

Satisfiability Encodings of DP Techniques for Maximal Completion

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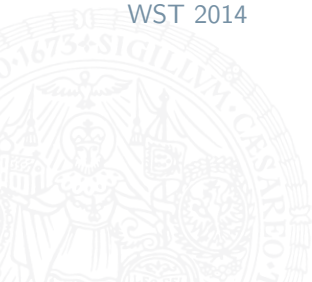




H. Sato and S. Winkler

A Satisfiability Encoding of Dependency Pair Techniques for
Maximal Completion

WST 2014



Outline

(Maximal) Completion

Encodings

Experiments

Conclusion



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Knuth-Bendix Completion: Aim

Given input equalities \mathcal{E} , construct complete TRS \mathcal{R} with $\leftrightarrow_{\mathcal{E}}^* = \leftrightarrow_{\mathcal{R}}^*$



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Group Theory

$$\mathcal{E}: \quad x \cdot 1 \approx x \qquad x \cdot x^{-} \approx 1 \qquad x \cdot (y \cdot z) \approx (x \cdot y) \cdot z$$

$$\mathcal{R}: \quad \begin{array}{lll} x \cdot 1 \rightarrow x & 1 \cdot x \rightarrow x & x \cdot (x^{-} \cdot y) \rightarrow y \\ x \cdot x^{-} \rightarrow 1 & x^{-} \cdot x \rightarrow 1 & (x \cdot y) \cdot z \rightarrow x \cdot (y \cdot z) \\ 1^{-} \rightarrow 1 & x^{-} \rightarrow x & x^{-} \cdot (x \cdot y) \rightarrow y \\ (x \cdot y)^{-} \rightarrow y^{-} \cdot x^{-} & & \end{array}$$

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$$\mathcal{R}: \quad \begin{array}{lll} x \cdot 1 \rightarrow x & 1 \cdot x \rightarrow x & x \cdot (x^{-1} \cdot y) \rightarrow y \\ x \cdot x^{-1} \rightarrow 1 & x^{-1} \cdot x \rightarrow 1 & (x \cdot y) \cdot z \rightarrow x \cdot (y \cdot z) \\ 1^{-1} \rightarrow 1 & x^{-1-1} \rightarrow x & x^{-1} \cdot (x \cdot y) \rightarrow y \\ (x \cdot y)^{-1} \rightarrow y^{-1} \cdot x^{-1} & Y \rightarrow Z \end{array}$$

PROPOSITION 1.31: (Cancellation laws) Let $(G, *)$ be a group, and let $x, y, z \in G$ such that either $x * y = x * z$ or $y * x = z * x$. Then $y = z$.

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WARNING 1.17: Let x and y be elements of a multiplicatively-written group G . Then $(xy)^2$ is not always the same as x^2y^2 .

Maximal Completion

Definition

for input ES \mathcal{E} and ES \mathcal{C} ,

$$\varphi(\mathcal{C})$$



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for input ES \mathcal{E} and ES \mathcal{C} ,

$$\varphi(\mathcal{C}) = \begin{cases} \mathcal{R} & \text{if } \mathcal{E} \cup \text{CP}(\mathcal{R}) \subseteq \downarrow_{\mathcal{R}} \text{ for some } \mathcal{R} \in \mathfrak{R}(\mathcal{C}) \end{cases}$$

where $\mathfrak{R}(\mathcal{C})$ consists of terminating TRSs \mathcal{R} such that $\mathcal{R} \subseteq \mathcal{C} \cup \mathcal{C}^{-1}$



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Theorem (KH11)

If $\varphi(\mathcal{E}) = \mathcal{R}$ then \mathcal{R} is complete and $\leftrightarrow_{\mathcal{E}}^* = \leftrightarrow_{\mathcal{R}}^*$

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Maxcomp

finds $\mathfrak{R}(\mathcal{C})$ using MAXSAT call on

$$\bigvee_{s \approx t \in \mathcal{C}} [s > t] \vee [t > s]$$

encoding of reduction order $>$
(e.g. LPO or KBO)

Maxcomp: Drawback

Experiments

115 systems from Maxcomp distribution

method	#	avg. time
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Commuting group endomorphisms (CGE₂)

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Aim: SAT encoding of termination techniques
with enough power

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find maximal terminating TRS \mathcal{R} in ES \mathcal{C}

symmetric: $\mathcal{C} = \mathcal{C} \cup \mathcal{C}^{-1}$

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find maximal terminating TRS \mathcal{R} in ES \mathcal{C} , and

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X_f^{def} for $f \in \mathcal{F}$

$\alpha(X_f^{\text{def}}) = \top$ if f is defined symbol

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rules trigger DPs

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reduction pair processors orient DPs weakly, and discard only strictly oriented ones

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finally no DP remains

Definitions

- ▶ dependency pair candidates

$$\text{DPC}(\mathcal{C}) = \{l^\# \rightarrow u^\# \mid l \approx r \in \mathcal{C} \text{ is rewrite rule, } r \triangleright u \text{ but } l \not\triangleright u\}$$

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- ▶ DP problem encoding $D = (\mathcal{S}, \mathcal{W}, \phi)$ consists of

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for pair of TRSs $(\mathcal{P}, \mathcal{R})$, and formula ϕ

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- ▶ for assignment α let TRS $\alpha(\mathcal{S}) = \{\ell \rightarrow r \mid S_{\ell \rightarrow r} \in \mathcal{S}, \alpha(S_{\ell \rightarrow r}) = \top\}$

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- ▶ **DP processor encoding Proc** maps D to set of DP problem encodings $\text{Proc}(D) = \{D_1, \dots, D_n\}$

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- ▶ for assignment α let TRS $\alpha(\mathcal{S}) = \{\ell \rightarrow r \mid S_{\ell \rightarrow r} \in \mathcal{S}, \alpha(S_{\ell \rightarrow r}) = \top\}$
- ▶ assignment α is **finite for** $D = (\mathcal{S}, \mathcal{W}, \phi)$ if $\alpha(\phi) = \top$ and DP problem $(\alpha(\mathcal{S}), \alpha(\mathcal{W}))$ is finite
- ▶ **DP processor encoding Proc** maps D to set of DP problem encodings $\text{Proc}(D) = \{D_1, \dots, D_n\}$
- ▶ DP processor encoding Proc is **sound** if

$$\left. \begin{array}{l} \text{Proc}(D) = \{D_1, \dots, D_n\} \\ \alpha \text{ finite for all } D_i \end{array} \right\} \implies \alpha \text{ finite for } D$$

Definition

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Example (CGE₂ snippet)

$$\begin{array}{ll} \mathcal{C}: & 1: f(x^-) \approx f(x)^- \quad \text{DPC}(\mathcal{C}): \quad 3: f^\#(x^-) \rightarrow f(x)^{-\#} \quad 4: f^\#(x^-) \rightarrow f^\#(x) \\ & 2: f(x)^- \approx f(x^-) \quad \quad \quad \quad 5: f(x)^{-\#} \rightarrow f^\#(x^-) \quad 6: f(x)^{-\#} \rightarrow x^{-\#} \end{array}$$

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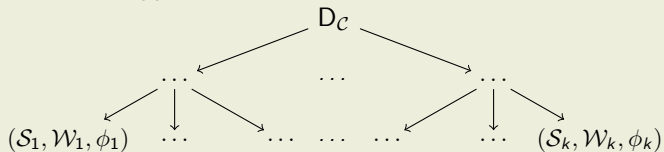
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$$\alpha_2(x) = \top \text{ iff } x \in \{X_1^{\text{rule}}, W_1, X_f^{\text{def}}, S_3, S_4, S_5\}$$

Lemma

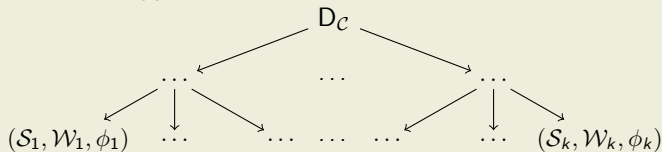
Let \mathcal{C} be an ES. Suppose there is a tree



where for every non-leaf node D with children D_1, \dots, D_n there is sound processor encoding Proc such that $\text{Proc}(D) = \{D_1, \dots, D_n\}$.

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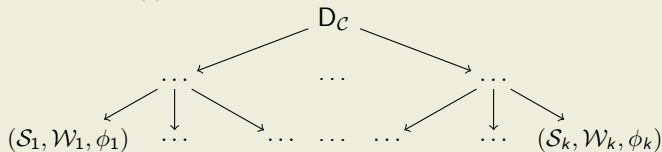
If α satisfies

$$\phi = \bigwedge_{i=1}^k \phi_i \wedge \bigwedge_{S_s \rightarrow t \in \mathcal{S}_i} \neg S_s \rightarrow t$$

then $\mathcal{R} = \{\ell \rightarrow r \mid \alpha(X_{\ell \rightarrow r}^{\text{rule}}) = \top\}$ is terminating.

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Proof Idea

α is finite for all nodes, in particular the root. Termination of \mathcal{R} follows from correctness of the DP framework [GTSK05].

Definition (Reduction pair processor)

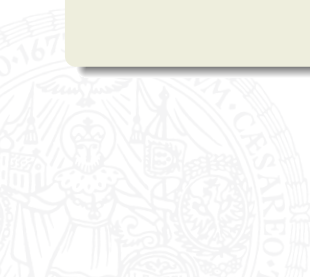
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Definition (Reduction pair processor **with rule removal**)

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Luckily ...

encodings $[\cdot \geq \cdot]$ for RPO, KBO, WPO (+ argument filterings, usable rules), polynomial, matrix interpretations, ... are well studied

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$$\begin{aligned}
 T_{\mathcal{W}} = & (W_1 \rightarrow [f(x^-) \geq f(x)^-] \wedge W'_1) \wedge \\
 & (W_2 \rightarrow [f(x)^- \geq f(x^-)] \wedge W'_2)
 \end{aligned}$$

DG with Cycle Analysis

Idea

assign number X_p^w to DP p such that cycle $p_1 \rightarrow p_2 \rightarrow \dots \rightarrow p_n \rightarrow p_1$ issues unsatisfiable constraint $X_{p_1}^w > X_{p_2}^w > \dots > X_{p_n}^w > X_{p_1}^w$



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where $X_{s \rightarrow t, u \rightarrow v}^{\text{edge}} = \text{root}(t) = \text{root}(u) ? \top : \perp$

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$$T_S = \bigwedge_{S_{p_1}, S_{p_2} \in \mathcal{S}} S_{p_1} \wedge S_{p_2} \wedge X_{p_1, p_2}^{\text{edge}} \wedge \neg S'_{p_1} \wedge \neg S'_{p_2} \rightarrow X_{p_1}^w > X_{p_2}^w$$

$$T_W = \bigwedge_{S_{\ell \rightarrow r} \in \mathcal{S}} S_{\ell \rightarrow r} \rightarrow W'_{\ell \rightarrow r} \wedge \bigwedge_{W_{\ell \rightarrow r} \in \mathcal{W}} W_{\ell \rightarrow r} \rightarrow W'_{\ell \rightarrow r}$$

where $X_{s \rightarrow t, u \rightarrow v}^{\text{edge}} = \text{root}(t) = \text{root}(u) ? \top : \perp$

Example (CGE₂ snippet)

\mathcal{C} : 1: $f(x^-) \approx f(x)^-$ DPC(\mathcal{C}): 3: $f^\#(x^-) \rightarrow f(x)^{-\#}$ 4: $f^\#(x^-) \rightarrow f^\#(x)$
 2: $f(x)^- \approx f(x^-)$ 5: $f(x)^{-\#} \rightarrow f^\#(x^-)$ 6: $f(x)^{-\#} \rightarrow x^{-\#}$

$D_{\mathcal{C}} = (\mathcal{S}, \mathcal{W}, \phi)$ for $\mathcal{S} = \{S_3, S_4, S_5, S_6\}$, $\mathcal{W} = \{W_1, W_2\}$

Example (CGE₂ snippet)

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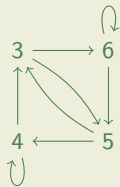
Proc_{DG} maps $D_{\mathcal{C}}$ to $D = (\mathcal{S}', \mathcal{W}', \phi \wedge T_S \wedge T_W)$, where
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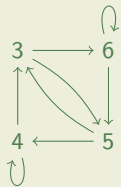
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$$T_S = (S_4 \wedge S_3 \wedge \neg S'_4 \wedge \neg S'_3 \rightarrow X_4^w > X_3^w)$$



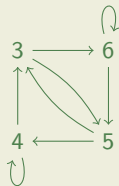
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$$\begin{aligned}
 T_S = & (S_4 \wedge S_3 \wedge \neg S'_4 \wedge \neg S'_3 \rightarrow X_4^w > X_3^w) \wedge \\
 & (S_3 \wedge S_5 \wedge \neg S'_3 \wedge \neg S'_5 \rightarrow X_3^w > X_5^w)
 \end{aligned}$$



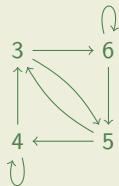
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$$\begin{aligned}
 T_S = & (S_4 \wedge S_3 \wedge \neg S'_4 \wedge \neg S'_3 \rightarrow X_4^w > X_3^w) \wedge \\
 & (S_3 \wedge S_5 \wedge \neg S'_3 \wedge \neg S'_5 \rightarrow X_3^w > X_5^w) \wedge \\
 & (S_5 \wedge S_4 \wedge \neg S'_5 \wedge \neg S'_4 \rightarrow X_5^w > X_4^w) \wedge \dots
 \end{aligned}$$



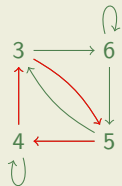
Example (CGE₂ snippet)

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 T_S = & (S_4 \wedge S_3 \wedge \neg S'_4 \wedge \neg S'_3 \rightarrow X_4^w > X_3^w) \wedge \\
 & (S_3 \wedge S_5 \wedge \neg S'_3 \wedge \neg S'_5 \rightarrow X_3^w > X_5^w) \wedge \\
 & (S_5 \wedge S_4 \wedge \neg S'_5 \wedge \neg S'_4 \rightarrow X_5^w > X_4^w) \wedge \dots
 \end{aligned}$$



Example (CGE₂ snippet)

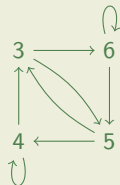
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 T_S = & (S_4 \wedge S_3 \wedge \neg S'_4 \wedge \neg S'_3 \rightarrow X_4^w > X_3^w) \wedge \\
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 & (S_5 \wedge S_4 \wedge \neg S'_5 \wedge \neg S'_4 \rightarrow X_5^w > X_4^w) \wedge \dots
 \end{aligned}$$

$$T_W = \bigwedge_{i=3}^6 S_i \rightarrow W'_i \wedge \bigwedge_{i=1}^2 W_i \rightarrow W'_i$$



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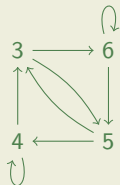
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$$T_S = (S_4 \wedge S_3 \wedge \neg S'_4 \wedge \neg S'_3 \rightarrow X_4^w > X_3^w) \wedge$$

$$(S_3 \wedge S_5 \wedge \neg S'_3 \wedge \neg S'_5 \rightarrow X_3^w > X_5^w) \wedge$$

$$(S_5 \wedge S_4 \wedge \neg S'_5 \wedge \neg S'_4 \rightarrow X_5^w > X_4^w) \wedge \dots$$

$$T_W = \bigwedge_{i=3}^6 S_i \rightarrow W'_i \wedge \bigwedge_{i=1}^2 W_i \rightarrow W'_i$$



Proc_k: DG with k SCCs

Idea: limit #SCCs to at most k , and assign SCC number X_p^{SCC} to DP p

Example (CGE₂ snippet)

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 2: $f(x)^- \approx f(x^-)$ 5: $f(x)^{-\#} \rightarrow f^\#(x^-)$ 6: $f(x)^{-\#} \rightarrow x^{-\#}$

$D_{\mathcal{C}} = (\mathcal{S}, \mathcal{W}, \phi)$ for $\mathcal{S} = \{S_3, S_4, S_5, S_6\}$, $\mathcal{W} = \{W_1, W_2\}$

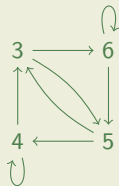
Proc_{DG} maps $D_{\mathcal{C}}$ to $D = (\mathcal{S}', \mathcal{W}', \phi \wedge T_S \wedge T_W)$, where
 $\mathcal{S}' = \{S_3, \dots, S_6\}$, $\mathcal{W}' = \{W_1, \dots, W_6\}$, and

$$T_S = (S_4 \wedge S_3 \wedge \neg S'_4 \wedge \neg S'_3 \rightarrow X_4^w > X_3^w) \wedge$$

$$(S_3 \wedge S_5 \wedge \neg S'_3 \wedge \neg S'_5 \rightarrow X_3^w > X_5^w) \wedge$$

$$(S_5 \wedge S_4 \wedge \neg S'_5 \wedge \neg S'_4 \rightarrow X_5^w > X_4^w) \wedge \dots$$

$$T_W = \bigwedge_{i=3}^6 S_i \rightarrow W'_i \wedge \bigwedge_{i=1}^2 W_i \rightarrow W'_i$$



Lemma

$\text{Proc}_{(>, \geq)}$, Proc_{DG} and Proc_k are sound

Example (CGE₂ in practice)

iteration 1: 6 equalities

$*(*(x,y),z) = *(x,*(y,z)), *(-x),x) = 1, *(1,x) = x, *(f(x),f(y)) = f(*(x,y)),$
 $*(g(x),g(y)) = g(*(x,y)), *(f(x),g(y)) = *(g(y),f(x))$

Example (CGE₂ in practice)

iteration 2: 19 equalities

$*(f(1),f(x)) = f(x)$, $*(f(-(x)),f(x)) = f(1)$, $*(x,y) = *(* (x,1),y)$, $*(* (-$
 $(*(x,y)),x),y) = 1$, $*(x,1) = *(* (x,-(y)),y)$, $*(f(g(x)),f(g(y))) = f(g(* (x,y)))$,
 $*(x,g(* (y,z))) = *(* (x,g(y)),g(z))$, $*(* (x,g(y)),f(z)) = *(* (x,f(z)),g(y))$,
 $*(f(f(x)),f(g(y))) = *(f(g(y)),f(f(x)))$, $x = *(-(y),*(y,x))$, $*(f(x),*(f(y),z)) =$
 $*(f(* (x,y)),z)$, $*(g(x),*(g(y),z)) = *(g(* (x,y)),z)$, $*(* (x,y),z) = *(x,*(y,z))$, $*(-$
 $(x),x) = 1$, $*(1,x) = x$, $*(f(x),f(y)) = f(* (x,y))$, $*(g(x),g(y)) = g(* (x,y))$,
 $*(f(x),g(y)) = *(g(y),f(x))$, $*(* (f(x),g(y)),f(z)) = *(* (f(x),f(z)),g(y))$

Example (CGE₂ in practice)

iteration 3: 45 equalities

$$\begin{aligned}
&*(-(g(x)),*(g(*(x,y)),z)) = *(g(y),z), *(f(x),*(g(y),z)) = *(g(y),*(f(x),z)), \\
&*(-(f(x)),*(f(*(x,y)),z)) = *(f(y),z), *(f(x),*(g(y),f(z))) = *(g(y),f(*(x,z))), \\
&*(x,*(-(g(y)),g(*(y,z)))) = *(x,g(z)), *(-(*(x,g(y))),*(x,g(*(y,z)))) = g(z), x \\
&= *(-1),x), x = *(-(-(x)),1), *(x,y) = *(-(-(x)),y), *(-(f(x)),f(*(x,y))) = \\
&f(y), *(-(g(x)),g(*(x,y))) = g(y), x = *(-(*(y,z)),*(y,*(z,x))), *(f(x),y) = \\
&*(f(1),*(f(x),y)), *(-(f(x)),*(g(y),f(x))) = g(y), *(f(1),x) = *(f(-y)),*(f(y),x)), \\
&*(x,y) = **(*(x,-1),y), *(x,*(-1),y) = *(x,y), x = **(-(*(y,1)),y),x), x \\
&= *(-(*(y,1)),*(y,x)), *(x,y) = **(*(x,-(-y))),1), *(x,*(-(-y)),1) = *(x,y), \\
&x = **(-(*(y,-(x))),y),1), x = *(-(*(y,-(x))),*(y,1)), **(-*(g(x)),f(y)),g(x)) \\
&= f(y), **(*(x,-1),-(*(x,y))),x),y) = 1, *(f(1),f(x)) = f(x), *(f(-x)),f(x) \\
&= f(1), *(x,y) = **(*(x,1),y), **(-(*(x,y)),x),y) = 1, *(x,1) = **(*(x,- \\
&(y)),y), *(f(g(x)),f(g(y))) = f(g(*(x,y))), *(x,g(*(y,z))) = **(*(x,g(y)),g(z)), \\
&**(*(x,g(y)),f(z)) = **(*(x,f(z)),g(y)), *(f(f(x)),f(g(y))) = *(f(g(y)),f(f(x))), \\
&x = *(-y),*(y,x)), *(f(x),*(f(y),z)) = *(f(*(x,y)),z), *(g(x),*(g(y),z)) = \\
&*(g(*(x,y)),z), **(*(x,y),z) = *(x,*(y,z)), *(-x),x) = 1, *(1,x) = x, *(f(x),f(y)) \\
&= f(*(x,y)), *(g(x),g(y)) = g(*(x,y)), *(f(x),g(y)) = *(g(y),f(x)), *(- \\
&(*(x,g(y))),*(x,*(g(y),f(z)))) = f(z), *(f(x),*(g(y),f(z))) = *(g(y),f(*(x,z)))
\end{aligned}$$

Example (CGE₂ in practice)

iteration 8: 282 equalities

$\mathcal{R}_{8,2}$ consists of 158 rules:

$\begin{aligned} & f(x, \neg(x, x)) \rightarrow 1 \\ & f(f(\neg(x, x)), x) \rightarrow 1 \\ & \neg(f(\neg(x, x)), y) \rightarrow \neg(f(y), f(x)) \\ * & (x, \neg(\neg(x, y), x)) * (y, x) \rightarrow 1 \\ & \neg(\neg(x, y)) \rightarrow \neg(\neg(y), \neg(x)) \\ & \neg(\neg(\neg(x, y), x)) \rightarrow y \\ & * f(x), * f(\neg(x, x)), x) \rightarrow y \\ & \neg(\neg(x)) \rightarrow x \\ & \neg(1) \rightarrow 1 \\ & f(1) \rightarrow 1 \\ & * (x, 1) \rightarrow x \\ & * (\neg(\neg(\neg(x, \neg(y)), 1)), y) \rightarrow * (x, 1) \\ & * (g(x), * f(y), z) \rightarrow * f(y), * (g(x), z) \\ & * (\neg(x, y)), * (x, * (y), z) \rightarrow x \\ & * (\neg(* (x, 1), x), y) \rightarrow y \\ & * (x, * (y, 1)) \rightarrow * (x, y) \\ & * (x, \neg(y), y) \rightarrow * (x, 1) \\ * & (x, * (y, \neg(x, * (x, y)))) \rightarrow \neg(x) \\ & * (x, \neg(\neg(x, * (x), z)) \rightarrow \neg(x, y, z) \\ & \neg(g(\neg(x, \neg(y))), x) \rightarrow * (x, \neg(x)) \\ & \neg(g(\neg(x, \neg(y))), y) \rightarrow g(\neg(x, \neg(y)), f(x)) \\ * & (g(x), * f(y), \neg(g(x))) \rightarrow f(\neg(y)) \\ & f(x, \neg(\neg(y, x))) \rightarrow f(\neg(y)) \\ * & (x, * (y, \neg(x, y))) \rightarrow 1 \\ & \neg(f(x)) \rightarrow \neg(x) \\ & f(\neg(g(x))) \rightarrow f(g(x)) \\ & * (\neg(f(x)), f(y)) \rightarrow f(x, y) \\ * f & (x, y), \neg(f(y)) \rightarrow f(x) \\ & * (\neg(f(x)), * f(x, y)) \rightarrow * f(x, y) \\ & * (\neg(f(x)), * f(y), z) \rightarrow * f(x, y), z) \\ & * (\neg(f(x)), * (g(y), f(x), z))) \rightarrow * (g(y), f(x)) \\ & * (\neg(x, \neg(x), x), y) \rightarrow 1 \\ & * (\neg(x, \neg(x)), 1) \rightarrow * (x, x) \\ & * (\neg(1), 1) \rightarrow f(x) \\ * (\neg & (x, g(y))), * (x, * (y), z)) \rightarrow g(z) \\ & * (\neg(x, g(y)), * (x, * (y), z)) \rightarrow * (x) \\ & * (x, \neg(\neg(x), y)) \rightarrow * (x, y) \\ & * (x, \neg(\neg(x), 1)) \rightarrow * (x, x) \\ & * (f(x), f(x)) \rightarrow f(1) \\ & * (f(x), f(y), z) \rightarrow * f(x, y), z) \\ & * (f(x), f(y)) \rightarrow * (x, y) * (g(x), g(y)) \end{aligned}$	$\begin{aligned} & g(\neg(\neg(x, x)) \rightarrow 1 \\ & * (\neg(x, * (y, \neg(z))), * (x, y), z) \rightarrow 1 \\ & \neg(* (x, f(y))) \rightarrow * f(\neg(y), \neg(x)) \\ & * (\neg(x, g(\neg(y))), * (g(y), \neg(x)) \\ & \neg(\neg(x), y)) \rightarrow \neg(\neg(y), x) \\ & * f(1), x) \rightarrow x \\ & * f(x), * f(1), y) \rightarrow * f(x), y \\ * f & (x, y), * (\neg(y), z) \rightarrow * f(x), z) \\ & \neg(1) \rightarrow 1 \\ & * (\neg(x, \neg(1)), * (x, y)) \rightarrow y \\ & * (\neg(x, \neg(y)), x) \rightarrow y \\ & f(x, * (\neg(y), y)) \rightarrow f(x) \\ & * (\neg(1), x) \rightarrow x \\ & * f(1), * f(x), y) \rightarrow * f(x), y \\ & * (\neg(x, 1), * (x, y)) \rightarrow y \\ & * (\neg(\neg(x, \neg(y)), x), 1) \rightarrow y \\ & * (x, * (g(y), g(z))) \rightarrow * (x, g(y, z)) \\ & \neg(g(f(f(f(x)))) \rightarrow g(f(f(f(\neg(x)))))) \\ & * (x, \neg(f(y, z, x))) \rightarrow \neg(x, y, z) \\ & * (\neg(* (g(x), f(y))), f(y)) \rightarrow \neg(g(x)) \\ & * (g(x), \neg(g(y)) \rightarrow g(\neg(x, \neg(y))) \\ & \neg(g(\neg(x, y))) \rightarrow g(\neg(\neg(y), x)) \\ & g(x, \neg(x, y)) \rightarrow g(\neg(y)) \\ & \neg(g(\neg(x))) \rightarrow g(x) \\ & \neg(g(x)) \rightarrow \neg(x) \\ & * (\neg(x, f(y)), x) \rightarrow f(\neg(y)) \\ & * f(x), * (f(1), y) \rightarrow * f(x, y) \\ & * (\neg(f(x)), \neg(f(y)) \rightarrow \neg(f(x)) \\ & * f(x, y), \neg(f(y)) \rightarrow f(x) \\ & * (\neg(f(x)), * f(x, y)) \rightarrow * f(x, y) \\ & * (\neg(f(x)), * f(y), z) \rightarrow * f(x, y), z) \\ & * (\neg(f(x)), * (g(y), f(x), z))) \rightarrow * (g(y), f(x)) \\ & * (\neg(x, \neg(x), x), y) \rightarrow 1 \\ & * (\neg(x, \neg(x)), 1) \rightarrow * (x, x) \\ & * (\neg(1), 1) \rightarrow f(x) \\ * (\neg & (x, g(y))), * (x, * (y), z)) \rightarrow g(z) \\ & * (\neg(x, g(y)), * (x, * (y), z)) \rightarrow * (x) \\ & * (x, \neg(\neg(x), y)) \rightarrow * (x, y) \\ & * (x, \neg(\neg(x), 1)) \rightarrow * (x, x) \\ & * (f(x), f(x)) \rightarrow f(1) \\ & * (f(x), f(y), z) \rightarrow * f(x, y), z) \\ & * (f(x), f(y)) \rightarrow * (x, y) * (g(x), g(y)) \end{aligned}$	$\begin{aligned} & * (x, * (\neg(y), y)) \rightarrow x \\ & * (\neg(* (x, y), * (x, z)) \rightarrow * (\neg(y), z) \\ & \neg(* f(x, y)) \rightarrow * (\neg(y), f(\neg(x))) \\ & f(x, * (\neg(\neg(x, x))) \rightarrow f(y) \\ & \neg(x, \neg(y)) \rightarrow * (y, \neg(x)) \\ & * (g(1), x) \rightarrow x \\ & g(1) \rightarrow 1 \\ * f & (x), * f(y), \neg(\neg(z)) \rightarrow * f(x), * f(y), z) \\ & g(\neg(1)) \rightarrow 1 \\ & * f(x, * (y), y) \rightarrow * f(x), y \\ & * (\neg(x, \neg(y)), * (x, z)) \rightarrow * (y, z) \\ & * (x, * (\neg(x), y)) \rightarrow y \\ & * (\neg(x, 1), 1) \rightarrow x \\ * f & (\neg(x)), * f(x), y) \rightarrow * f(1), y \\ & * (x, * (\neg(y), 1)) \rightarrow * (x, y) \\ & * (\neg(x, \neg(y)), * (x, 1)) \rightarrow y \\ & * (\neg(x), * (x, y)) \rightarrow y \\ * (g & (x), \neg(* (g(x), f(y))) \rightarrow f(\neg(y)) \\ & * (\neg(x), x) \rightarrow 1 \\ & * (\neg(g(x), f(y)) \rightarrow * f(y), \neg(g(x)) \\ & * (\neg(g(x, y)), g(x)) \rightarrow \neg(g(y)) \\ & * (g(\neg(x, y)), * (x, x)) \rightarrow g(y) \\ & * (g(x, y), \neg(g(x))) \rightarrow g(x) \\ & \neg(g(\neg(x))) \rightarrow g(f(x)) \\ & \neg(f(x)) \rightarrow f(f(\neg(x))) \\ & \neg(f(x)) \rightarrow f(x) \\ & * (\neg(f(x)), f(x, y, z)) \rightarrow f(x, * (y, z)) \\ & * f(x, y), \neg(\neg(z)) \rightarrow * f(x, y), z) \\ & * (\neg(x)), * f(x), y) \rightarrow * f(x), y \\ & * (\neg(f(x)), * f(x, y), z) \rightarrow * f(x, y), z) \\ & * (\neg(f(x)), * (g(y), f(x), z))) \rightarrow * (g(y), f(x)) \\ & * (\neg(x, \neg(x), x), y) \rightarrow 1 \\ & * (\neg(x, \neg(x)), 1) \rightarrow * (x, x) \\ & * (\neg(1), 1) \rightarrow f(x) \\ * (\neg & (x, g(y))), * (x, * (y), z)) \rightarrow g(z) \\ & * (\neg(x, g(y)), * (x, * (y), z)) \rightarrow * (x) \\ & * (x, \neg(\neg(x), y)) \rightarrow * (x, y) \\ & * (x, \neg(\neg(x), 1)) \rightarrow * (x, x) \\ & * (f(x), f(x)) \rightarrow f(1) \\ & * (f(x), f(y), z) \rightarrow * f(x, y), z) \\ & * (f(x), f(y)) \rightarrow * (x, y) * (g(x), g(y)) \end{aligned}$	$\begin{aligned} & f(g(\neg(\neg(x, x))) \rightarrow 1 \\ & * (\neg(x, f(\neg(y))) \rightarrow * f(y), \neg(x)) \\ * (x, & * (y, \neg(\neg(x, y), z)) \rightarrow y \\ & * (x, * (y, \neg(x), y)) \rightarrow * (y, x) \\ & * (x, \neg(1)) \rightarrow x \\ & f(1) \rightarrow 1 \\ * f & (x, y), * (\neg(f(x), z)) \rightarrow * (\neg(f(\neg(x))), z) \\ & \neg(f(1)) \rightarrow 1 \\ & * f(x, * (\neg(x), x)) \rightarrow * (x, y) \\ & * f(x, * (y), z) \rightarrow * (x, y) \\ & * (x, 1), y) \rightarrow * (x, y) \\ & * (\neg(x), y) \rightarrow * (x, y) \\ & * (x, \neg(1), y) \rightarrow * (x, y) \\ & * (\neg(y), y) \rightarrow * (x, y) \\ & * (x, 1), y) \rightarrow * (x, y) \\ & * (x, \neg(y), y) \rightarrow * (x, g(x)) \\ & * (\neg(g(x), g(x))) \rightarrow 1 \\ & \neg(g(x, f(\neg(y)))) \rightarrow g(\neg(x, y), \neg(x)) \\ * (g & (x), \neg(g(x, y))) \rightarrow \neg(g(y)) \\ & * (x, \neg(y), 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Example (CGE₂ in practice)

iteration 8: 282 equalities

 $\mathcal{R}_{8,2}$ reduced:

$$\begin{array}{lll}
 1 \cdot x \rightarrow x & f(x) \cdot f(y) \rightarrow f(x \cdot y) & x \cdot (y \cdot z) \rightarrow (x \cdot y) \cdot z \\
 x \cdot 1 \rightarrow x & f(1) \rightarrow 1 & (x \cdot y) \cdot y^- \rightarrow x \\
 x^- \cdot x \rightarrow 1 & f(x)^- \rightarrow f(x^-) & (x \cdot y^-) \cdot y \rightarrow x \\
 x \cdot x^- \rightarrow 1 & g(x) \cdot g(y) \rightarrow g(x \cdot y) & f(x) \cdot (f(y) \cdot z) \rightarrow f(x \cdot y) \cdot z \\
 1^- \rightarrow 1 & (x \cdot y)^- \rightarrow y^- \cdot x^- & g(x) \cdot (g(y) \cdot z) \rightarrow g(x \cdot y) \cdot z \\
 x^{--} \rightarrow x & g(x)^- \rightarrow g(x^-) & g(x) \cdot (f(y) \cdot z) \rightarrow f(x) \cdot (g(y) \cdot z) \\
 g(1) \rightarrow 1 & g(x) \cdot f(y) \rightarrow f(y) \cdot g(x) &
 \end{array}$$

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 x^{-} \cdot x \rightarrow 1 & f(x)^{-} \rightarrow f(x^{-}) & (x \cdot y^{-}) \cdot y \rightarrow x \\
 x \cdot x^{-} \rightarrow 1 & g(x) \cdot g(y) \rightarrow g(x \cdot y) & f(x) \cdot (f(y) \cdot z) \rightarrow f(x \cdot y) \cdot z \\
 1^{-} \rightarrow 1 & (x \cdot y)^{-} \rightarrow y^{-} \cdot x^{-} & g(x) \cdot (g(y) \cdot z) \rightarrow g(x \cdot y) \cdot z \\
 x^{-} \rightarrow x & g(x)^{-} \rightarrow g(x^{-}) & g(x) \cdot (f(y) \cdot z) \rightarrow f(x) \cdot (g(y) \cdot z) \\
 g(1) \rightarrow 1 & g(x) \cdot f(y) \rightarrow f(y) \cdot g(x) &
 \end{array}$$

50 arithmetic and 16K boolean variables, 86K clauses ... 5.6 seconds

Outline

(Maximal) Completion

Encodings

Experiments

Conclusion



Implementation in Maxcomp

Supported Termination Strategies

- ▶ base orders: linear polynomials, KBO, LPO with argument filterings
- ▶ DPs, optionally DGs (cycle and SCC variants), followed by arbitrary sequence of reduction pair processors



Implementation in Maxcomp

Supported Termination Strategies

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Optimizations

- ▶ avoid re-submitting same clauses: maintain state of SMT solver throughout completion run, pushing and popping clauses
- ▶ to extend active equation set A , add n new critical pairs and m reduced passive equations, where n and m depend on $|A|$
- ▶ restarts
- ▶ combined `auto` strategy: start with both LPO and DP strategy, but after first iteration keep only one strategy

Experiments

115 systems from Maxcomp distribution, $\infty := 180$ seconds

	method	#	avg. time	CGE ₂	CGE ₃	pr	equiv_proofs
(1)	Maxcomp	85	3.8	∞	∞	∞	∞
(2)	DPs	83	14.7	6.4	∞	∞	1.6
(3)	DG	40	2.8	∞	∞	∞	∞
(4)	DG/2SCCs	25	2.5	∞	∞	∞	∞
(5)	auto	92	10.2	5.6	135.1	150.2	1.5
(6)	KBCV	88	4.4	∞	∞	∞	∞
(7)	mkbTT	83	6.1	6.7	45.17	∞	2.0

- (1) Maxcomp using LPO
- (2) DPs, linear polynomials * 3, LPO + AF
- (3) DPs, DG, linear polynomials * 3, LPO + AF
- (4) DPs, DG/2 SCCs, linear polynomials * 3, LPO + AF
- (5) start with LPO and (2), abandon less successful branch

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<http://cl-informatik.uibk.ac.at/software/maxcompdp/>

Summary

- ▶ Maxcomp can complete CGE_2 , CGE_3 , equiv_proofs, proofreduction
- ▶ lightweight DP techniques are feasible, but graphs (currently) not
- ▶ with some tweaking, Maxcomp wins some problems over other tools

Future Work

- ▶ WPO
- ▶ (try to) improve graph encoding
- ▶ certifiable output
- ▶ more input systems?