

All Solutions

- 1 No. The initial configuration has an odd number of black coffee beans, and this property is preserved during the game:

$$\begin{array}{ll}
 \bullet \bullet \rightarrow \circ & b \rightsquigarrow b - 2 \\
 \circ \circ \rightarrow \circ & b \rightsquigarrow b \\
 \bullet \circ \rightarrow \bullet & b \rightsquigarrow b \\
 \circ \bullet \rightarrow \bullet & b \rightsquigarrow b
 \end{array}$$

In each move the number of beans decreases by one, and hence the game ends with a single black bean. It follows that the player making the last move wins the game. Since there are 15 beans in the initial configuration, the second player always makes the last move, and thus wins the game.

- 2 (b) The total number of A's and T's in the milk gene is odd (7) whereas it is even (6) in the mad cow retrovirus. Since the five DNA substitutions preserve the parity of this number

$$\begin{array}{ll}
 \text{TCAT} \leftrightarrow \text{T} & \pm 2 \\
 \text{GAG} \leftrightarrow \text{AG} & 0 \\
 \text{CTC} \leftrightarrow \text{TC} & 0 \\
 \text{AGTA} \leftrightarrow \text{A} & \pm 2 \\
 \text{TAT} \leftrightarrow \text{CT} & \pm 2
 \end{array}$$

the milk gene cannot be transformed into the mad cow retrovirus.

- 3 (a) Consider the well-founded relation $R = \{(a, f(a))\}$. We have $a R f(a) \triangleright a$, so $R \cup \triangleright$ is not well-founded. Observe that R lacks closure under contexts because $(f(a), f(f(a))) \notin R$.
- (b) Because $R \cdot \triangleright \subseteq (R \cup \triangleright)^+$, the relation $R \cdot \triangleright$ inherits well-foundedness from $R \cup \triangleright$. Hence $R \cdot \triangleright$ is also irreflexive. It remains to show that $R \cdot \triangleright$ is transitive. In the proof of Lemma 2.1.25 we observed $\triangleright \cdot R \subseteq R \cdot \triangleright$. Therefore $R \cdot \triangleright \cdot R \cdot \triangleright \subseteq R \cdot R \cdot \triangleright \cdot \triangleright \subseteq R \cdot \triangleright$.
- 4 (a) We have $t\sigma = y + (y + (y + y))$ and $\text{Dom}(\sigma) = \{x\}$.
- (b) We have $t\sigma = (y + x) + ((y + y) + ((y + x) + (y + y)))$ and $\text{Dom}(\sigma) = \{x, y, z\}$.
- (c) We have $t\sigma = (0 + z) + (s(0) + ((0 + z) + s(0)))$ and $\text{Dom}(\sigma) = \{x, y, z\}$.

- 5 (a) The term $s(x) + y$ does not match t :

$$\begin{aligned}
 \{s(x) + y \mapsto t\} &\implies \{s(x) \mapsto s(x) + x, y \mapsto s(s(0 + y))\} \\
 &\implies \perp
 \end{aligned}$$

- (b) The term $x + s(y)$ matches t :

$$\begin{aligned}
 \{x + s(y) \mapsto t\} &\implies \{x \mapsto s(x) + x, s(y) \mapsto s(s(0 + y))\} \\
 &\implies \{x \mapsto s(x) + x, y \mapsto s(0 + y)\}
 \end{aligned}$$

(c) The term $(x + y) + x$ does not match t :

$$\begin{aligned} \{(x + y) + x \mapsto t\} &\implies \{x + y \mapsto s(x) + x, x \mapsto s(s(0 + y))\} \\ &\implies \{x \mapsto s(x), y \mapsto x, x \mapsto s(s(0 + y))\} \\ &\implies \perp \end{aligned}$$

6 (a) There are two rewrite sequences leading to the same normal form:

$$\begin{aligned} 2 + (4 + (6 + (4 + 2))) &\rightarrow 2 + (4 + (6 + 6)) \\ &\rightarrow 2 + (4 + (1 : 2)) \\ &\rightarrow 2 + (1 : (4 + 2)) \\ &\rightarrow 2 + (1 : 6) \\ &\rightarrow 1 : (2 + 6) \\ &\rightarrow 1 : 8 \end{aligned}$$

and

$$\begin{aligned} 2 + (4 + (6 + (4 + 2))) &\rightarrow 2 + (4 + (6 + 6)) \\ &\rightarrow 2 + (4 + (1 : 2)) \\ &\rightarrow 2 + (1 : (4 + 2)) \\ &\rightarrow 1 : (2 + (4 + 2)) \\ &\rightarrow 1 : (2 + 6) \\ &\rightarrow 1 : 8 \end{aligned}$$

(b) We add the following 102 rewrite rules:

$0 \times 0 \rightarrow 0$	$1 \times 0 \rightarrow 0$	\dots	$9 \times 0 \rightarrow 0$
$0 \times 1 \rightarrow 0$	$1 \times 1 \rightarrow 1$	\dots	$9 \times 1 \rightarrow 9$
$0 \times 2 \rightarrow 0$	$1 \times 2 \rightarrow 2$	\dots	$9 \times 2 \rightarrow 1 : 8$
$0 \times 3 \rightarrow 0$	$1 \times 3 \rightarrow 3$	\dots	$9 \times 3 \rightarrow 2 : 7$
$0 \times 4 \rightarrow 0$	$1 \times 4 \rightarrow 4$	\dots	$9 \times 4 \rightarrow 3 : 6$
$0 \times 5 \rightarrow 0$	$1 \times 5 \rightarrow 5$	\dots	$9 \times 5 \rightarrow 4 : 5$
$0 \times 6 \rightarrow 0$	$1 \times 6 \rightarrow 6$	\dots	$9 \times 6 \rightarrow 5 : 4$
$0 \times 7 \rightarrow 0$	$1 \times 7 \rightarrow 7$	\dots	$9 \times 7 \rightarrow 6 : 3$
$0 \times 8 \rightarrow 0$	$1 \times 8 \rightarrow 8$	\dots	$9 \times 8 \rightarrow 7 : 2$
$0 \times 9 \rightarrow 0$	$1 \times 9 \rightarrow 9$	\dots	$9 \times 9 \rightarrow 8 : 1$
$x \times (y : z) \rightarrow (x \times y) : (x \times z)$			$(x : y) \times z \rightarrow (x \times z) : (y \times z)$

We compute $11 \times 46 = 506$ as follows:

$$\begin{aligned} (1:1) \times (4:6) &\rightarrow ((1:1) \times 4) : ((1:1) \times 6) \rightarrow ((1 \times 4) : (1 \times 4)) : ((1:1) \times 6) \\ &\rightarrow (4 : (1 \times 4)) : ((1:1) \times 6) \rightarrow (4:4) : ((1:1) \times 6) \rightarrow (4:4) : ((1 \times 6) : (1 \times 6)) \\ &\rightarrow (4:4) : (6 : (1 \times 6)) \rightarrow (4:4) : (6:6) \rightarrow ((4:4) + 6) : 6 \rightarrow (4 : (4 + 6)) : 6 \\ &\rightarrow (4 : (1 : 0)) : 6 \rightarrow ((4 + 1) : 0) : 6 \rightarrow (5 : 0) : 6 \end{aligned}$$

(c) Computing a normal form $2 + (4 + (6 + (4 + 2)))$ with the rules in Example 5 requires 20 steps. To shorten the notation, we write $s^n(t)$ for $s(\dots(s(t))\dots)$ with n occurrences of s . Here t stands for an

arbitrary expression. We further abbreviate $s^n(0)$ to n . An example computation is shown below:

$$\begin{aligned}
& a(2, a(4, a(6, a(4, 2)))) \rightarrow s(a(1, a(4, a(6, a(4, 2))))) \rightarrow s^2(a(0, a(4, a(6, a(4, 2))))) \\
& \rightarrow s^2(a(4, a(6, a(4, 2)))) \rightarrow s^3(a(3, a(6, a(4, 2)))) \rightarrow s^4(a(2, a(6, a(4, 2)))) \\
& \rightarrow s^5(a(1, a(6, a(4, 2)))) \rightarrow s^6(a(0, a(6, a(4, 2)))) \rightarrow s^6(a(6, a(4, 2))) \\
& \rightarrow s^7(a(5, a(4, 2))) \rightarrow s^8(a(4, a(4, 2))) \rightarrow s^9(a(3, a(4, 2))) \rightarrow s^{10}(a(2, a(4, 2))) \\
& \rightarrow s^{11}(a(1, a(4, 2))) \rightarrow s^{12}(a(0, a(4, 2))) \rightarrow s^{12}(a(4, 2)) \rightarrow s^{13}(a(3, 2)) \\
& \rightarrow s^{14}(a(2, 2)) \rightarrow s^{15}(a(1, 2)) \rightarrow s^{16}(a(0, 2)) \rightarrow s^{18}(0)
\end{aligned}$$

To model multiplication on natural numbers in unary notation we add the following two rules:

$$m(0, y) \rightarrow 0 \qquad m(s(x), y) \rightarrow a(m(x, y), y)$$

Calculating 11×46 amounts to computing a normal form of $m(11, 46)$. The following computation has 2553 steps:

$$\begin{aligned}
& m(11, 46) \\
& \rightarrow a(m(10, 46), 46) \\
& \rightarrow a(a(m(9, 46), 46), 46) \\
& \rightarrow a(a(a(m(8, 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(m(7, 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(m(6, 46), 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(a(m(5, 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(a(a(m(4, 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(a(a(a(m(3, 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(a(a(a(a(m(2, 46), 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(a(a(a(a(a(m(1, 46), 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(a(a(a(a(a(a(m(0, 46), 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(m(0, 46), 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(a(46, 46), 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow \dots \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(a(92, 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow \dots \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(a(138, 46), 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow \dots \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(a(184, 46), 46), 46), 46), 46), 46), 46) \\
& \rightarrow \dots \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(a(230, 46), 46), 46), 46), 46), 46) \\
& \rightarrow \dots \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(a(276, 46), 46), 46), 46), 46), 46) \\
& \rightarrow \dots \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(a(322, 46), 46), 46), 46), 46) \\
& \rightarrow \dots \rightarrow a(a(a(a(a(a(a(a(a(a(a(a(a(368, 46), 46), 46) \\
& \rightarrow \dots \rightarrow a(a(414, 46), 46) \\
& \rightarrow \dots \rightarrow a(460, 46) \\
& \rightarrow \dots \rightarrow 506
\end{aligned}$$