ 0 :: \text{real}$  rule: linorder-cases) (auto simp: f-def)
qed auto

```

1.2 The Inequality

We first prove the equality under the assumption that all the a_i and w_i are positive.

lemma weighted-arithmetic-geometric-mean-pos:

```

fixes  $a\ w :: 'a \Rightarrow \text{real}$ 
assumes  $\text{finite } X$ 
assumes  $\text{pos1}: \bigwedge x. x \in X \implies a\ x > 0$ 
assumes  $\text{pos2}: \bigwedge x. x \in X \implies w\ x > 0$ 
assumes  $\text{sum-weights}: (\sum x \in X. w\ x) = 1$ 
shows  $(\prod x \in X. a\ x^{\text{powr } w\ x}) \leq (\sum x \in X. w\ x * a\ x)$ 
proof -
  note  $\text{nonneg1} = \text{less-imp-le}[OF\ \text{pos1}]$ 
  note  $\text{nonneg2} = \text{less-imp-le}[OF\ \text{pos2}]$ 
  define  $A$  where  $A = (\sum x \in X. w\ x * a\ x)$ 
  define  $r$  where  $r = (\lambda x. a\ x / A - 1)$ 
  from  $\text{sum-weights}$  have  $X \neq \{\}$  by  $\text{auto}$ 
  hence  $A \neq 0$ 
    unfolding  $A\text{-def}$  using  $\text{nonneg1}\ \text{nonneg2}\ \text{pos1}\ \text{pos2}\ \langle \text{finite } X \rangle$ 
    by  $(\text{subst}\ \text{sum-nonneg-eq-0-iff})\ \text{force+}$ 
  moreover from  $\text{nonneg1}\ \text{nonneg2}$  have  $A \geq 0$ 
    by  $(\text{auto}\ \text{simp: } A\text{-def}\ \text{intro!}: \text{sum-nonneg})$ 
  ultimately have  $A > 0$  by  $\text{simp}$ 

  have  $(\prod x \in X. (1 + r\ x)^{\text{powr } w\ x}) = (\prod x \in X. (a\ x / A)^{\text{powr } w\ x})$ 
    by  $(\text{simp}\ \text{add: } r\text{-def})$ 
  also have  $\dots = (\prod x \in X. a\ x^{\text{powr } w\ x}) / (\prod x \in X. A^{\text{powr } w\ x})$ 
    unfolding  $\text{prod-dividef}$   $[\text{symmetric}]$ 
    using  $\langle \text{assms}\ \text{pos2}\ \langle A > 0 \rangle \rangle$  by  $(\text{intro}\ \text{prod.cong}\ \text{powr-divide})\ (\text{auto}\ \text{intro: } \text{less-imp-le})$ 
  also have  $(\prod x \in X. A^{\text{powr } w\ x}) = \exp((\sum x \in X. w\ x) * \ln A)$ 
    using  $\langle A > 0 \rangle$  and  $\langle \text{finite } X \rangle$  by  $(\text{simp}\ \text{add: } \text{powr-def}\ \text{exp-sum}\ \text{sum-distrib-right})$ 
  also have  $(\sum x \in X. w\ x) = 1$  by  $\text{fact}$ 
  also have  $\exp(1 * \ln A) = A$ 
    using  $\langle A > 0 \rangle$  by  $\text{simp}$ 
  finally have  $\text{lhs: } (\prod x \in X. (1 + r\ x)^{\text{powr } w\ x}) = (\prod x \in X. a\ x^{\text{powr } w\ x}) / A$ 

  have  $(\prod x \in X. \exp(w\ x * r\ x)) = \exp(\sum x \in X. w\ x * r\ x)$ 
    using  $\langle \text{finite } X \rangle$  by  $(\text{simp}\ \text{add: } \text{exp-sum})$ 
  also have  $(\sum x \in X. w\ x * r\ x) = (\sum x \in X. a\ x * w\ x) / A - 1$ 
    using  $\langle A > 0 \rangle$  by  $(\text{simp}\ \text{add: } r\text{-def}\ \text{algebra-simps}\ \text{sum-subtractf}\ \text{sum-divide-distrib}\ \text{sum-weights})$ 
  also have  $(\sum x \in X. a\ x * w\ x) / A = 1$ 
    using  $\langle A > 0 \rangle$  by  $(\text{simp}\ \text{add: } A\text{-def}\ \text{mult.commute})$ 
  finally have  $\text{rhs: } (\prod x \in X. \exp(w\ x * r\ x)) = 1$  by  $\text{simp}$ 

  have  $(\prod x \in X. a\ x^{\text{powr } w\ x}) / A = (\prod x \in X. (1 + r\ x)^{\text{powr } w\ x})$ 
    by  $(\text{fact}\ \text{lhs}\ [\text{symmetric}])$ 
  also have  $(\prod x \in X. (1 + r\ x)^{\text{powr } w\ x}) \leq (\prod x \in X. \exp(w\ x * r\ x))$ 
proof  $(\text{intro}\ \text{prod-mono}\ \text{conjI})$ 
  fix  $x$  assume  $x: x \in X$ 
  have  $1 + r\ x \leq \exp(r\ x)$ 
    by  $(\text{rule}\ \text{exp-ge-add-one-self})$ 
  hence  $(1 + r\ x)^{\text{powr } w\ x} \leq \exp(r\ x)^{\text{powr } w\ x}$ 

```

```

    using nonneg1[of x] nonneg2[of x] x <A > 0>
    by (intro powr-mono2) (auto simp: r-def field-simps)
  also have ... = exp (w x * r x)
    by (simp add: powr-def)
  finally show (1 + r x) powr w x ≤ exp (w x * r x) .
qed auto
also have (∏ x∈X. exp (w x * r x)) = 1 by (fact rhs)
finally show (∏ x∈X. a x powr w x) ≤ A
  using <A > 0> by (simp add: field-simps)
qed

```

We can now relax the positivity assumptions to non-negativity: if one of the a_i is zero, the theorem becomes trivial (note that $0^0 = 0$ by convention for the real-valued power operator (*powr*)).

Otherwise, we can simply remove all the indices that have weight 0 and apply the above auxiliary version of the theorem.

theorem *weighted-arithmetic-geometric-mean:*

```

  fixes a w :: 'a ⇒ real
  assumes finite X
  assumes nonneg1: ⋀x. x ∈ X ⇒ a x ≥ 0
  assumes nonneg2: ⋀x. x ∈ X ⇒ w x ≥ 0
  assumes sum-weights: (∑ x∈X. w x) = 1
  shows (∏ x∈X. a x powr w x) ≤ (∑ x∈X. w x * a x)
proof (cases ∃ x∈X. a x = 0)
  case True
  hence (∏ x∈X. a x powr w x) = 0
    using <finite X> by simp
  also have ... ≤ (∑ x∈X. w x * a x)
    by (intro sum-nonneg mult-nonneg-nonneg assms)
  finally show ?thesis .

```

next

```

  case False
  have (∑ x∈X-⋅{x. w x = 0}. w x) = (∑ x∈X. w x)
    by (intro sum.mono-neutral-left assms) auto
  also have ... = 1 by fact
  finally have sum-weights': (∑ x∈X-⋅{x. w x = 0}. w x) = 1 .

```

```

  have (∏ x∈X. a x powr w x) = (∏ x∈X-⋅{x. w x = 0}. a x powr w x)
    using <finite X> False by (intro prod.mono-neutral-right) auto
  also have ... ≤ (∑ x∈X-⋅{x. w x = 0}. w x * a x) using assms False
    by (intro weighted-arithmetic-geometric-mean-pos sum-weights')
    (auto simp: order.strict-iff-order)
  also have ... = (∑ x∈X. w x * a x)
    using <finite X> by (intro sum.mono-neutral-left) auto
  finally show ?thesis .

```

qed

We can derive the regular arithmetic/geometric mean inequality from this by simply setting all the weights to $\frac{1}{n}$:

corollary *arithmetic-geometric-mean:*

```

fixes  $a :: 'a \Rightarrow \text{real}$ 
assumes  $\text{finite } X$ 
defines  $n \equiv \text{card } X$ 
assumes  $\text{nonneg}: \bigwedge x. x \in X \implies a\ x \geq 0$ 
shows  $\text{root } n\ (\prod_{x \in X}. a\ x) \leq (\sum_{x \in X}. a\ x) / n$ 
proof ( $\text{cases } X = \{\}$ )
  case  $\text{False}$ 
  with  $\text{assms have } n: n > 0$ 
    by  $\text{auto}$ 
  have  $(\prod_{x \in X}. a\ x\ \text{powr } (1 / n)) \leq (\sum_{x \in X}. (1 / n) * a\ x)$ 
    using  $n\ \text{assms by (intro weighted-arithmetic-geometric-mean) auto}$ 
  also have  $(\prod_{x \in X}. a\ x\ \text{powr } (1 / n)) = (\prod_{x \in X}. a\ x)\ \text{powr } (1 / n)$ 
    using  $\text{nonneg by (subst powr-sum-distrib-real-left) auto}$ 
  also have  $\dots = \text{root } n\ (\prod_{x \in X}. a\ x)$ 
    using  $\langle n > 0 \rangle\ \text{nonneg by (subst root-powr-inverse') (auto simp: prod-nonneg)}$ 
  also have  $(\sum_{x \in X}. (1 / n) * a\ x) = (\sum_{x \in X}. a\ x) / n$ 
    by (subst sum-distrib-left [symmetric]) auto
  finally show  $?thesis$  .
qed ( $\text{auto simp: n-def}$ )

```

1.3 The Equality Case

Next, we show that weighted arithmetic and geometric mean are equal if and only if all the a_i are equal.

We first prove the more difficult direction as a lemmas and again first assume positivity of all a_i and w_i and will relax this somewhat later.

lemma *weighted-arithmetic-geometric-mean-eq-iff-pos:*

```

fixes  $a\ w :: 'a \Rightarrow \text{real}$ 
assumes  $\text{finite } X$ 
assumes  $\text{pos1}: \bigwedge x. x \in X \implies a\ x > 0$ 
assumes  $\text{pos2}: \bigwedge x. x \in X \implies w\ x > 0$ 
assumes  $\text{sum-weights}: (\sum_{x \in X}. w\ x) = 1$ 
assumes  $\text{eq}: (\prod_{x \in X}. a\ x\ \text{powr } w\ x) = (\sum_{x \in X}. w\ x * a\ x)$ 
shows  $\forall x \in X. \forall y \in X. a\ x = a\ y$ 
proof –
  note  $\text{nonneg1} = \text{less-imp-le}[OF\ \text{pos1}]$ 
  note  $\text{nonneg2} = \text{less-imp-le}[OF\ \text{pos2}]$ 
  define  $A$  where  $A = (\sum_{x \in X}. w\ x * a\ x)$ 
  define  $r$  where  $r = (\lambda x. a\ x / A - 1)$ 
  from  $\text{sum-weights have } X \neq \{\}$  by  $\text{auto}$ 
  hence  $A \neq 0$ 
    unfolding  $A\text{-def using nonneg1 nonneg2 pos1 pos2 } \langle \text{finite } X \rangle$ 
    by (subst sum-nonneg-eq-0-iff) force+
  moreover from nonneg1 nonneg2 have  $A \geq 0$ 
    by (auto simp: A-def intro!: sum-nonneg)
  ultimately have  $A > 0$  by  $\text{simp}$ 

```

have $r\text{-ge}: r\ x \geq -1$ **if** $x: x \in X$ **for** x
using $\langle A > 0 \rangle$ $\text{pos1}[OF\ x]$ **by** (*auto simp: r-def field-simps*)

have $(\prod_{x \in X}. (1 + r\ x) \text{ powr } w\ x) = (\prod_{x \in X}. (a\ x / A) \text{ powr } w\ x)$
by (*simp add: r-def*)
also have $\dots = (\prod_{x \in X}. a\ x \text{ powr } w\ x) / (\prod_{x \in X}. A \text{ powr } w\ x)$
unfolding *prod-dividef [symmetric]*
using *assms pos2* $\langle A > 0 \rangle$ **by** (*intro prod.cong powr-divide*) (*auto intro: less-imp-le*)
also have $(\prod_{x \in X}. A \text{ powr } w\ x) = \exp((\sum_{x \in X}. w\ x) * \ln A)$
using $\langle A > 0 \rangle$ **and** $\langle \text{finite } X \rangle$ **by** (*simp add: powr-def exp-sum sum-distrib-right*)
also have $(\sum_{x \in X}. w\ x) = 1$ **by** *fact*
also have $\exp(1 * \ln A) = A$
using $\langle A > 0 \rangle$ **by** *simp*
finally have *lhs*: $(\prod_{x \in X}. (1 + r\ x) \text{ powr } w\ x) = (\prod_{x \in X}. a\ x \text{ powr } w\ x) / A$.

have $(\prod_{x \in X}. \exp(w\ x * r\ x)) = \exp(\sum_{x \in X}. w\ x * r\ x)$
using $\langle \text{finite } X \rangle$ **by** (*simp add: exp-sum*)
also have $(\sum_{x \in X}. w\ x * r\ x) = (\sum_{x \in X}. a\ x * w\ x) / A - 1$
using $\langle A > 0 \rangle$ **by** (*simp add: r-def algebra-simps sum-subtractf sum-divide-distrib sum-weights*)
also have $(\sum_{x \in X}. a\ x * w\ x) / A = 1$
using $\langle A > 0 \rangle$ **by** (*simp add: A-def mult.commute*)
finally have *rhs*: $(\prod_{x \in X}. \exp(w\ x * r\ x)) = 1$ **by** *simp*

have $a\ x = A$ **if** $x: x \in X$ **for** x
proof –
have $(1 + r\ x) \text{ powr } w\ x = \exp(w\ x * r\ x)$
proof (*rule prod-ge-pointwise-le-imp-pointwise-eq*
 $[\text{where } f = \lambda x. (1 + r\ x) \text{ powr } w\ x \text{ and } g = \lambda x. \exp(w\ x * r\ x)]$)
show $(1 + r\ x) \text{ powr } w\ x \leq \exp(w\ x * r\ x)$ **if** $x: x \in X$ **for** x
proof –
have $1 + r\ x \leq \exp(r\ x)$
by (*rule exp-ge-add-one-self*)
hence $(1 + r\ x) \text{ powr } w\ x \leq \exp(r\ x) \text{ powr } w\ x$
using *nonneg1[of x] nonneg2[of x] x* $\langle A > 0 \rangle$
by (*intro powr-mono2*) (*auto simp: r-def field-simps*)
also have $\dots = \exp(w\ x * r\ x)$
by (*simp add: powr-def*)
finally show $(1 + r\ x) \text{ powr } w\ x \leq \exp(w\ x * r\ x)$.
qed

next
show $(\prod_{x \in X}. (1 + r\ x) \text{ powr } w\ x) \geq (\prod_{x \in X}. \exp(w\ x * r\ x))$
proof –
have $(\prod_{x \in X}. (1 + r\ x) \text{ powr } w\ x) = (\prod_{x \in X}. a\ x \text{ powr } w\ x) / A$
by (*fact lhs*)
also have $\dots = 1$
using $\langle A \neq 0 \rangle$ **by** (*simp add: eq A-def*)
also have $\dots = (\prod_{x \in X}. \exp(w\ x * r\ x))$

```

    by (simp add: rhs)
  finally show ?thesis by simp
qed
qed (use x ⟨finite X⟩ in auto)

also have  $\exp (w x * r x) = \exp (r x) \text{ powr } w x$ 
  by (simp add: powr-def)
finally have  $1 + r x = \exp (r x)$ 
  using x pos2[of x] r-ge[of x] by (subst (asm) powr-left-real-eq-iff) auto
hence  $r x = 0$ 
  using exp-real-eq-one-plus-iff[of r x] by auto
hence  $a x = A$ 
  using ⟨A > 0⟩ by (simp add: r-def field-simps)
thus ?thesis
  by (simp add: )
qed
thus  $\forall x \in X. \forall y \in X. a x = a y$ 
  by auto
qed

```

We can now show the full theorem and relax the positivity condition on the a_i to non-negativity. This is possible because if some a_i is zero and the two means coincide, then the product is obviously 0, but the sum can only be 0 if *all* the a_i are 0.

```

theorem weighted-arithmetic-geometric-mean-eq-iff:
  fixes a w :: 'a  $\Rightarrow$  real
  assumes finite X
  assumes nonneg1:  $\bigwedge x. x \in X \implies a x \geq 0$ 
  assumes pos2:  $\bigwedge x. x \in X \implies w x > 0$ 
  assumes sum-weights:  $(\sum x \in X. w x) = 1$ 
  shows  $(\prod x \in X. a x \text{ powr } w x) = (\sum x \in X. w x * a x) \longleftrightarrow X \neq \{\} \wedge (\forall x \in X. \forall y \in X. a x = a y)$ 
proof
  assume *:  $X \neq \{\} \wedge (\forall x \in X. \forall y \in X. a x = a y)$ 
  from * have  $X \neq \{\}$ 
    by blast

  from * obtain c where  $c: \bigwedge x. x \in X \implies a x = c \wedge c \geq 0$ 
proof (cases  $X = \{\}$ )
  case False
    then obtain x where  $x \in X$  by blast
    thus ?thesis using * that[of a x] nonneg1[of x] by metis
  next
  case True
    thus ?thesis
      using that[of 1] by auto
qed

```

```

have  $(\prod x \in X. a x \text{ powr } w x) = (\prod x \in X. c \text{ powr } w x)$ 

```



```

    by (simp add: c)
  also have ... = c
    using assms c <math>X \neq \{\}> by (cases c = 0) (auto simp: powr-sum-distrib-real-right)
  also have ... =  $(\sum_{x \in X}. w\ x * a\ x)$ 
    using sum-weights by (simp add: c(1) flip: sum-distrib-left sum-distrib-right)
  finally show  $(\prod_{x \in X}. a\ x\ powr\ w\ x) = (\sum_{x \in X}. w\ x * a\ x)$  .
next
  assume *:  $(\prod_{x \in X}. a\ x\ powr\ w\ x) = (\sum_{x \in X}. w\ x * a\ x)$ 
  have  $X \neq \{\}$ 
    using * by auto
  moreover have  $(\forall x \in X. \forall y \in X. a\ x = a\ y)$ 
  proof (cases  $\exists x \in X. a\ x = 0$ )
    case False
    with nonneg1 have pos1:  $\forall x \in X. a\ x > 0$ 
      by force
    thus ?thesis
      using weighted-arithmetic-geometric-mean-eq-iff-pos[of X a w] assms *
      by blast
  next
    case True
    hence  $(\prod_{x \in X}. a\ x\ powr\ w\ x) = 0$ 
      using assms by auto
    with * have  $(\sum_{x \in X}. w\ x * a\ x) = 0$ 
      by auto
    also have ?this  $\longleftrightarrow (\forall x \in X. w\ x * a\ x = 0)$ 
      using assms by (intro sum-nonneg-eq-0-iff mult-nonneg-nonneg) (auto intro:
less-imp-le)
    finally have  $(\forall x \in X. a\ x = 0)$ 
      using pos2 by force
    thus ?thesis
      by auto
  qed
  ultimately show  $X \neq \{\} \wedge (\forall x \in X. \forall y \in X. a\ x = a\ y)$ 
    by blast
qed

```

Again, we derive a version for the unweighted arithmetic/geometric mean.

corollary *arithmetic-geometric-mean-eq-iff*:

```

  fixes a :: 'a  $\Rightarrow$  real
  assumes finite X
  defines n  $\equiv$  card X
  assumes nonneg:  $\bigwedge x. x \in X \implies a\ x \geq 0$ 
  shows root n  $(\prod_{x \in X}. a\ x) = (\sum_{x \in X}. a\ x) / n \longleftrightarrow (\forall x \in X. \forall y \in X. a\ x = a\ y)$ 
  proof (cases X =  $\{\}$ )
    case False
    with assms have n > 0
      by auto
    have  $(\prod_{x \in X}. a\ x\ powr\ (1 / n)) = (\sum_{x \in X}. (1 / n) * a\ x) \longleftrightarrow$ 

```

```

       $X \neq \{\}$   $\wedge (\forall x \in X. \forall y \in X. a\ x = a\ y)$ 
    using assms False by (intro weighted-arithmetic-geometric-mean-eq-iff) auto
  also have  $(\prod_{x \in X}. a\ x\ \text{powr}\ (1 / n)) = (\prod_{x \in X}. a\ x)\ \text{powr}\ (1 / n)$ 
    using nonneg by (subst powr-sum-distrib-real-left) auto
  also have  $\dots = \text{root}\ n\ (\prod_{x \in X}. a\ x)$ 
    using  $\langle n > 0 \rangle$  nonneg by (subst root-powr-inverse') (auto simp: prod-nonneg)
  also have  $(\sum_{x \in X}. (1 / n) * a\ x) = (\sum_{x \in X}. a\ x) / n$ 
    by (subst sum-distrib-left [symmetric]) auto
  finally show ?thesis using False by auto
qed (auto simp: n-def)

```

1.4 The Binary Version

For convenience, we also derive versions for only two numbers:

corollary *weighted-arithmetic-geometric-mean-binary*:

```

  fixes  $w1\ w2\ x1\ x2 :: \text{real}$ 
  assumes  $x1 \geq 0\ x2 \geq 0\ w1 \geq 0\ w2 \geq 0\ w1 + w2 = 1$ 
  shows  $x1\ \text{powr}\ w1 * x2\ \text{powr}\ w2 \leq w1 * x1 + w2 * x2$ 
proof -
  let  $?a = \lambda b. \text{if } b \text{ then } x1 \text{ else } x2$ 
  let  $?w = \lambda b. \text{if } b \text{ then } w1 \text{ else } w2$ 
  from assms have  $(\prod_{x \in UNIV}. ?a\ x\ \text{powr}\ ?w\ x) \leq (\sum_{x \in UNIV}. ?w\ x * ?a\ x)$ 
    by (intro weighted-arithmetic-geometric-mean) (auto simp add: UNIV-bool)
  thus ?thesis by (simp add: UNIV-bool add-ac mult-ac)
qed

```

corollary *weighted-arithmetic-geometric-mean-eq-iff-binary*:

```

  fixes  $w1\ w2\ x1\ x2 :: \text{real}$ 
  assumes  $x1 \geq 0\ x2 \geq 0\ w1 > 0\ w2 > 0\ w1 + w2 = 1$ 
  shows  $x1\ \text{powr}\ w1 * x2\ \text{powr}\ w2 = w1 * x1 + w2 * x2 \longleftrightarrow x1 = x2$ 
proof -
  let  $?a = \lambda b. \text{if } b \text{ then } x1 \text{ else } x2$ 
  let  $?w = \lambda b. \text{if } b \text{ then } w1 \text{ else } w2$ 
  from assms have  $(\prod_{x \in UNIV}. ?a\ x\ \text{powr}\ ?w\ x) = (\sum_{x \in UNIV}. ?w\ x * ?a\ x)$ 
     $\longleftrightarrow (UNIV :: \text{bool set}) \neq \{\} \wedge (\forall x \in UNIV. \forall y \in UNIV. ?a\ x =$ 
     $?a\ y)$ 
    by (intro weighted-arithmetic-geometric-mean-eq-iff) (auto simp add: UNIV-bool)
  thus ?thesis by (auto simp: UNIV-bool add-ac mult-ac)
qed

```

corollary *arithmetic-geometric-mean-binary*:

```

  fixes  $x1\ x2 :: \text{real}$ 
  assumes  $x1 \geq 0\ x2 \geq 0$ 
  shows  $\text{sqrt}\ (x1 * x2) \leq (x1 + x2) / 2$ 
  using weighted-arithmetic-geometric-mean-binary[of  $x1\ x2\ 1/2\ 1/2$ ] assms
  by (simp add: powr-half-sqrt field-simps real-sqrt-mult)

```

corollary *arithmetic-geometric-mean-eq-iff-binary*:

```

  fixes  $x1\ x2 :: \text{real}$ 

```

```

assumes  $x1 \geq 0 \ x2 \geq 0$ 
shows  $\text{sqrt } (x1 * x2) = (x1 + x2) / 2 \longleftrightarrow x1 = x2$ 
using weighted-arithmetic-geometric-mean-eq-iff-binary[of x1 x2 1/2 1/2] assms
by (simp add: powr-half-sqrt field-simps real-sqrt-mult)

end

```

References

- [1] J. M. Steele. *The Cauchy–Schwarz Master Class: An Introduction to the Art of Mathematical Inequalities*. Cambridge University Press, 2004.