

Solved exercises must be marked and solutions (as a single PDF file) uploaded in [OLAT](#). Solutions for bonus exercises must be submitted separately. The (strict) deadline is 7 am on April 30.

Exercises

- (2) 1. Consider the sentence

$$\varphi = \forall x \exists y (P(x, y) \wedge \forall z (P(y, z) \rightarrow \neg P(z, x)))$$

Which of the following models satisfies φ ?

- (a) The model \mathcal{M}_1 consisting of \mathbb{N} and the interpretation $P^{\mathcal{M}_1} = \{(a, b) \mid b = a + 1\}$.
- (b) The model \mathcal{M}_2 consisting of \mathbb{N} and the interpretation $P^{\mathcal{M}_2} = \{(a, b) \mid 2a \neq b\}$.
- (c) The model \mathcal{M}_3 consisting of \mathbb{N} and the interpretation $P^{\mathcal{M}_3} = \{(a, b) \mid a < b\}$.

- (2) 2. Prove the validity of the following sequents using natural deduction:

- (a) $Q(a) \vdash \forall x (x = a \rightarrow Q(x))$
- (b) $\exists x \exists y (P(x, y) \vee P(y, x)), \neg \exists x P(x, x) \vdash \exists x \exists y \neg(x = y)$

For each application of $=e$, make sure to specify φ , t_1 and t_2 .

- (3) 3. For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.

- (a) $\exists x (\neg P(x) \wedge \forall y (x = y \rightarrow \neg Q(y))) \vdash \neg \forall x (P(x) \vee Q(x))$
- (b) $\exists x (\neg P(x) \wedge \forall y (\neg(x = y) \vee \neg Q(y))) \vdash \exists x (P(x) \vee Q(x))$
- (c) $\exists x (P(x) \wedge \exists y (\neg(x = y) \wedge \neg Q(y))) \vdash \neg \forall x (P(x) \wedge Q(x))$

- (3) 4. Prove the validity of the following sequents using natural deduction:

- (a) $\exists x (P \rightarrow Q(x)) \vdash P \rightarrow \exists x Q(x)$
- (b) $\forall x Q(x) \rightarrow P \vdash \exists x (Q(x) \rightarrow P)$
- (c) $P \rightarrow \exists x Q(x) \vdash \exists x (P \rightarrow Q(x))$

Bonus Exercises

- (2) 5. Prove the soundness of propositional resolution: A set S of clauses that admits a refutation is unsatisfiable.
- (3) 6. Prove the completeness of propositional resolution: Every unsatisfiable set S of clauses admits a refutation.